Continuous Program Reasoning via Differential Bayesian Inference

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Abstract

Programs often evolve continuously by integrating changes from multiple programmers. The effective adoption of program analysis tools in this continuous integration setting is hindered by the need to only report alarms relevant to a particular program change. We present a probabilistic framework, Drake, to apply program analyses to continuously evolving programs. Drake is applicable to a broad range of analyses that are based on deductive reasoning. The key insight underlying Drake is to compute a graph that concisely and precisely captures differences between the derivations of alarms produced by the given analysis on the program before and after the change. Performing Bayesian inference on the graph thereby enables to rank alarms by relevance to the change. We evaluate Drake using Sparrow—a static analyzer that targets buffer-overflow, format-string, and integer-overflow errors—on a suite of ten widely-used C programs each comprising 13k–112k lines of code. Drake enables to discover all true bugs by inspecting only 30 alarms per benchmark on average, compared to 85 (3X more) alarms by the same ranking approach in batch mode, and 118 (4X more) alarms by a differential approach based on syntactic masking of alarms which also misses 4 of the 26 bugs overall.

1 Introduction

The application of program analysis tools such as Astrée [5], SLAM [2], Coverity [4], FindBugs [42] and Infer [7] to large software projects has highlighted research challenges at the intersection of program reasoning theory and software engineering practice. An important aspect of long-lived, multi-developer projects is the practice of continuous integration, where the codebase evolves through multiple versions which are separated by incremental changes. In this context, programmers are typically less worried about the possibility of bugs in existing code—which has been in active use in the field—and in parts of the project which are unrelated to their immediate modifications. They specifically want to know whether the present commit introduces new bugs, regressions, or breaks assumptions made by the rest of the codebase [4, 39, 45]. How do we determine whether a static analysis alarm is relevant for inspection given a small change to a large program?

A common approach is to suppress alarms that have already been reported on previous versions of the program [4, 11, 13]. Unfortunately, such syntactic masking of alarms has a great risk of missing bugs, especially when the commit modifies code in library routines or in commonly used helper methods, since the new code may make assumptions that are not satisfied by the rest of the program [34]. Therefore, even alarms previously reported and marked as false positives may potentially need to be inspected again.

In this paper, we present a probabilistic framework to apply program analyses to continuously evolving programs. The framework, called Drake, has to address four key challenges to be effective. First, it must overcome the limitation of syntactic masking by reasoning about semantic changes that produce alarms. For this purpose, it employs derivations of alarms—logical chains of cause-and-effect—produced by the given analysis on the program before and after the change. Such derivations are naturally obtained from analyses whose reasoning can be expressed or instrumented in the form of deductive rules. As such, Drake is applicable to a broad range of analyses, including those commonly specified in the logic programming language Datalog [6, 46].

Second, Drake must relate abstract states of the two program versions which do not share a common vocabulary. We build upon previous syntactic program differencing work by setting up a matching function which maps source locations, variable names, and other syntactic entities of the old version of the program to the corresponding entities of the new version. The matching function allows to not only relate alarms but also the derivations that produce them.

Third, Drake must efficiently and precisely compute the relevance of each alarm to the program change. For this purpose, it constructs a differential derivation graph that captures differences between the derivations of alarms produced by the given analysis on the program before and after the change. For a fixed analysis, this graph construction takes
```c
#define CMP_SIZE 529200

void read_value_long(FILE *file, long *val) {
    char buf[5];
    fread(buf, 1, 4, file); // Input Source
    header_size = min(info->header_size, HEADER_SIZE);
    read_value_long(f, info->data_size);
    return info;
}

 unions.c


```
Figure 2. The Drake system. By applying the analysis to each version of the program, we start with the grounded constraints for each version. The syntactic matching function $\delta$ allows us to compare the constraints and derivation trees of $P_{\text{old}}$ and $P_{\text{new}}$, which we merge to obtain the differential derivation graph $GC_\Delta$. We treat the resulting structure as a probabilistic model, and interactively reprioritize the list of alarms as the user triages them and labels their ground truth.

The first alarm concerns the command line option "t", which trims periods of silence from the ends of an audio file. The program reads unsanitized data into the field $info->header_size$ at line 25, and allocates a buffer of proportional size at line 30. Drake observes this data flow, concludes that the multiplication could overflow, and subsequently raises an alarm at the allocation site. However, this data has been sanitized at line 29, so that the expression $header_size * \text{sizeof(char)}$ cannot overflow. This is therefore a false alarm in both $P_{\text{old}}$ and $P_{\text{new}}$.

The second alarm is triggered by the command line option "c", which compares the contents of two audio files. The first version has source-sink flows from the untrusted fields $info1->data_size$ and $info2->data_size$, but this is a false alarm since the value of bytes cannot be larger than $\text{CMP_SIZE}$. On the other hand, the new version of the program includes an option to offset the contents of one file by $\text{shift_secs}$ seconds. This value is used without sanitization to compute $\text{cmp_size}$, leading to a possible integer overflow at line 42, which would then result in a buffer of unexpected size being allocated at line 45. Thus, while Drake raises an alarm at the same allocation site for both versions of the program, this is a false alarm in $P_{\text{old}}$ but a real bug in $P_{\text{new}}$.

We now restate the central question of this paper: How do we alert the user to the possibility of a bug at line 45, while not forcing them to inspect all the alarms of the "batch mode" analysis, such as that at line 30?

Figure 2 presents an overview of our system, Drake. First, the system extracts static analysis results from both the old and new versions of the program. Since these results are described in terms of syntactic entities (such as source locations) from different versions of the program, it uses a syntactic matching function $\delta$ to translate the old version of the constraints into the setting of the new program. Drake then merges the two sets of constraints into a unified differential derivation graph. These differential derivations highlight the relevance of the changed code to the static analysis alarms. Moreover, the differential derivation graph also enables us to perform marginal inference with the feedback from the user as well as previously labeled alarms from the old version.

We briefly explain the reasoning performed by Sparrow in Section 2.1, and explain our ideas in Sections 2.2–2.3.

2.1 Reflecting on the Integer Overflow Analysis

Sparrow detects harmful integer overflows by performing a flow-, field-, and context-sensitive taint analysis from untrusted data sources to sensitive sinks [15]. While the actual implementation includes complex details to ensure performance and accuracy, it can be approximated by inference rules such as those shown in Figure 3.

The input tuples indicate elementary facts about the program which the analyzer determines from the program text. For example, the tuple $\text{DUEdge}(7, 9)$ indicates that there is a one-step data flow from line 7 to line 9 of the program. The inference rules, which we express here as Datalog programs, provide a mechanism to derive new conclusions about the program being analyzed. For example, the rule $r_2$, $\text{DUPath}(c_1, c_3) \leftarrow \text{DUPath}(c_1, c_2), \text{DUEdge}(c_2, c_3)$, indicates that for each triple $(c_1, c_2, c_3)$ of program points, whenever there is a multi-step data flow from $c_1$ to $c_2$ and an immediate data flow from $c_2$ to $c_3$, then there may be a multi-step data flow from $c_1$ to $c_3$. Starting from the input tuples, we repeatedly apply these inference rules to reach new conclusions, until we reach a fixpoint. This process may be visualized...
Input relations
DUEdge(c1, c2) : Immediate data flow from c1 to c2
Src(c) : Origin of potentially erroneous traces
Dst(c) : Potential program crash point

Output relations
DUPath(c1, c2) : Transitive data flow from c1 to c2
Alarm(c) : Potentially erroneous trace reaching c

Analysis rules
r1 : DUPath(c1, c2) => DUEdge(c1, c2).
r2 : DUPath(c1, c3) => DUPath(c1, c2), DUEdge(c2, c3).
r3 : Alarm(c2) => DUPath(c1, c2), Src(c1), Dst(c2).

Figure 3. Approximating a complex taint analysis with simple inference rules. All variables c1, c2, etc. range over the set of program points.

as discovering the nodes of a derivation graph such as that shown in Figure 4.

We use derivation graphs to determine alarm relevance. As we have just shown, such derivation graphs can be naturally described by inference rules. These inference rules are straightforward to obtain if the analysis is written in a declarative language such as Datalog. If the analysis is written in a general-purpose language, we define a set of inference rules that approximate the reasoning processes of the original analyzer. The degree of approximation does not affect the accuracy of the analysis but only affects the accuracy of subsequent probabilistic reasoning. Furthermore, in practice, it requires only a small amount of effort to implement by instrumenting the original analyzer. We explain this instrumentation for a general class of analyses in Section 4.2.

2.2 Classifying Derivations by Alarm Transfer

Traditional approaches such as syntactic alarm masking will de-prioritize both Alarm(30) and Alarm(45) as they occur in both versions of the program. Concretely then, our problem is to provide a mechanism by which to continue to de-prioritize Alarm(30), but highlight Alarm(45) as needing re-inspection.

Translating clauses. For each grounded clause g in the derivation from the new program Pnew, we can ask whether g also occurs in the old program Pold. For example, the clauses in Figure 4(a) commonly exist in both of the versions, but the clauses in Figure 4(c) are only present in Pnew. Such questions presuppose the existence of some correspondence between program points, variables, functions, and other syntactic entities of Pold, and the corresponding entities of Pnew. In Section 4.3, we will construct a matching function δ to perform this translation, but for the purpose of this example, it can be visualized as simply being a translation between line numbers, such as that obtained using diff.

Translating derivation trees. The graph of Figure 4 can be viewed as encoding a set of derivation trees for each alarm. A derivation tree is an inductive structure which culminates in the production of a tuple t. It is either: (a) an input tuple, or (b) a grounded clause t1 ∧ t2 ∧ · · · ∧ tk => t together with a derivation tree τi for each antecedent tuple ti.

Let us focus on two specific derivation trees from this graph: first, the sequence τ30 in Figure 4(a):

DUPath(7, 9) => DUPath(7, 18) => · · · => Alarm(30),

and second, the sequence τ45 in Figure 4(c):

DUPath(54, 42) => DUPath(54, 43) => · · · => Alarm(45),

and where each sequence is supplemented with appropriate input tuples. Observe that each clause of the first tree, τ30, is common to both Pold and Pnew. More generally, every derivation tree of Alarm(30) from Pnew is already present in Pold. As a result, Alarm(30) is unlikely to represent a real bug. On the other hand, the second tree, τ45, exclusively occurs in the new version of the program. Therefore, since there are more reasons to suspect the presence of a bug at Alarm(45) in Pnew than in Pold, we conclude that it is necessary to re-inspect this alarm.

The first step to identifying relevant alarms is therefore to determine which alarms have new derivation trees. As we show in Figure 5, where the new t2 => t3 derivation for t3 transitively extends to t4, this question inherently involves non-local reasoning. Other approaches based on enumerating derivation trees by exhaustive unrolling of the fixpoint graph will fail in the presence of loops, i.e., when the number of derivation trees is infinite. For a fixed analysis, we will now describe a technique to answer this question in time linear in the size of the new graph.

The differential derivation graph. From Figures 4(b) and 4(c), consider the tuple DUPath(7, 39) which is the starting point of the diverged derivations of Alarm(45) between the two versions. Say we want to know whether it itself has at least one derivation tree exclusive to Pnew. Observe that only one clause g = r2(7, 19, 39) derives this tuple:

DUPath(7, 19) ∧ DUEdge(19, 39) => r2 DUPath(7, 39). (1)

Since the clauses are computed to fixpoint, the only way for a derivation tree to be exclusive to Pnew is by using a newly appearing derivation of either DUPath(7, 19) or DUEdge(19, 39).

A derivation tree τ is either an input tuple t or a grounded clause t1 ∧ t2 ∧ · · · ∧ tk => τ applied to a set of smaller derivation trees τ1, τ2, . . . , τk. If τ is an input tuple, then it is exclusive to the new analysis run if it did not appear in the old program. In the inductive case, τ is exclusive to the new analysis run if τi is a new derivation tree for some i.

Our key insight is that we can therefore separate these derivation trees by splitting each tuple t itself into two variants, τα and τβ. We set this up so that the derivations of τα
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Figure 4. Portions of the old and new derivation graphs by which the analysis identifies suspicious source-sink flows in the two versions of the program. The numbers indicate line numbers of the corresponding code in Figure 1. Nodes corresponding to grounded clauses, such as \( r_1(7, 9) \), indicate the name of the rule and the instantiation of its variables, i.e., \( r_1 \) with \( c_1 = 7 \) and \( c_2 = 9 \). Notice that in the new derivation graph the analysis has discovered two suspicious flows—from lines 7 and 54 respectively—which both terminate at line 45.

Figure 5. Naive approaches, based on tree or graph differences, may fail to recognize the presence of new derivation tree for \( \Delta t_3 \): \( t_2 \to t_3 \to t_4 \).

Figure 6. Differentiating the clause \( g \) from Equation 1.

Observe that the only way to derive \( \text{DUPath}_{\alpha}(7, 39) \) is by applying a clause to a set of tuples all of which are themselves the \( \alpha \)-variety. The use of even a single \( \beta \)-variant tuple always results in the production of \( \text{DUPath}_{\beta}(7, 39) \).

We visualize this process in Figure 6. By similarly splitting each clause \( g \) of the analysis fixpoint, we produce the clauses of the differential derivation graph \( GC_\Delta \).

At the base case, let the set of merged input tuples \( I_\Delta \) be the \( \alpha \)-variants of input tuples which occur in common, and the \( \beta \)-variants of all input tuples which only occur in \( P_{\text{new}} \). We wanted to know whether \( \text{DUPath}(7, 39) \) has any new derivation trees. In the differential derivation graph, we can answer this question by considering the derivability of \( \text{DUPath}_{\beta}(7, 39) \), which can be derived by any of the three clauses \( g_{\alpha \beta}, g_{\beta \alpha}, g_{\beta \beta} \) in Equations 3–5. Since input tuple \( \text{DUEdge}(19, 39) \) occurs in both versions, \( \text{DUEdge}(19, 39) \notin I_\Delta \) is not derivable, so the only way to produce \( \text{DUPath}_{\beta}(7, 39) \) is by using the clause \( g_{\beta \beta} \) which in turn involves producing \( \text{DUPath}_{\beta}(7, 19) \). By a similar reasoning process, we find \( \text{DUPath}_{\beta}(7, 19) \) to be underviable. We therefore conclude that \( \text{DUPath}_{\beta}(7, 39) \) is underviable in the merged graph.
which coincides with our intuition that the original tuple DUPath(7, 39) has no new derivation trees in $P_{\text{new}}$.

2.3 A Probabilistic Model of Alarm Relevance

We build our system on the idea of highlighting alarms $\text{Alarm}(c)$ whose $\beta$-variants, $\text{Alarm}_\beta(c)$, are derivable in the differential derivation graph. By leveraging recent work on probabilistic alarm ranking, we can also transfer feedback across program versions and highlight alarms which are both relevant and likely to be real bugs. The idea is that since alarms share root causes and intermediate tuples, labelling one alarm as true or false should change our confidence in closely related alarms.

Differential derivation graphs, probabilistically. The inference rules of the analysis are frequently designed to be sound, but deliberately incomplete. Let us say that a rule misfires if it takes a set of true hypotheses, and produces an output tuple which is actually false. In practice, in large real-world programs, rules misfire in statistically regular ways.

We therefore associate each rule $r$ with the probability $p_r$ of its producing valid conclusions despite assuming valid hypotheses.

Consider the rule $r_2$, and its instantiation as the grounded clause, $g_{\alpha\beta} = r_2(t_1, t_2)$, with $t_1 = \text{DUPath}_\alpha(7, 19)$ and $t_2 = \text{DUEdge}_\beta(19, 39)$ as its antecedent tuples, and with $t_3 = \text{DUPath}_\beta(7, 39)$ as its conclusion. We define:

$$\Pr(t_3 \mid t_1 \wedge t_2) = p_{r_2}, \quad (6)$$

$$\Pr(t_3 \mid \neg t_1 \vee \neg t_2) = 0, \quad (7)$$

so that $t_3$ is true only if $t_1$ and $t_2$ are both true, and even in that case, only with probability $p_{r_2}$.

Since multiple alarms share portions of their derivation trees, labelling $\text{Alarm}(c)$ as $v = \text{True}$ or $v = \text{False}$ gives us additional information about the truth and relevance of the remaining alarms, which the Bayesian network automatically incorporates into our updated probabilities, $\Pr(\text{Alarm}(c') \mid \text{Alarm}(c) = v)$.

Interaction Model. Drake presents the user with a list of alarms, sorted according to $\Pr(\text{Alarm}(c) \mid e)$, i.e., the probability that $\text{Alarm}(c)$ is both relevant and a true bug, conditioned on the current feedback set $e$. After each round of user feedback, we update $e$ to include the user label for the last triaged alarm, and rerank the remaining alarms according to $\Pr(\text{Alarm}(c) \mid e)$.

Furthermore, $e$ can also be initialized by applying any feedback that the user has provided to the old program, pre-commit, say to $\text{Alarm}(45)$, to the old versions of the corresponding tuples in $\text{GC}_\alpha$, i.e., to $\text{Alarm}_\alpha(45)$. We note that this combination of differential relevance computation and probabilistic generalization of feedback is dramatically effective in practice: while the original analysis produces an average of 563 alarms in each of our benchmarks, after relevance-based ranking, the last real bug is at rank 94; the initial feedback transfer reduces this to rank 78, and through the process of interactive reranking, all true bugs are discovered within just 30 rounds of interaction on average.

3 A Framework for Alarm Transfer

We formally describe the Drake workflow in Algorithm 1, and devote this section to our core technical contributions: the constraint merging algorithm $\text{Merge}$ in step 3 and enabling feedback transfer in step 5. We begin by setting up preliminary details regarding the analysis and reviewing the use of Bayesian inference for interactive alarm ranking.

**Algorithm 1** DRAKE($P_{\text{old}}, P_{\text{new}}$), where $\mathcal{A}$ is an analysis, and $P_{\text{old}}$ and $P_{\text{new}}$ are the old and new versions of the program to be analyzed.

1. Compute $\mathbf{R}_{\text{old}} = \mathcal{A}(P_{\text{old}})$ and $\mathbf{R}_{\text{new}} = \mathcal{A}(P_{\text{new}})$. Analyze both programs.

2. Define $\mathbf{R}_3 = \delta(\mathbf{R}_{\text{old}})$. Translate the analysis results and feedback from $P_{\text{old}}$ to the setting of $P_{\text{new}}$.

3. Compute the differential derivation graph:

$$\mathbf{R}_\Delta = \text{Merge}(\mathbf{R}_3, \mathbf{R}_{\text{new}}).$$

4. Pick a bias $\epsilon$ according to Section 3.2 and convert $\mathbf{R}_\Delta$ into a Bayesian network, $\text{BNET}(\mathbf{R}_\Delta)$. Let $\Pr$ be its joint probability distribution.

5. Initialize the feedback set $e$ according to the chosen feedback transfer mode (see Section 3.3).

6. While there exists an unlabelled alarm:

a. Let $A_u$ be the set of unlabelled alarms.

b. Present the highest probability unlabelled alarm for user inspection:

$$a = \arg \max_{a \in A_u} \Pr(a_{\beta} \mid e) .$$

If the user marks it as true, update $e := e \land a_{\beta}$.

Otherwise update $e := e \land \neg a_{\beta}$.

3.1 Preliminaries

**Declarative program analysis.** Drake assumes that the analysis result $\mathcal{A}(P)$ is a tuple, $\mathbf{R} = (I, C, A, GC)$, where $I$ is the set of input facts, $C$ is the set of output tuples, $A$ is the set of alarms, and $GC$ is the set of grounded clauses which connect them. We obtain $I$ by instrumenting the original analysis $(A, I) = \mathcal{A}_{\text{orig}}(P)$. For example, in our experiments, Sparrow outputs all immediate dataflows, DUEdge($c_1, c_2$) and potential source and sink locations, Src($c$) and Dst($c$). We obtain $C$ and $GC$ by approximating the analysis with a Datalog program.
A Datalog program [1]—such as that in Figure 3—consumes a set of input relations and produces a set of output relations. Each relation is a set of tuples, and the computation of the output relations is specified using a set of rules. A rule \( r \) is an expression of the form \( R_h(u_1) \leftarrow R_1(v_1), R_2(v_2), \ldots, R_k(u_k) \), where \( R_1, R_2, \ldots, R_k \) are relations, \( R_h \) is an output relation, \( u_1, v_2, \ldots, u_k \) and \( v_h \) are vectors of variables of appropriate arity. The rule \( r \) encodes the following universally quantified logical formula: "For all values of \( u_1, v_2, \ldots, u_k \) and \( v_h \), if \( R_1(v_1) \land R_2(v_2) \land \cdots \land R_k(u_k) \), then \( R_h(v_h) \).

To evaluate the Datalog program, we initialize the set of conclusions \( C \) := \( \emptyset \) and the set of grounded clauses \( GC := \emptyset \), and repeatedly instantiate each rule to add tuples to \( C \) and grounded clauses to \( GC \); i.e., whenever \( R_1(c_1), R_2(c_2), \ldots, R_k(c_k) \in C \), we update \( C := C \cup \{ R_5(c_k) \} \) and \( GC := GC \cup \{ R_1(c_1) \land R_2(c_2) \land \cdots \land R_k(c_k) \implies r \land R_h(c_k) \} \).

For each grounded clause \( g \) of the form \( H_g \implies c_g \), we refer to \( H_g \) as the set of antecedents of \( g \), and \( c_g \) as its conclusion.

We repeatedly add tuples to \( C \) and grounded clauses to \( GC \) until a fixpoint is reached.

Bayesian alarm ranking. The main observation behind Bayesian alarm ranking [43] is that alarms are correlated in their ground truth: labelling one alarm as true or false should change our confidence in the tuples involved in its production, and transitively, affect our confidence in a large number of other related alarms. Concretely, these correlations are encoded by converting the set of grounded clauses \( GC \) into a Bayesian network: we will now describe this process.

Let \( G \) be the derivation graph formed by all tuples \( t \in C \) and grounded clauses \( g \in GC \). Figure 4 is an example. Consider a grounded clause \( g \in GC \) of the form \( t_1 \land t_2 \land \cdots \land t_k \implies r \). Observe that \( g \) requires \( a \) all its antecedents to be true to be able to successfully derive its output tuple. In particular, if any of the antecedents fails, then the clause is definitely inoperative. Let us assume a function \( p \) which maps each rule \( r \) to the probability of its successful firing, \( p_r \).

Then, we associate \( g \) with the following conditional probability distribution (CPD) using an assignment \( P \):

\[
P(g \mid t_1 \land t_2 \land \cdots \land t_k) = p_r, \text{ and } P(g \mid \neg(t_1 \land t_2 \land \cdots \land t_k)) = 0.
\]

The conditional probabilities of an event and its complement sum to one, so that \( Pr(\neg g \mid t_1 \land t_2 \land \cdots \land t_k) = 1 - p_r \) and \( Pr(\neg g \mid \neg(t_1 \land t_2 \land \cdots \land t_k)) = 1 \).

On the other hand, consider some tuple \( t \) which is produced by the clauses \( g_1, g_2, \ldots, g_l \). If there exists some clause \( g_j \) which is derivable, then \( t \) is itself derivable. If none of the clauses is derivable, then neither is \( t \). We therefore associate \( t \) with the CPD for a deterministic disjunction:

\[
P(t \mid g_1 \lor g_2 \lor \cdots \lor g_l) = 1, \text{ and } P(t \mid \neg(g_1 \lor g_2 \lor \cdots \lor g_l)) = 0.
\]

Let us also assume a function \( p_{in} \) which maps input tuples \( t \) to their prior probabilities. In the simplest case, input tuples are known with certainty, so that \( p_{in}(t) = 1 \). In Section 3.2, we will see that the choice of \( p_{in} \) allows us to uniformly generalize both relevance-based and traditional batch-mode ranking. We define the CPD of each input tuple \( t \) as:

\[
P(t) = p_{in}(t).
\]

By definition, a Bayesian network is a pair \((G, P)\), where \( G \) is an acyclic graph and \( P \) is an assignment of CPDs to each node [23]. We have already defined the CPDs in Equations 9–13; the challenge is that the derivation graph \( G \) may have cycles. Raghothaman et al. [43] presents an algorithm to extract an acyclic subgraph \( G_c \subseteq G \) which still preserves derivability of all tuples. Using this, we may define the final Bayesian network, \( BN(R) = (G_c, P) \).

3.2 The Constraint Merging Process
As motivated in Section 2.2, we combine the constraints from the old and new analysis runs into a single differential derivation graph \( R_{\Delta} \). Every derivation tree \( r \) of a tuple from \( R_{new} \) is either common to both \( R_S \) and \( R_{new} \), or is exclusive to the new analysis run.

Recall that a derivation tree is inductively defined as either:

(a) an individual input tuple, or
(b) a grounded clause \( t_1 \land t_2 \land \cdots \land t_k \implies r \) together with derivation trees \( t_1, t_2, \ldots, t_k \) for each of the antecedent tuples. Since the grounded clauses are collected until fixpoint, the only way for a derivation tree to be exclusive to the new program is if it is either: (a) a new input tuple \( t \in I_{new} \setminus I_S \), or (b) a clause \( t_1 \land t_2 \land \cdots \land t_k \implies r \) with a new derivation tree for at least one child \( t_i \).

The idea behind the construction of \( R_{\Delta} \) is therefore to split each tuple \( t \) into two variants, \( t_a \) and \( t_b \), where \( t_a \) precisely captures the common derivation trees and \( t_b \) exactly captures the derivation trees which only occur in \( R_{new} \). We formally describe its construction in Algorithm 2. Theorem 3.1 is a straightforward consequence.

Theorem 3.1 (Separation). Let the combined analysis results from \( P_{old} \) and \( P_{new} \) be \( R_{\Delta} = Merge(R_S, R_{new}) \). Then, for each tuple \( t \),

1. \( t_a \) is derivable from \( R_{\Delta} \) iff \( t \) has a derivation tree which is common to both \( R_S \) and \( R_{new} \), and
2. \( t_b \) is derivable from \( R_{\Delta} \) iff \( t \) has a derivation tree which is absent from \( R_S \) but present in \( R_{new} \).

Proof: In each case, by induction on the tree which is given to exist. All base cases are all immediate. We will now explain the inductive cases.

Of part 1, in the \( \Rightarrow \) direction. Let \( t_a \) be the result of a clause \( t'_1 \land t'_2 \land \cdots \land t'_k \implies r \). By construction, it is the case that each \( t'_i \) is of the form \( t_{i,a} \), and by IH, it must already have a derivation tree \( r_i \) which is common to both analysis results. It follows that \( t_a \) also has a derivation tree \( r(t_1, t_2, \ldots, t_k) \) in common to both results.
Distinguishing abstract derivations. One detail is that since the output tuples indicate program behaviors in the abstract domain, it may be possible for \( P_{\text{new}} \) to have a new concrete behavior, while the analysis continues to produce the same set of tuples. This could conceivably affect ranking performance by suppressing real bugs in \( R_\Delta \). Therefore, instead of using \( I_\Delta \) as the set of input tuples in BINET(\( R_\Delta \)), we use the set of all input tuples \( t \in \{ t_a, t_\beta \mid t \in I_{\text{new}} \} \), with prior probability: if \( t \in I_{\text{new}} \setminus I_\delta \), then \( p_\text{inc}(t_\alpha) = 1 - p_\text{inc}(t_\beta) = 1.0 \), and otherwise, if \( t \in I_{\text{new}} \cap I_\delta \), then \( p_\text{inc}(t_\alpha) = 1 - p_\text{inc}(t_\beta) = \epsilon \). Here, \( \epsilon \) is our belief that the same abstract state has new concrete behaviors. The choice of \( \epsilon \) also allows us to interpolate between purely change-based (\( \epsilon = 0 \)) and purely batch-mode ranking (\( \epsilon = 1 \)).

3.3 Bootstrapping by Feedback Transfer

It is often the case that the developer has already inspected some subset of the analysis results on the program from before the code change. By applying this old feedback \( e_{\text{old}} \) to the new program, as we will now explain, the differential derivation graph also allows us to further improve the alarm rankings beyond just the initial estimates of relevance.

Conservative mode. Consider some negatively labelled alarm \( a \in e_{\text{old}} \). The programmer has therefore indicated that all of its derivation trees in \( R_{\text{old}} \) are false. If \( a' \equiv \delta(a) \), since the derivation trees of \( a' \) in \( R_\Delta \) correspond to a subset of the derivation trees of \( a \) in \( R_{\text{old}} \), we can additionally deprioritize these derivation trees by initializing:

\[
e := \{-a_a \mid \forall \text{negative labels } -a_a \in \delta(e_{\text{old}})\}.
\]

Strong mode. In many cases, programmers have a lot of trust in \( P_{\text{old}} \) since it has been tested in the field. We can then make the strong assumption that \( P_{\text{old}} \) is bug-free, and extend inter-version feedback transfer, by initializing:

\[
e := \{-a_a \mid \forall a \in A_\Delta\}.
\]

Our experiments in Section 5 are primarily conducted with this setting.

Aggressive mode. Finally, if the programmer is willing to accept a greater risk of missed bugs, then we can be more aggressive in transferring inter-version feedback:

\[
e := \{-a_a, -a_\beta \mid \forall a \in A_\Delta\}.
\]

In this case, we not only assume that all common derivations of the alarms are false, but also additionally assume that the new alarms are false. It may be thought of as a combination of syntactic alarm masking and Bayesian alarm prioritization. We also performed experiments with this setting and, as expected, observed that it misses 4 real bugs (15%), but additionally reduces the average number of alarms to be inspected before finding all true bugs from 30 to 22.

Algorithm 2 MERGE(\( R_\Delta, R_{\text{new}} \)), where \( R_\Delta \) is the translated analysis result from \( P_{\text{old}} \) and \( R_{\text{new}} \) is the result from \( P_{\text{new}} \).

1. Unpack the input-, output-, alarm tuples, and grounded clauses from each version of the analysis result. Let \( (I_\Delta, C_\Delta, A_\Delta, G_\Delta) = R_\Delta \) and \( (I_{\text{new}}, C_{\text{new}}, A_{\text{new}}, G_{\text{new}}) = R_{\text{new}} \).
2. Form two versions, \( t_a, t_\beta \), of each output tuple in \( R_{\text{new}} \):
   \[ C_\Delta = \{ t_a, t_\beta \mid t \in C_{\text{new}} \}, \]
   \[ A_\Delta = \{ t_a, t_\beta \mid t \in A_{\text{new}} \}. \]
3. Classify the input tuples into those which are common to both versions and those which are exclusively new:
   \[ I_\Delta = \{ t_a \mid t \in I_{\text{new}} \cap I_\delta \} \cup \{ t_\beta \mid t \in I_{\text{new}} \setminus I_\delta \}. \]
4. Populate the clauses of \( G_\Delta \): For each clause \( G \) in \( C_{\text{new}} \) of the form \( t_1 \land t_2 \land \cdots \land t_k \Rightarrow t_h \), and for each \( H_g' \in \{ (t_1, t_\beta) \times (t_2, t_\beta) \times \cdots \times (t_k, t_\beta) \} \) for exactly one \( i \), so that \( t_i \) has an exclusively new derivation tree \( t_i \).
   a. If \( H'_g = (t_1, t_2, \ldots, t_{i-1}, t_a, t_{i+1}, \ldots, t_k) \) consists entirely of “\( a \)”-tuples, produce the clause:
   \[ H'_g \Rightarrow t_h].\]
   b. Otherwise, if there is at least one “\( \beta \)”-tuple, then emit the clause:
   \[ H'_g \Rightarrow t_h].\]
5. Output the merged result \( R_\Delta = (I_\Delta, C_\Delta, A_\Delta, G_\Delta) \).
4 Implementation

In this section, we discuss key implementation aspects of 
Drake, in particular: (a) extracting derivation trees from pro-
gram analyzers that are not necessarily written in a declara-
tive language, and (b) comparing two versions of a program.
In Section 4.2, we explain how we extract derivation trees
from complex, black-box static analyses, while Section 4.3
describes the syntactic matching function $\delta$ for a pair of
program versions.

4.1 Setting

We assume that the analysis is implemented on top of a
framework which is based on abstract interpretation [10] and sup-
ports relational as well as non-relational semantic properties
for various programming languages.

Program. A program is represented as a control flow graph
$(C, \rightarrow, c_0)$ where $C$ denotes the set of program points, $(\rightarrow)$
$\subseteq C \times C$ denotes the control flow relation, and $c_0$ is the entry
node of the program. Each program point is associated with a
command.

Program analysis. We target a class of analyses whose ab-
stract domain maps program points to abstract states:

$$D = C \rightarrow \mathbb{S}.$$ 

An abstract state maps abstract locations to abstract values:

$$\mathbb{S} = L \rightarrow \mathbb{V}.$$ 

The analysis produces alarms for each potentially erroneous
program points.

The data dependency relation $(\rightsquigarrow) \subseteq C \times L \times C$ is defined
as follows:

$$c_0 \rightsquigarrow c_n = \exists I \subseteq \{c_0, c_1, \ldots, c_n\} \subseteq \text{Paths}, \exists I \subseteq L, I \subseteq D(c_0) \cap \text{U}(c_n) \land \forall i \in (0, n). l \notin D(c_i)$$

where $D(c) \subseteq L$ and $U(c) \subseteq L$ denote the def and use sets
of abstract locations at program point $c$. A data dependency
$c_0 \rightsquigarrow c_n$ represents that abstract location $I$ is defined at
program point $c_0$ and used at $c_n$ through path $[c_0, c_1, \ldots, c_n]$, and no intermediate program points on the path re-define $I$.

4.2 Extracting derivation trees from complex,
non-declarative program analyses

To extract the Bayesian network, the analysis additionally
computes derivation trees for each alarm. In general, instru-
menting a program analyzer to do bookkeeping at each
reasoning step would impose a high engineering burden. We
instead abstract the reasoning steps using dataflow relations.
Figure 3 shows the relations and deduction rules to de-
scribe the reasoning steps of the analysis. Data flow relation
$DUEdge \subseteq C \times C$ which is a variant of data dependency [37]
is defined as follows:

$$DUEdge(c_0, c_n) = \exists I \subseteq L. c_0 \rightsquigarrow c_n.$$ 

A dataflow relation $DUEdge(c_0, c_n)$ represents that an ab-
stract location is defined at program point $c_0$ and used at $c_n$.
Relation $DUPath(c_1, c_n)$ represents transitive dataflow rela-
tion from point $c_1$ to $c_n$. Relation $Alarm(c_1, c_n)$ describes an
erroneous dataflow from point $c_1$ to $c_n$ where $c_1$ and $c_n$ are
the potential origin and crash point of the error, respectively.

For a conventional source-sink property (i.e., taint analysis),
program points $c_1$ and $c_n$ correspond to the source and sink
points for the target class of errors. For other properties such
as buffer-overrun that do not fit the source-sink problem
formulation, the origin $c_1$ is set to the entry point $c_0$ of the
program and $c_n$ is set to the alarm point.

4.3 Syntactic matching function

To relate program points of the old version $P_1$ and the new
version $P_2$ of the program, we compute function $\delta \in \mathbb{C}P_1 \rightarrow$
$(\mathbb{C}P_1 \cup \mathbb{C}P_2)$:

$$\delta(c_1) = \begin{cases} 
  c_2 & \text{if } c_1 \text{ corresponds to a unique point } c_2 \in \mathbb{C}P_2, \\
  c_1 & \text{otherwise}
\end{cases}$$

where $\mathbb{C}P_1$ and $\mathbb{C}P_2$ denote the sets of program points in $P_1$
and $P_2$, respectively. The function $\delta$ translates program point $c_1$ in the old version to the corresponding program point $c_2$ in the new version. If no corresponding program point exists,
or multiple possibilities exist, then $c_1$ is not translated. In
our implementation, we check the correspondence between
two program points $c_1$ and $c_2$ through the following steps:

1. Check whether $c_1$ and $c_2$ are from the matched file. Our
implementation matches the old file with the new file
if their names match. This assumption can be relaxed if
renaming history is available in a version control system.

2. Check whether $c_1$ and $c_2$ are from the matched lines. Our
implementation matches the old line with the new line
using the GNU diff utility.

3. Check whether $c_1$ and $c_2$ have the same program com-
mands. In practice, one source code line can be translated
into multiple commands in the intermediate representa-
ton of program analyzer.

5 Experimental Evaluation

Our evaluation aims to answer the following questions:

Q1. How effective is Drake for continuous and interactive
reasoning?

Q2. How do different parameter settings of Drake affect
the quality of ranking?

Q3. Does Drake scale to large programs?
Table 1. Benchmark characteristics. Old and New denote program versions before and after introducing the bugs. Size reports the lines of code before preprocessing. Δ reports the percentage of changed lines of code across versions.

<table>
<thead>
<tr>
<th>Program</th>
<th>Version</th>
<th>Size (KLOC)</th>
<th>Δ</th>
<th>#Bugs</th>
<th>Bug Type</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Old</td>
<td>New</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>shntool</td>
<td>3.0.4</td>
<td>3.0.5</td>
<td>13</td>
<td>1</td>
<td>6</td>
</tr>
<tr>
<td></td>
<td>13</td>
<td>13</td>
<td>100</td>
<td></td>
<td>Integer overflow</td>
</tr>
<tr>
<td>latex2rtf</td>
<td>2.1.0</td>
<td>2.1.1</td>
<td>27</td>
<td>2</td>
<td>2</td>
</tr>
<tr>
<td></td>
<td>27</td>
<td>27</td>
<td>100</td>
<td></td>
<td>Format string</td>
</tr>
<tr>
<td>urjtag</td>
<td>0.7</td>
<td>0.8</td>
<td>45</td>
<td>46</td>
<td>16</td>
</tr>
<tr>
<td></td>
<td>46</td>
<td>46.18</td>
<td>100</td>
<td></td>
<td>Format string</td>
</tr>
<tr>
<td>optipng</td>
<td>0.5.2</td>
<td>0.5.3</td>
<td>60</td>
<td>61</td>
<td>2</td>
</tr>
<tr>
<td></td>
<td>61</td>
<td>61</td>
<td>100</td>
<td></td>
<td>Integer overflow</td>
</tr>
<tr>
<td>wget</td>
<td>1.11.4</td>
<td>1.12</td>
<td>42</td>
<td>65</td>
<td>6</td>
</tr>
<tr>
<td></td>
<td>65</td>
<td>65.47</td>
<td>100</td>
<td></td>
<td>Buffer overrun</td>
</tr>
<tr>
<td>readdir</td>
<td>2.23.2</td>
<td>2.24</td>
<td>63</td>
<td>65</td>
<td>1</td>
</tr>
<tr>
<td></td>
<td>65</td>
<td>65</td>
<td>100</td>
<td></td>
<td>Buffer overrun</td>
</tr>
<tr>
<td>grep</td>
<td>2.18</td>
<td>2.19</td>
<td>68</td>
<td>68</td>
<td>1</td>
</tr>
<tr>
<td></td>
<td>68</td>
<td>68</td>
<td>100</td>
<td></td>
<td>Buffer overrun</td>
</tr>
<tr>
<td>sed</td>
<td>4.2.2</td>
<td>4.3</td>
<td>48</td>
<td>83</td>
<td>40</td>
</tr>
<tr>
<td></td>
<td>83</td>
<td>83.40</td>
<td>100</td>
<td></td>
<td>Buffer overrun</td>
</tr>
<tr>
<td>sort</td>
<td>7.1</td>
<td>7.2</td>
<td>96</td>
<td>98</td>
<td>3</td>
</tr>
<tr>
<td></td>
<td>98</td>
<td>98</td>
<td>100</td>
<td></td>
<td>Buffer overrun</td>
</tr>
<tr>
<td>tar</td>
<td>1.27</td>
<td>1.28</td>
<td>108</td>
<td>112</td>
<td>4</td>
</tr>
<tr>
<td></td>
<td>112</td>
<td>112.04</td>
<td>100</td>
<td></td>
<td>Buffer overrun</td>
</tr>
</tbody>
</table>

5.1 Experimental Setup

All experiments were conducted on Linux machines with i7 processors running at 3.4 GHz and with 16 GB memory. We implemented the Bayesian inference engine based on libDAI [35].

Instance analyses. We have implemented our system with Sparrow, a static analysis framework for C programs [38]. Sparrow is designed to be soundy [32] and its analysis is flow-, field-sensitive and partially context-sensitive. It basically computes both numeric and pointer values using the interval domain and allocation-site-based heap abstraction. Sparrow has two analysis engines: an interval analysis for buffer-overrun errors, and a taint analysis for format-string and integer-overflow errors. The taint analysis checks whether unchecked user inputs and overflowed integers are used as arguments of printf-like functions and malloc-like functions, respectively. Since each engine is based on different abstract semantics, we run Drake separately on the analysis results of each engine.

We instrumented Sparrow to generate the elementary dataflow relations (DUEdge, Src, and Dst) in Section 4 and used an off-the-shelf Datalog solver Soufflé [18] to compute derivation trees. The dataflow relations are straightforwardly extracted from the sparse analysis framework [37] on which Sparrow is based. Our instrumentation comprises 0.5K lines while the original Sparrow tool comprises 15K lines of OCaml code.

Benchmarks. We evaluated Drake on the suite of 10 benchmarks shown in Table 1. The benchmarks include those from previous work applying Sparrow [15] as well as GNU open source packages with recently submitted bug-fix commits. Since commit-level source code changes typically introduce modest semantic differences, we ran our differential reasoning process on two consecutive minor versions of the programs before and after the bugs were introduced.

Baselines. We compare Drake to two baseline techniques: Bingo [43] and SynMask, modeled after Facebook Infer’s strategy [39]. Bingo is an interactive alarm ranking system for batch-mode analysis. It computes the ranking of alarms using the Bayesian network extracted from only the new version of the program. SynMask, on the other hand, provides differential reasoning ability, but with a syntactic alarm masking process: it suppresses alarms that have already been reported in the old version of the program using the syntactic matching algorithm described in Section 4.3.

5.2 Effectiveness

This section evaluates the effectiveness of Drake’s ranking compared to the baseline systems. We instantiate Drake with two different settings, DrakeSound and DrakeUnsound as described in Section 3.3. DrakeSound is bootstrapped by assuming the old variants of common alarms to be false (strong mode in Section 3.3) and its input parameter ϵ is set to 0.001. DrakeUnsound aggressively deprioritizes the alarms by assuming both of the old and new variants of common alarms to be false (aggressive mode in Section 3.3), and setting ϵ to 0. For each setting, we measure three metrics: (a) the quality of the initial ranking based on the differential derivation graph, (b) the quality of ranking after transferring old feedback, and (c) the quality of the interactive ranking process. For Bingo, we show the number of user interactions on the alarms only from the new version. For SynMask, we report the number of alarms and missed bugs after syntactic masking.

Table 2 shows the performance of each system. The Initial and Feedback columns report the positions of last true alarm in the initial ranking before and after feedback transfer (corresponding to metrics (a) and (b) above, respectively). The #Iters columns report the number of iterations within which all bugs in each benchmark were discovered (corresponding to metric (c) above).
Table 2. Effectiveness of Drake. **Batch** reports the number of alarms from old and new version of programs. **Bingo** reports the number of interactions until all of the bugs have been inspected. For **SynMask**, we show the number of missed bugs and alarms after the syntactic alarm masking. **DrakeUnsound** and **DrakeSound** show the initial ranking and the number of iterations of Drake with sound and unsound settings.

<table>
<thead>
<tr>
<th>Program</th>
<th>Batch #Old</th>
<th>Batch #New</th>
<th>Batch #Iters</th>
<th>SynMask #Missed</th>
<th>SynMask #Diff</th>
<th>DrakeUnsound Initial</th>
<th>DrakeUnsound Feedback</th>
<th>DrakeUnsound #Iters</th>
<th>DrakeSound Initial</th>
<th>DrakeSound Feedback</th>
<th>DrakeSound #Iters</th>
</tr>
</thead>
<tbody>
<tr>
<td>shntool</td>
<td>20</td>
<td>23</td>
<td>13</td>
<td>3</td>
<td>3</td>
<td>N/A</td>
<td>N/A</td>
<td>N/A</td>
<td>8</td>
<td>21</td>
<td>19</td>
</tr>
<tr>
<td>latex2rtf</td>
<td>7</td>
<td>13</td>
<td>6</td>
<td>0</td>
<td>6</td>
<td>5</td>
<td>6</td>
<td>5</td>
<td>12</td>
<td>9</td>
<td>6</td>
</tr>
<tr>
<td>urjtag</td>
<td>15</td>
<td>35</td>
<td>22</td>
<td>0</td>
<td>27</td>
<td>25</td>
<td>16</td>
<td>18</td>
<td>28</td>
<td>25</td>
<td>21</td>
</tr>
<tr>
<td>optipng</td>
<td>50</td>
<td>67</td>
<td>14</td>
<td>0</td>
<td>17</td>
<td>11</td>
<td>5</td>
<td>4</td>
<td>26</td>
<td>5</td>
<td>9</td>
</tr>
<tr>
<td>wget</td>
<td>850</td>
<td>792</td>
<td>167</td>
<td>0</td>
<td>217</td>
<td>122</td>
<td>139</td>
<td>54</td>
<td>392</td>
<td>317</td>
<td>122</td>
</tr>
<tr>
<td>readelf</td>
<td>841</td>
<td>882</td>
<td>80</td>
<td>0</td>
<td>108</td>
<td>28</td>
<td>4</td>
<td>4</td>
<td>216</td>
<td>182</td>
<td>25</td>
</tr>
<tr>
<td>grep</td>
<td>916</td>
<td>913</td>
<td>53</td>
<td>1</td>
<td>204</td>
<td>N/A</td>
<td>N/A</td>
<td>N/A</td>
<td>15</td>
<td>10</td>
<td>9</td>
</tr>
<tr>
<td>sed</td>
<td>572</td>
<td>818</td>
<td>102</td>
<td>0</td>
<td>398</td>
<td>262</td>
<td>209</td>
<td>60</td>
<td>154</td>
<td>118</td>
<td>41</td>
</tr>
<tr>
<td>sort</td>
<td>684</td>
<td>715</td>
<td>177</td>
<td>0</td>
<td>41</td>
<td>14</td>
<td>14</td>
<td>10</td>
<td>33</td>
<td>9</td>
<td>13</td>
</tr>
<tr>
<td>tar</td>
<td>1,229</td>
<td>1,369</td>
<td>219</td>
<td>0</td>
<td>156</td>
<td>23</td>
<td>29</td>
<td>15</td>
<td>56</td>
<td>82</td>
<td>32</td>
</tr>
<tr>
<td><strong>Total</strong></td>
<td>5,184</td>
<td>5,627</td>
<td>853</td>
<td>4</td>
<td>1,177</td>
<td>490</td>
<td>422</td>
<td>173</td>
<td>940</td>
<td>778</td>
<td>297</td>
</tr>
</tbody>
</table>

Figure 7. The normalized number of iterations until the last true alarm has been discovered with different values of parameter $\epsilon$ for DrakeSound.

In general, the number of alarms of the batch-mode analyses (the “Batch” columns) are proportional to the size of program. Likewise, the number of syntactically new alarms by SynMask is proportional to the amount of syntactic difference. Counterintuitive examples are wget, grep, and readelf. In case of wget, the number of alarms decreased even though the code size increased. It is mainly because a part of user-defined functionalities which reported many alarms has been replaced with library calls. Furthermore, a large part of the newly added code consists of simple wrappers of library calls that do not have buffer accesses. On the other hand, small changes of grep and readelf introduced many new alarms because the changes are mostly in core functionalities that heavily use buffer accesses. When such a complex code change happens, SynMask cannot suppress false alarms effectively and can even miss real bugs. In case of grep, SynMask still reports 22.3% of alarms compared to the batch mode and misses the newly introduced bug.

On the other hand, Drake consistently shows effectiveness in the various cases. For example, DrakeUnsound initially shows the bug in readelf at rank 28, and this ranking rises to 4 after transferring the old feedback. Finally the bug is presented at the top only within 4 iterations out of 108 syntactically new alarms. Furthermore, DrakeSound requires only 9 iterations to detect the bug in grep that is missed by the syntactic approach, which was initially ranked at 15. An exceptional case is shntool and tar. Initially, the ranking of the last true alarm of tar was 56 but it drops to 82 after the feedback transfer. This phenomenon occurs due to false generalization of user feedback resulting from various sources of imprecision in practice such as abstract semantics, approximate derivation graphs, or approximate marginal inference. However, the iterative process gradually improves the quality of the ranking, and eventually finds the bug within 32 rounds out of 1,369 alarms in the new version.

In total, Drake dramatically reduces manual effort for inspecting alarms. The original analysis in the batch mode reports 5,184 and 5,627 alarms for old and new versions of programs, respectively. Applying Bingo on the alarms from new versions requires the user to inspect 853 (15.2%) alarms. SynMask suppresses all the previous alarms and reports 1,177 (20.9%) alarms. However, SynMask misses 4 bugs that were previously false alarms in the old version. DrakeUnsound misses the same 4 bugs because it also suppresses the old alarms. Instead, DrakeSound presents the remaining bugs only within 173 (3.1%) iterations. DrakeSound finds all the...
Table 3. Sizes of the old, new and merged Bayesian networks in terms of the number of tuples (#T) and clauses (#C), and the average iteration time on the merged network.

<table>
<thead>
<tr>
<th>Program</th>
<th>Old #T</th>
<th>Old #C</th>
<th>New #T</th>
<th>New #C</th>
<th>Merged #T</th>
<th>Merged #C</th>
<th>Time(s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>shntool</td>
<td>208</td>
<td>296</td>
<td>236</td>
<td>341</td>
<td>895</td>
<td>1,800</td>
<td>21</td>
</tr>
<tr>
<td>latex2rtf</td>
<td>152</td>
<td>179</td>
<td>710</td>
<td>943</td>
<td>1,674</td>
<td>2,550</td>
<td>17</td>
</tr>
<tr>
<td>urjtag</td>
<td>547</td>
<td>765</td>
<td>676</td>
<td>920</td>
<td>1,463</td>
<td>2,251</td>
<td>23</td>
</tr>
<tr>
<td>optipng</td>
<td>492</td>
<td>561</td>
<td>633</td>
<td>730</td>
<td>1,901</td>
<td>3,313</td>
<td>7</td>
</tr>
<tr>
<td>wget</td>
<td>3,959</td>
<td>4,484</td>
<td>3,285</td>
<td>3,594</td>
<td>8,528</td>
<td>13,561</td>
<td>23</td>
</tr>
<tr>
<td>grep</td>
<td>4,265</td>
<td>4,802</td>
<td>4,346</td>
<td>4,091</td>
<td>10,470</td>
<td>16,599</td>
<td>31</td>
</tr>
<tr>
<td>readelf</td>
<td>3,702</td>
<td>4,283</td>
<td>3,952</td>
<td>4,565</td>
<td>10,628</td>
<td>16,658</td>
<td>31</td>
</tr>
<tr>
<td>sed</td>
<td>1,887</td>
<td>2,030</td>
<td>2,971</td>
<td>3,265</td>
<td>6,698</td>
<td>9,730</td>
<td>15</td>
</tr>
<tr>
<td>sort</td>
<td>2,672</td>
<td>2,951</td>
<td>2,796</td>
<td>3,085</td>
<td>8,308</td>
<td>13,597</td>
<td>31</td>
</tr>
<tr>
<td>tar</td>
<td>5,620</td>
<td>6,197</td>
<td>6,096</td>
<td>6,708</td>
<td>17,484</td>
<td>29,129</td>
<td>47</td>
</tr>
<tr>
<td>Total</td>
<td>23,504</td>
<td>26,548</td>
<td>25,701</td>
<td>28,242</td>
<td>68,049</td>
<td>109,188</td>
<td>246</td>
</tr>
</tbody>
</table>

bugs within 297 (5.3%) iterations, a significant improvement over the baseline approaches.

5.3 Sensitivity analysis on different configurations

This section conducts a sensitivity study with different values of parameter $\epsilon$ for $\text{Drake}_{\text{Sound}}$. Recall that $\epsilon$ represents the degree of belief that the same abstract derivation tree from two versions has different concrete behaviors. Therefore, the higher $\epsilon$ is set, the more conservatively $\text{Drake}$ behaves.

Figure 7 shows the normalized number of iterations until all the bugs have been found by $\text{Drake}_{\text{Sound}}$ with different values for $\epsilon$. We observe that the overall number of iterations generally increases as $\epsilon$ increases because $\text{Drake}_{\text{Sound}}$ conservatively suppresses the old information. However, the rankings move opposite to this trend in some cases such as $\text{latex2rtf}$, $\text{readelf}$, and $\text{tar}$. In practice, various kinds of factors are involved in the probability of each alarm such as structure of the network. For example, when bugs are closely related to many false alarms that were transformed from the old versions, an aggressive approach (i.e., small $\epsilon$) can introduce negative effects. In fact, the bugs in the three benchmarks are closely related to huge functions or recursive calls that hinder precise static analysis. In such cases, aggressive assumptions on the previous derivations can be harmful for the ranking.

5.4 Scalability

The scalability of the iterative ranking process mostly depends on the size of the Bayesian network. $\text{Drake}$ optimizes the Bayesian networks using optimization techniques described in previous work [43]. We measure the network size in terms of the number of tuples and clauses in derivation trees after the optimizations, and report the average time for each marginal inference computation where $\epsilon$ is set to 0.001.

Table 3 show the size and average computation time for each iteration. The merged networks have 3x more tuples and 4x more clauses compared to the old and new versions of networks. The average iteration time for all benchmarks is less than 1 minute which is reasonable for user interaction.

6 Related Work

Our work is inspired by recent industrial scale deployments of program analysis tools such as Coverity [4], Facebook Infer [39], Google Tricorder [45], and SonarQube [8]. These tools primarily employ syntactic masking to suppress reporting alarms that are likely irrelevant to a particular code commit. Indeed, syntactic program differencing goes back to the classic Unix diff algorithm proposed by Hunt and McIlroy in 1976 [16]. Our work builds upon these works and uses syntactic matching to identify abstract states before and after a code commit.

Program differencing techniques have been developed by the software engineering community [17, 22]. Their goal is to effectively summarize, to a human developer, the semantic code changes using dependency analysis or logical rules. The reports are typically based on syntactic features of the code change. On the other hand, our goal is to identify newly introduced bugs, and $\text{Drake}$ captures deep semantic changes indicated by the program analysis in the derivation graph.

The idea of checking program properties using information obtained from its previous versions has also been studied by the program verification community, as the problem of differential static analysis [28]. Differential assertion checking [27], verification modulo versions [33], and the SymDiff project [14] are prominent examples of research in this area.

The SafeMerge system [47] considers the problem of detecting bugs introduced while merging code changes. These systems typically analyze the old version of the program to obtain the environment conditions that preclude buggy behavior, and subsequently verify that the new version is bug-free under the same environment assumptions. Therefore, these approaches usually need general-purpose program verifiers, significant manual annotations, and do not consider the problems of user interaction or alarm ranking.

Research on hyperproperties [9] and on relational verification [3] relates the behaviors of a single program on multiple inputs or of multiple programs on the same input. Typical problems studied include equivalence checking [21, 26, 40, 44], information flow security [36], and verifying the correctness of code transformations [20]. Various logical formulations, such as Hoare-style partial equivalence [12], and techniques such as differential symbolic execution [41, 44] have been explored. In contrast to our work, such systems focus on identifying divergent behaviors between programs. On the other hand, in our case, it is almost certain that the programs are semantically equivalent, and our focus is instead on differential bug-finding.
Finally, there is a large body of work in the literature on leveraging probabilistic methods and machine learning to improve static analysis accuracy [19, 24, 29, 30] and find bugs in programs [25, 31]. To the best of our knowledge, our work is the first to apply such techniques to the problem of continuous reasoning.

7 Conclusion

We have presented a system, Drake, for the analysis of continuously evolving programs. Drake prioritizes alarms according to their likely relevance relative to the last code change, and reranks alarms in response to user feedback. Drake operates by comparing the results of the static analysis runs from each version of the program, and builds a probabilistic model of alarm relevance using a differential derivation graph. Our experiments on a suite of ten widely-used C programs demonstrate that Drake dramatically reduces the alarm inspection burden compared to other state-of-the-art techniques without missing any bugs.

References


