

21 Lecture 04.13

On 04.13, we continued our study of deduction, and began by giving a classic “argument by cases.”

$\{(\forall x)Fx \vee (\forall x)Gx\}$ implies $(\forall x)(Fx \vee Gx)$.

{1}	(1) $(\forall x)Fx \vee (\forall x)Gx$	P
{2}	(2) $(\forall x)Fx$	P
{2}	(3) Fx	(2) UI
{2}	(4) $Fx \vee Gx$	(3) TF
{2}	(5) $(\forall x)(Fx \vee Gx)$	(4) UG
{}	(6) $(\forall x)Fx \supset (\forall x)(Fx \vee Gx)$	{2}(5) D
{7}	(7) $(\forall x)Gx$	P
{7}	(8) Gx	(7) UI
{7}	(9) $Fx \vee Gx$	(8) TF
{7}	(10) $(\forall x)(Fx \vee Gx)$	(9) UG
{}	(11) $(\forall x)Gx \supset (\forall x)(Fx \vee Gx)$	{7}(10) D
{1}	(12) $(\forall x)(Fx \vee Gx)$	(1, 6, 11) TF

We followed with an example of argument by *reductio ad absurdum*, that, in addition, illustrated use of the “conversion of quantifiers” rule, also known as “driving a negation across a quantifier.”

$(\exists y)(Py \supset (\forall x)Px)$ is valid

{1}	(1) $\neg(\exists y)(Py \supset (\forall x)Px)$	P
{1}	(2) $(\forall y)\neg(Py \supset (\forall x)Px)$	(1) CQ
{1}	(3) $\neg(Py \supset (\forall x)Px)$	(2) UI
{1}	(4) Py	(3) TF
{1}	(5) $(\forall x)Px$	(4) UG
{1}	(6) $\neg(\forall x)Px \wedge (\forall x)Px$	(3)(5) TF
{}	(7) $\neg(\exists y)(Py \supset (\forall x)Px) \supset$ $(\neg(\forall x)Px \wedge (\forall x)Px)$	{1}(6) D
{}	(8) $(\exists y)(Py \supset (\forall x)Px)$	(7) TF

We next gave the first of a pair of deductions that legitimate extending the “conversion of quantifiers rule” to allow passing directly from $\neg(\forall x)S$ to $(\exists x)\neg S$ and *vice versa*. I include the second here, though we did not cover it in class.

$\{\neg(\forall x)Fx\}$ implies $(\exists x)\neg Fx$.

{1}	(1) $\neg(\forall x)Fx$	P
{2}	(2) $\neg(\exists x)\neg Fx$	(1) P
{2}	(3) $(\forall x)\neg\neg Fx$	(2) CQ
{2}	(4) $\neg\neg Fx$	(3) UI
{2}	(5) Fx	(4) TF
{2}	(6) $(\forall x)Fx$	(5) UG
{1, 2}	(7) $(\forall x)Fx \wedge \neg(\forall x)Fx$	(6) TF
{1}	(8) $\neg(\exists x)\neg Fx \supset ((\forall x)Fx \wedge \neg(\forall x)Fx)$	{2}(7) D
{1}	(9) $(\exists x)\neg Fx$	(8) TF

$\{(\exists x)\neg Fx\}$ implies $\neg(\forall x)Fx$.

{1}	(1) $(\forall x)Fx$	P
{2}	(2) $(\exists x)\neg Fx$	(1) P
{1}	(3) Fx	(1) UI
{1}	(4) $\neg\neg Fx$	(3) TF
{1}	(5) $(\forall x)\neg\neg Fx$	(4) UG
{1}	(6) $\neg(\exists x)\neg Fx$	(5) CQ
{1, 2}	(7) $\neg(\exists x)\neg Fx \wedge (\exists x)\neg Fx$	(6) TF
{2}	(8) $(\forall x)Fx \supset (\neg(\exists x)\neg Fx \wedge (\exists x)\neg Fx)$	{1}(7) D
{1}	(9) $\neg(\forall x)Fx$	(8) TF