



CIS 5530: Networked Systems

Components of a Network

January 25, 2023



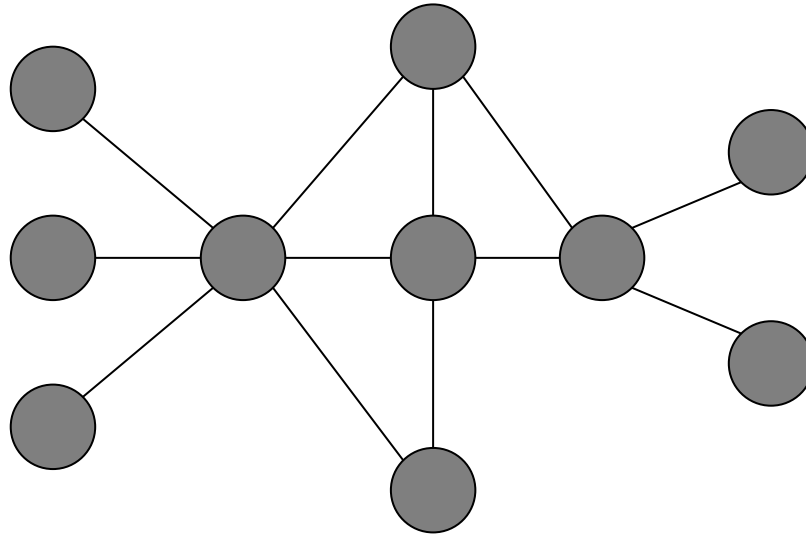
Agenda

- Layering and the end-to-end principle ✓
- Performance metrics ✓
- Components of a Network ← NEXT
 - Physical-layer links
 - Framing
 - L1/L2 nodes
 - VLANs



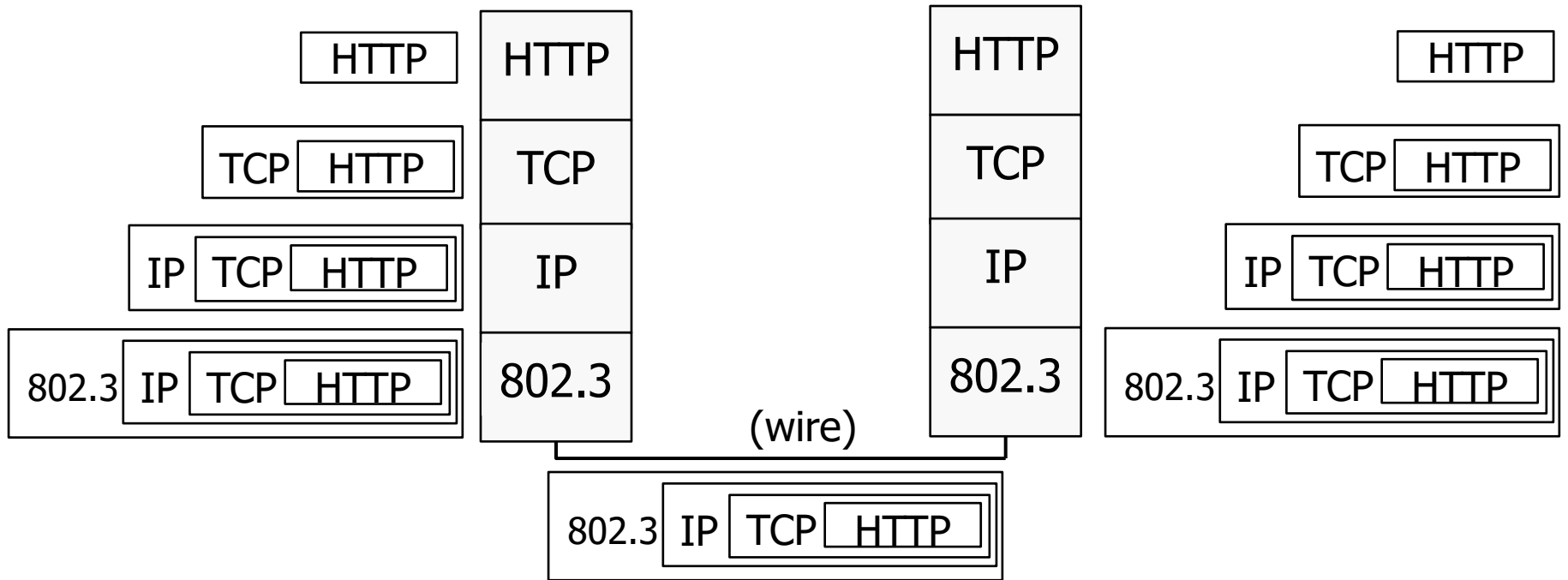
Networked Systems

- Study of systems of **links** that interconnect **computers** in order to move **data**



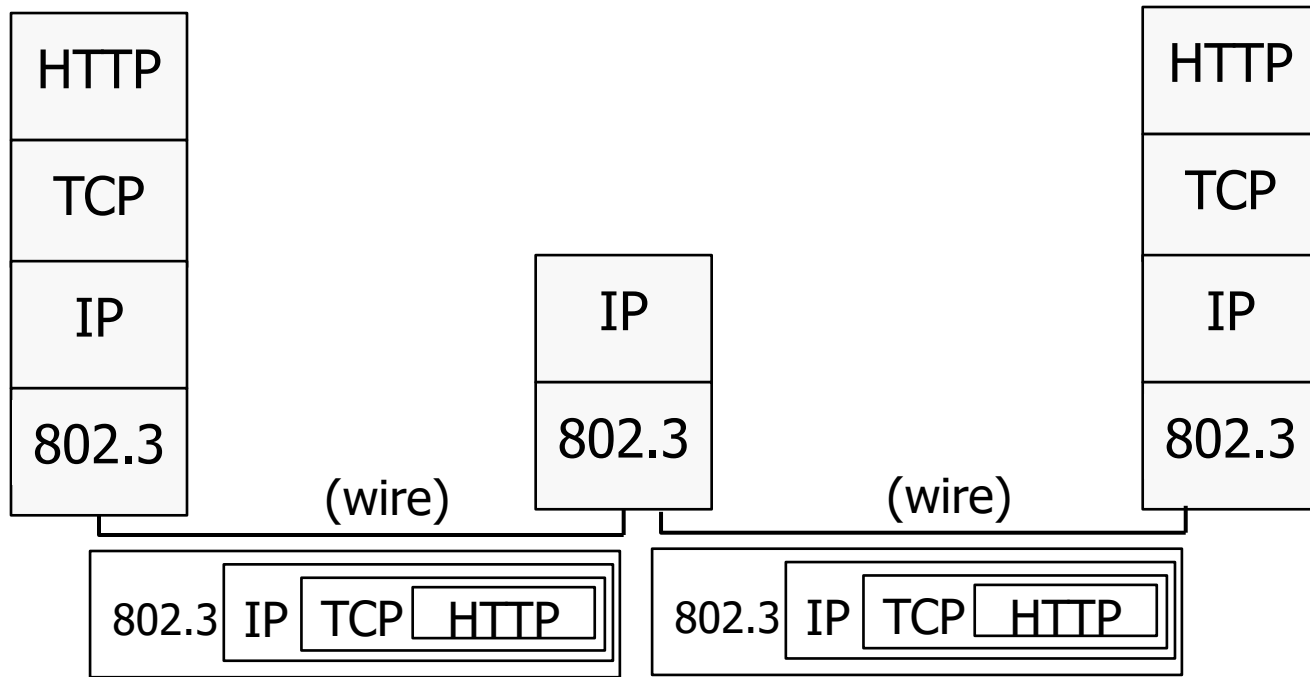


Encapsulation





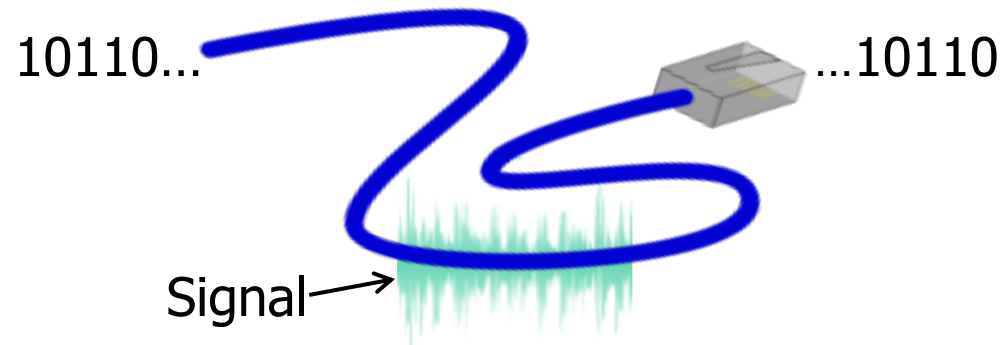
Encapsulation (multihop)





Layer 1: Physical layer

- Connects: Transmitters and receivers
- Lower interface: Analog signals
- Upper interface: Stream of digital bits
- Name of network: Physical links
- Name of message: Bits





Layer 2: Data link layer

- Connects: physical interfaces
 - Usually all the same physical layer protocol
- Lower interface: Stream of bits
- Upper interface: Messages with payloads

- Name of network: Local Area Network (LAN)
- Name of message: Frame





Layer 3: Network layer

- Connects: LANs
 - Possibly different technologies (e.g., WiFi and Ethernet)
- Lower interface: A packet of data
- Upper interface: A segment of data
 - Same thing, just has an extra header
- Name of network: Internet
- Name of message: Packet

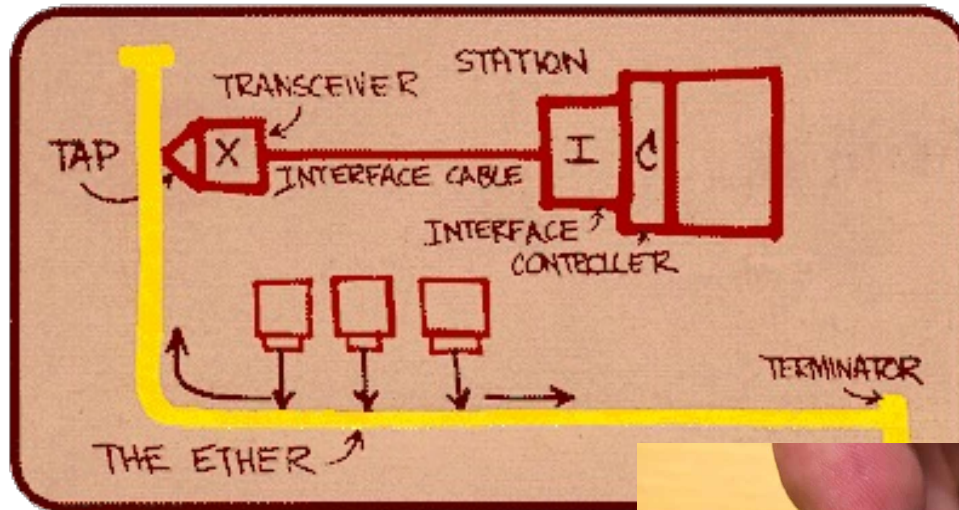


Agenda

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Links: Ethernet

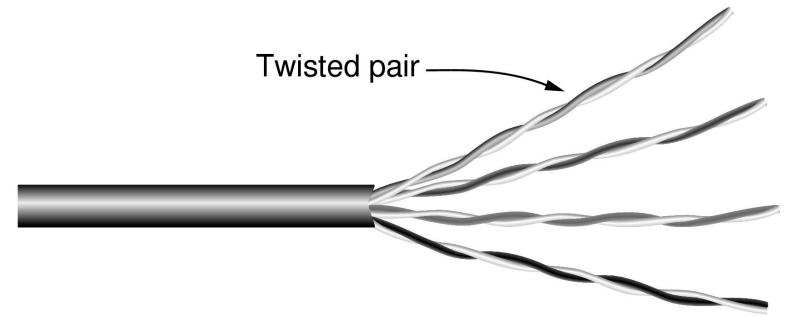


Metcalfe's Ethernet sketch

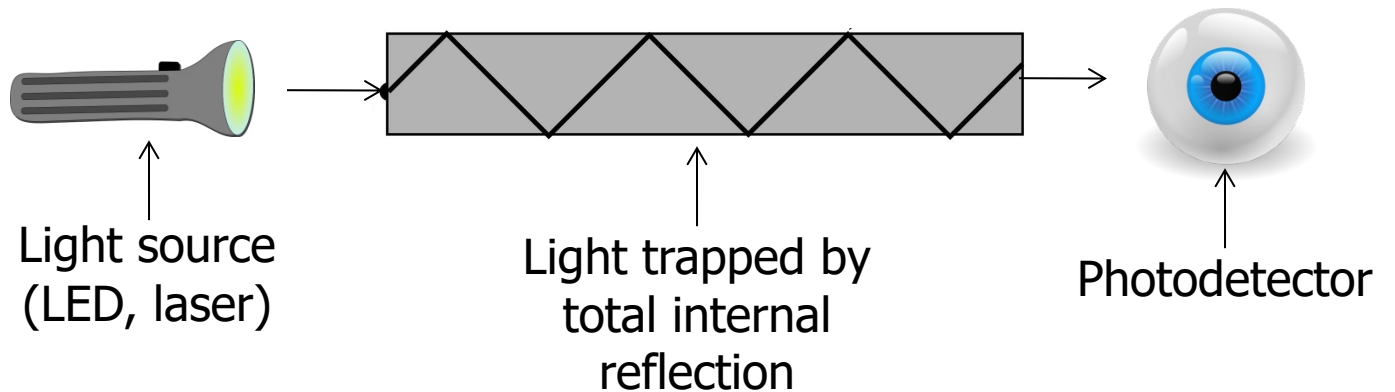
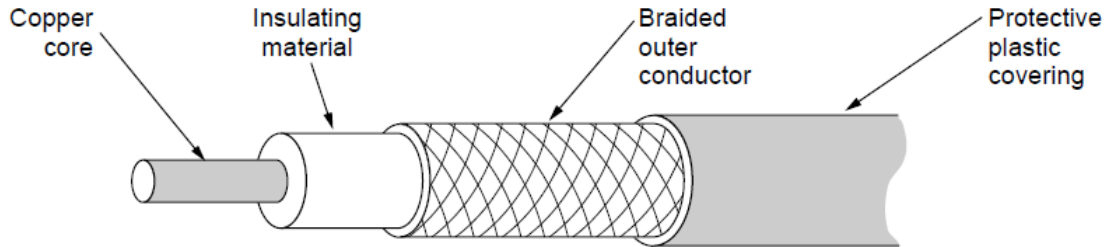




Links: Wires



Twisted pair





Links: Wireless

- Sender radiates signal over a region
 - In many directions, unlike a wire, to potentially many receivers
 - Nearby signals (same freq.) **interfere** at a receiver; need to coordinate use



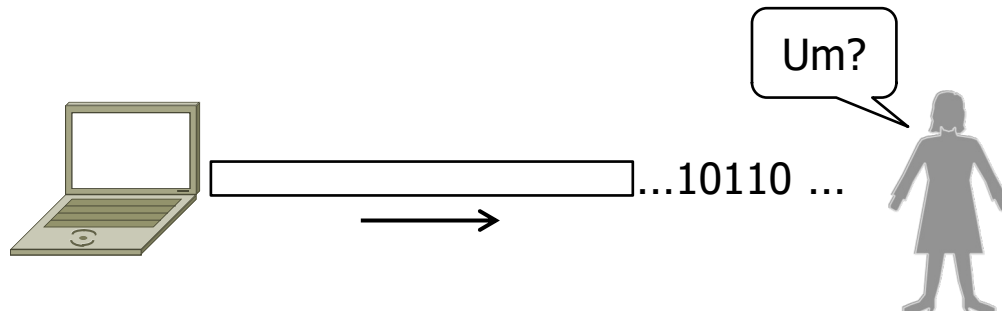


How to convert bits to frames?

- Ex: A message in Ethernet:



- Physical layer puts bits on a link
- **Framing problem**: how does the link layer determine where each frame begins and ends?





Framing methods

- We'll look at two:
 - Byte count (motivation)
 - Byte stuffing
- In practice, the physical layer often provides hints
 - E.g., if the signal on the wire goes flat, that's probably the end of a frame



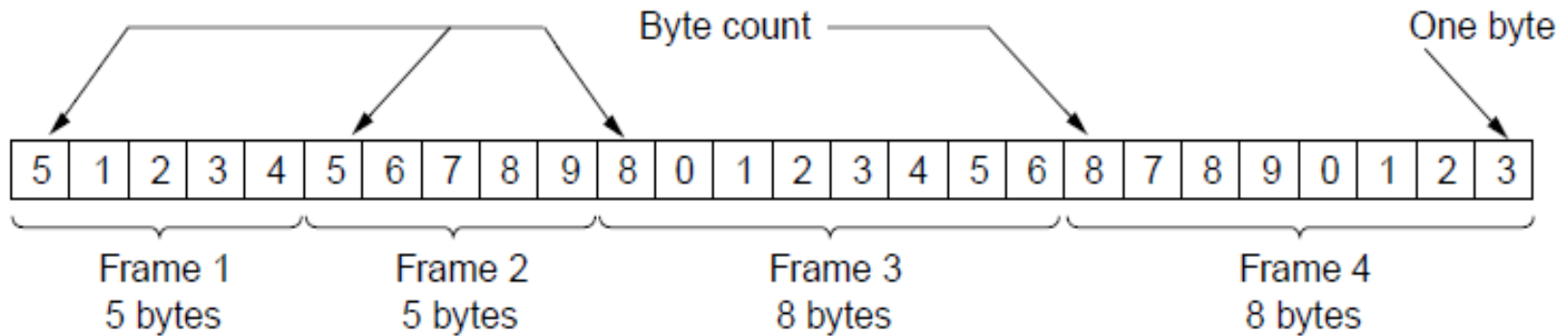
Simple approach: byte count

- First try:
 - Let's start each frame with a length field!
 - It's simple, and hopefully good enough ...



Simple approach: byte count

- First try:
 - Let's start each frame with a length field!

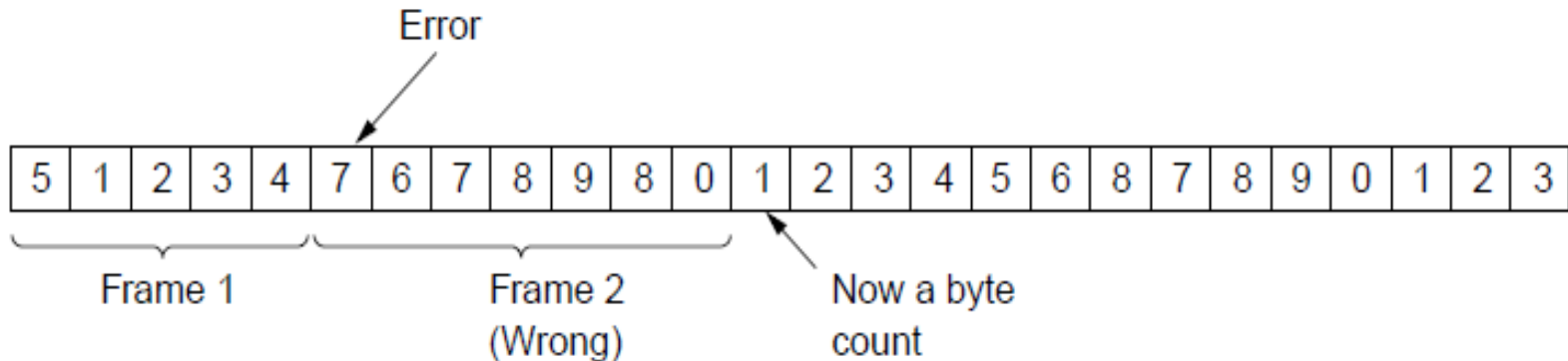


Are we done?



What if count is corrupted?

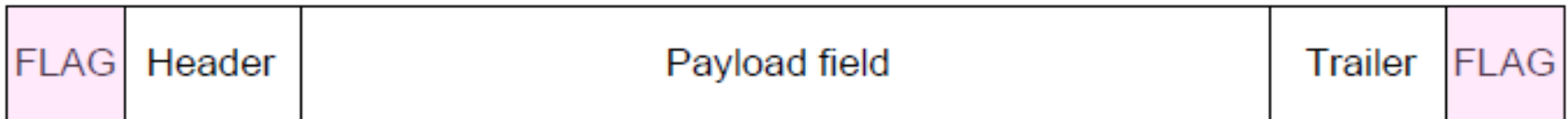
- We will frame the wrong bytes: **framing error**
 - Once framing on a link is desynchronized, it can stay that way
 - Need a method to **resynchronize**





Byte stuffing

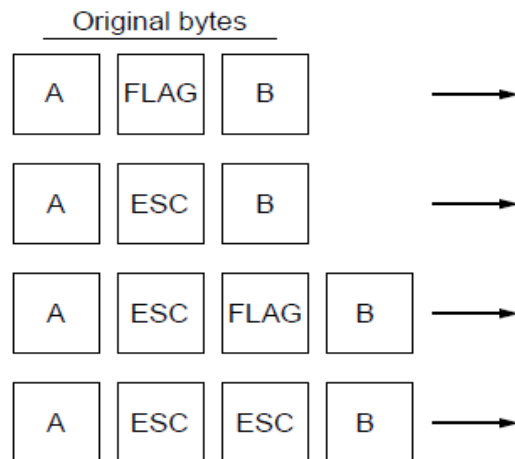
- Better idea:
 - Special FLAG byte that marks start/end of frame
 - When FLAG appears in the payload, “stuff” an escape character into the payload before the FLAG
 - **Complication**: must escape the escape code too!





Byte stuffing examples

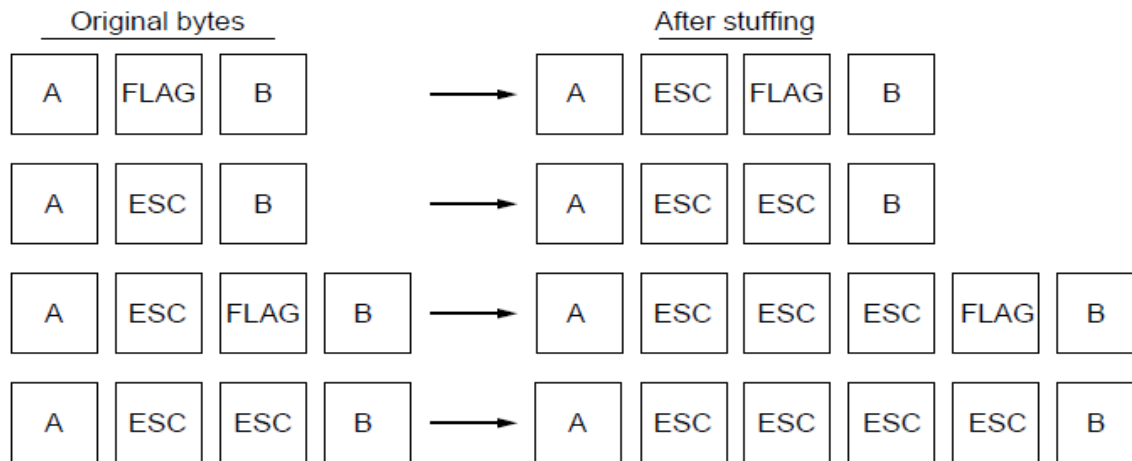
- Rules:
 - Replace each FLAG in data with ESC FLAG
 - Replace each ESC in data with ESC ESC





Byte stuffing examples

- Rules:
 - Replace each FLAG in data with ESC FLAG
 - Replace each ESC in data with ESC ESC



- Now any unescaped FLAG Is the start/end of a frame



Bit Stuffing

- Can stuff at the bit level too
 - Call a flag six consecutive 1s
 - On transmit, after five 1s in the data, insert a 0
 - On receive, a 0 after five 1s is deleted



Bit Stuffing Example

Data bits 0 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 1 0

Transmitted bits
with stuffing



Bit Stuffing Example

Data bits 0 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 1 0

Transmitted bits with stuffing 0 1 1 0 1 1 1 1 1 0 1 1 1 1 1 0 1 1 1 1 1 0 1 0 0 1 0

Stuffed bits

How does this compare with byte stuffing?



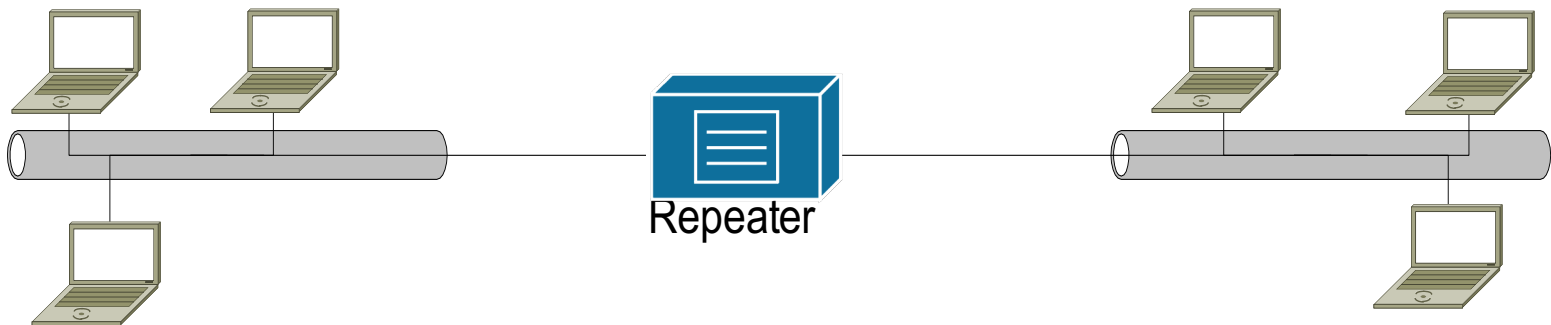
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Physical Layer: Repeaters

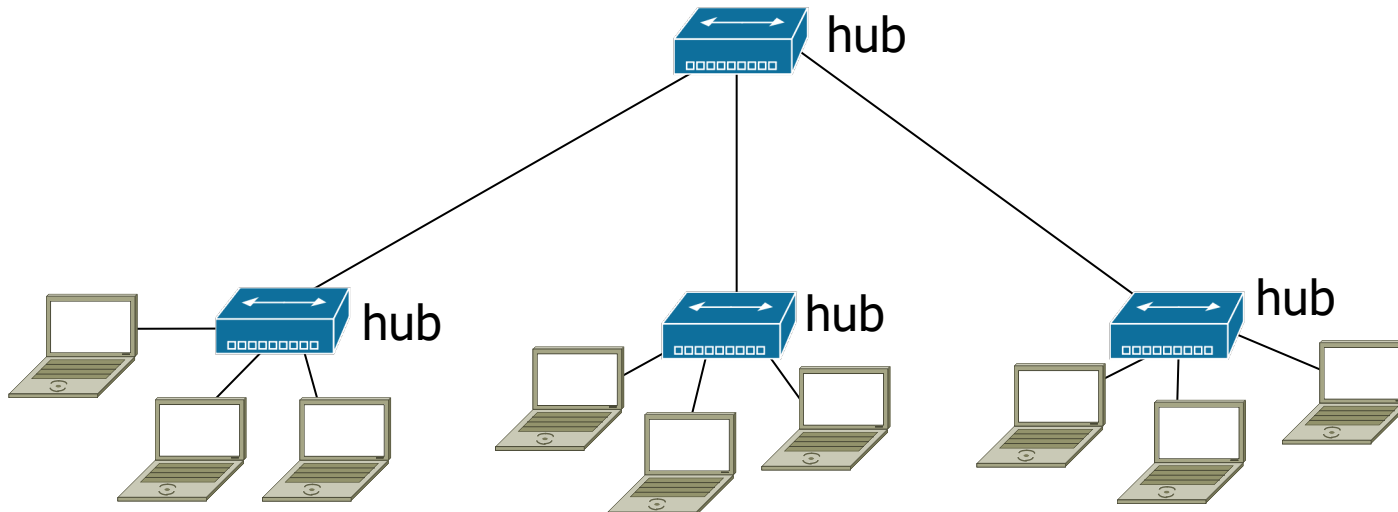
- Distance limitations
 - Signals become weaker as they travel
 - Imposes a limit on the length of each **segment** of a LAN
- Repeaters join two segments together
 - Often analog electronic device
 - Continuously monitors electrical signals
 - Transmits an amplified copy





Physical Layer: Hubs

- Fancy name for multi-port repeater
- Joins multiple segments
 - Does not necessarily amplify the signal
 - Broadcasts traffic that arrives on every other port
- Very similar to repeaters
 - Also operates at the physical layer





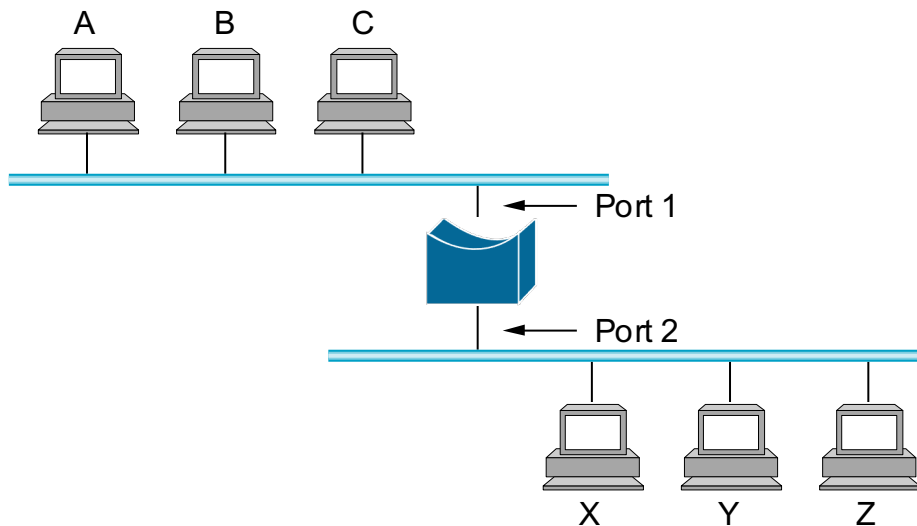
Limitations of Repeaters and Hubs

- One large shared LAN
 - Each bit is sent everywhere
 - So, aggregate throughput is limited
- Cannot support multiple technologies
 - Does not buffer or interpret frames
 - Can't interconnect between different rates/formats
- Limitations on maximum nodes and distances
 - Shared medium imposes length limits
 - E.g., cannot go beyond 2500 meters on Ethernet



Solution: Link Layer Bridges

- Connects two or more segments at the link layer
- Based around a **forwarding table**
 - Extracts destination address from the frame
 - Looks up the destination in a table
 - Forwards the frame to the appropriate segment



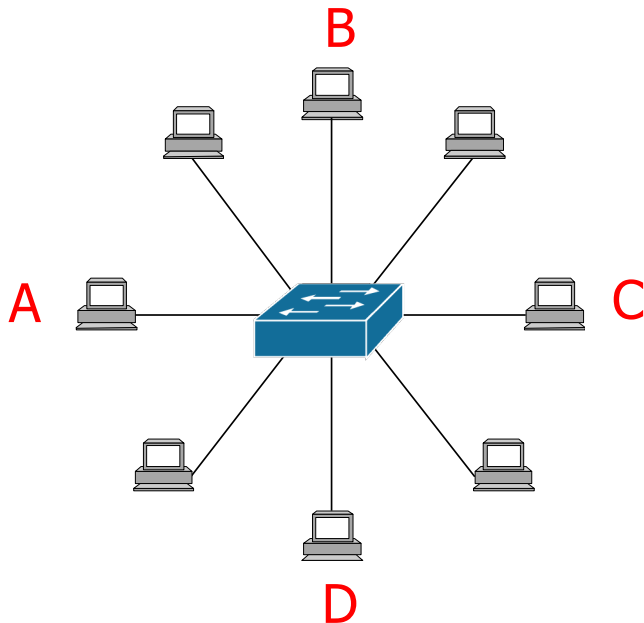
Host	Port
A	1
B	1
C	1
X	2
Y	2
Z	2

Forwarding table



Link Layer Switches

- Typically connects individual computers
 - A switch is essentially the same as a bridge
 - ... though typically used to connect hosts, not segments
- Like bridges, support concurrent communication
 - Host A can talk to C, while B talks to D



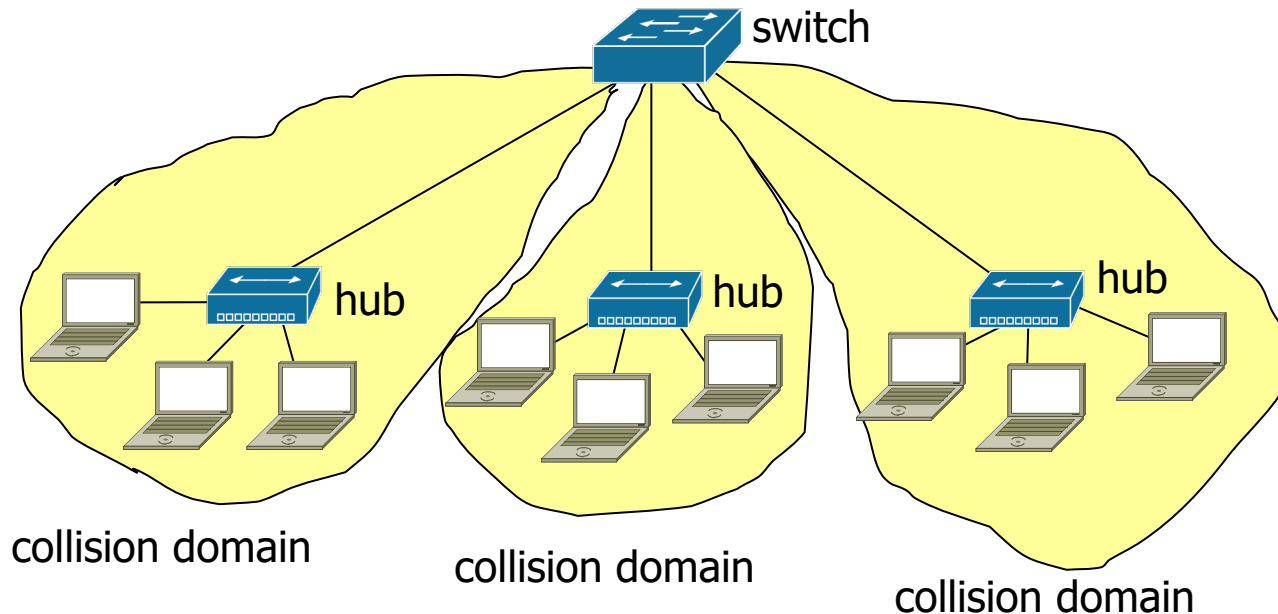
Cisco Gigabit Ethernet switch





Bridges and Switches: Traffic Isolation

- Breaks network into segments
- Filters packets
 - Frame only forwarded to the necessary segments
 - Segments become separate collision domains





Advantages Over Hubs/Repeaters

- Only forwards frames as needed
 - Filters frames to avoid unnecessary load on segments
 - Sends frames only to segments that need to see them
 - Separate collision domains allow longer distances
- Joins segments using different technologies
 - E.g. Ethernet (10MBps, 100MBps) – made possible with buffering
 - Not complete heterogeneity
 - E.g. Ethernet and ATM not compatible