

CNF Satisfiability

CNF Boolean Formula

- Variables, x, y, z
- Literals, positive and negative variables
- Clauses, disjunctions of literals
- CNF formula, a conjunction of clauses

A formula is *satisfiable* if there is some truth assignment such that the formula evaluates to true.



Satisfiability examples

•
$$(x \vee \neg x)$$

- (x) ^ (¬ x)
- $(x \lor \neg x) \land (y \lor z)$
- $(x \lor y) \land (\neg y \lor z)$
- $(x \lor y) \land (\neg y \lor z) \land (\neg y \lor \neg z)$ $\land (y \lor \neg z) \land (\neg y \lor x)$

Resolution based satisfiability

 $(x \lor y) \land (\neg x \lor z)$ is satisfiable when $(y \lor z)$ is satisfiable

Use this idea to create resolution based theorem prover:

- Start with a set of clauses S.
- Find two clauses to resolve. (x \lor y1 ..) & (\neg y \lor z1 ...)
- If the result (y1 .. V z1 ...) is empty, then the formula is unsatisfiable.
- If the result is a tautology, ignore it.
- Otherwise add result to S, and continue until no new information can be derived.

The Formula

- 1. D V $\neg AE$ V $\neg U$ 2. $\neg R \lor \neg D$ 3. $\neg B \lor \neg O \lor G$ 4. $\neg C \lor \neg AD \lor \neg A$ 5. $\neg T \lor AE \lor E$ 6. C V Q V \neg E 7. T V \neg B V \neg M 8. AD V \neg N V U 9. $\neg G \lor A \lor S$
- 10. D V \neg G V M 11. AE $\vee \neg P \vee X$ 12. $\neg N \lor \neg D \lor O$ 13. Y V C V \neg X 14. $\neg N \lor Y \lor R$ 15. $S \vee G \vee N$ 16. ¬T V ¬Y 17. ¬AE ∨ P 18. X V \neg Y

Activity

- 1. Find your clause (other people with same #)
- 2. With your clause, find another clause with complementary literal
- 3. RESOLVE!
- 4. Resolved literals out, sit and chat
- 5. If you find a contradiction, let me know!
- 6. If resolved clause is a tautology, let me know your clause number(s). Everyone out.
- 7. Otherwise, got back to step 2 until no more resolutions.

- Random formula has residual of
- AE $\vee \neg P \vee C \vee W \vee \neg T$
- $\neg G \lor \neg O \lor \neg T$ (never resolves)