

## Recitation Guide - Week 10

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**Topics Covered:** Graphs, Variance

**Problem 1:**

We've proven that in any graph, the number of vertices of odd degree is even with the Handshake Lemma. Prove this again, using induction.

**Solution:**

We will prove this statement by induction on the number of edges,  $m$ . Let  $P(m)$  be defined as:

In any graph with  $m$  edges, the number of vertices of odd degree is even.

Base Case:  $P(0)$  holds, because a graph with no edge has only isolated vertices, that is, vertices of degree 0. Hence, there are an even number (0) of vertices with odd degree.

Induction Step: Assume  $P(k)$  is true, for an arbitrary  $k \in \mathbb{N}$ . Now, we want to prove  $P(k+1)$  is true.

Let  $G$  be a graph with  $k+1$  edges. Remove an arbitrary edge  $e = \{u, v\}$  from  $G$  (note that it could be any edge), so that we now have a graph  $G'$  with  $k$  edges. By the Induction Hypothesis, the number of vertices with odd degree in  $G'$  is even. Denote the number of vertices with odd degree in  $G'$  to be  $2a$ , where  $a \in \mathbb{N}$ . Now put back the edge  $e$  that we removed earlier. Observe that doing so increases the degree of vertices  $u$  and  $v$  by one each. We consider the following three cases:

*Case 1:* Both  $u$  and  $v$  have odd degree in  $G'$ . Adding  $e$  back would make the degree of both  $u$  and  $v$  even. Hence, the number of vertices with odd degree becomes  $2a - 2$ .

*Case 2:* Both  $u$  and  $v$  have even degree in  $G'$ . Adding  $e$  back would make the degree of both  $u$  and  $v$  odd. Hence, the number of vertices with odd degree becomes  $2a + 2$ .

*Case 3:* Exactly one of  $u$  and  $v$  has odd degree in  $G'$ . WLOG, assume  $u$  has an odd degree and  $v$  has an even degree in  $G'$ . Adding  $e$  back would result in  $u$  with an even degree and  $v$  with an odd degree. Hence, the number of vertices with odd degree would stay unchanged ( $2a$ ).

In all cases, the number of odd degree vertices in  $G$  is even. Thus, we have shown our claim is true when  $m = k + 1$ , concluding our Induction Step and completing our proof.

ALTERNATE:

We will prove this statement by induction on the number of vertices,  $n$ . Let  $P(n)$  be defined as:

In any graph with  $n$  vertices, the number of vertices of odd degree is even.

Base Case:  $P(1)$  holds, because a graph with 1 vertex is an edgeless graph, that is, the graph contains one vertex of degree 0. Hence, there are an even number (0) of vertices with odd degree.

Induction Step: Assume  $P(k)$  is true, for an arbitrary  $k \in \mathbb{Z}^+$ . Now, we want to prove  $P(k+1)$  is true.

Let  $G$  be a graph with  $k+1$  vertices. We partition  $V$  into the following sets:

- $X$  = the neighbors of  $v$  with even degree in  $G$ .
- $Y$  = the neighbors of  $v$  with odd degree in  $G$ .
- $R$  = vertices not adjacent to  $v$  with even degree in  $G$ .
- $S$  = vertices not adjacent to  $v$  with odd degree in  $G$ .

Since  $Y$  and  $S$  partition the set of odd-degree vertices in  $G$ , want to show  $|Y| + |S|$  is even.

Remove an arbitrary vertex  $v$  with degree  $\alpha$  from  $G$ , so that we now have a graph  $G'$  with  $k$  vertices.

*Case 1:*  $\alpha$  is even. Then  $\alpha = 2k, k \in \mathbb{N}$ . This means that  $|X| + |Y| = 2k$ , so  $|X|$  and  $|Y|$  have the same parity.

The degree of each neighbor of  $v$  decreased by 1, so all  $x \in X$  in  $G$  became odd degree vertices in  $G'$ . By the Induction Hypothesis, the number of vertices with odd degree in  $G'$  is even. Denote the number of vertices with odd degree in  $G'$  to be  $2a$ , where  $a \in \mathbb{N}$ . In  $G'$  this gives  $|X| + |S| = 2a$ , since  $|X|$  now represents the new odd vertices in  $G'$ , and  $|S|$  represents the odd-degree vertices unaffected by the removal of  $v$ . Note that  $|X|$  and  $|S|$  must have the same parity (think why this is true!).

Now put back the vertex  $v$  that we removed earlier. Observe that doing so increases the degree of each of the neighbors of  $v$  by 1. All vertices in  $X$  now have even degree, and all vertices in  $Y$  have odd degree again.

Since  $|X|$  and  $|S|$  must have the same parity, we have  $|X|, |Y|, |S|$  all have the same parity. This means  $|Y| + |S|$  is even (since odd+odd is even, and even + even is even).

*Case 2:*  $\alpha$  is odd. (We include  $v$  in  $S$ ). Then  $\alpha = 2k+1, k \in \mathbb{N}$ . This means that  $|X| + |Y| = 2k+1$ , so  $|X|$  and  $|Y|$  have the opposite parity.

The degree of each neighbor of  $v$  decreased by 1, so all  $x \in X$  in  $G$  became odd degree vertices in  $G'$ . By the Induction Hypothesis, the number of vertices with odd degree in  $G'$  is even. Denote the number of vertices with odd degree in  $G'$  to be  $2a$ , where  $a \in \mathbb{N}$ . In  $G'$  this gives  $|X| + |S| = 2a$ , since  $|X|$  now represents the new odd vertices in  $G'$ , and  $S' = S \setminus \{v\}$  since  $v$  had odd degree, removing it means  $|S'| = |S| - 1$  since all other vertices in  $S$  are unaffected. Note that  $|X| + |S'| = 2a$ , or  $|X| + |S| = 2a + 1$ . This means  $|X|$  and  $|S|$  have opposite parity.

Now put back the vertex  $v$  that we removed earlier. Observe that doing so increases the degree of each of the neighbors of  $v$  by 1. All vertices in  $X$  now have even degree, and all vertices in  $Y$  have odd degree again.

Since  $|X|$  and  $|S|$  must have the opposite parity, and  $|X|$  and  $|Y|$  have opposite parity, we have note  $|Y|$  and  $|S|$  have the same parity. (Can you prove this?). This means  $|Y| + |S|$  is even (since odd+odd is even, and even + even is odd).

In all cases, the number of odd degree vertices in  $G$  is even. Thus, we have shown our claim is true when  $m = k + 1$ , concluding our Induction Step and completing our proof.

**Problem 2:**

Prove that a graph  $G = (V, E)$  is connected iff for every partition of  $V$  into two disjoint, non-empty sets  $S$  and  $T$ , there exists an edge between some vertex in  $S$  and some vertex in  $T$ .

**Solution:**

( $\implies$ ): We first show that if a graph  $G = (V, E)$  is connected, then for every partition of  $V$  into two disjoint, non-empty sets  $S$  and  $T$ , there exists an edge between some vertex in  $S$  and some vertex in  $T$ . Consider an arbitrary partition  $V = S \cup T$  into two disjoint, non-empty sets  $S$  and  $T$ . Let  $x \in S$  and  $y \in T$ ; since  $G$  is connected, there must be a path  $x \rightsquigarrow y$ , say:

$$P = x - v_1 - v_2 - \dots - v_{k-1} - y$$

We claim that there must be some edge from  $S$  to  $T$  in this path. Suppose towards contradiction that all edges are between two vertices in  $S$  or two vertices in  $T$ . Since  $x \in S$ , we must have  $v_1 \in S$ . Similarly, we must then have  $v_2 \in S$ . We may continue this process to show that  $v_{k-1} \in S$  (see if you can formally prove this with induction!), and  $y \in S$ , a contradiction.

( $\impliedby$ ): We now show that, given a graph  $G = (V, E)$ , if for every partition of  $V$  into two disjoint, non-empty sets  $S$  and  $T$ , there exists an edge between some vertex in  $S$  and some vertex in  $T$ , then  $G$  is connected. We proceed by proving the contrapositive, namely, that if  $G$  is not connected, then there exists a partition of  $V$  into disjoint nonempty sets  $S$  and  $T$  with no edges between the two.

Since  $G$  is not connected, it must have at least two connected components. Let  $S$  be a connected component of  $G$  and let  $T = V \setminus S$ . By definition of connected component, there is no edge from a vertex of  $S$  to one in  $T$  (if there were, we would violate the maximality condition). This gives us our desired partition.

**Problem 3:**

Let  $X, Y$  be two random variables defined on the same probability space. The covariance of  $X$  and  $Y$  is defined to be

$$\text{Cov}(X, Y) = E[XY] - E[X]E[Y]$$

Note that by Proposition 17.5 if  $X \perp Y$  then  $\text{Cov}(X, Y) = 0$ . Two random variables are uncorrelated if  $\text{Cov}(X, Y) = 0$  (equivalently, such that  $E[XY] = E[X]E[Y]$ ). Therefore independence implies uncorrelation.

- (a) Give an example of two random variables that are uncorrelated yet they are not independent
- (b) Prove that  $\text{Cov}(X, Y) = E[(X - E[X])(Y - E[Y])]$ . Conclude that  $\text{Var}(X) = \text{Cov}(X, X)$ .
- (c) Then show that  $\text{Var}[X + Y] = \text{Var}[X] + \text{Var}[Y] + 2\text{Cov}(X, Y)$ .

**Solution:**

- (a) Let our probability space consist of the outcomes  $\{(0, 1), (0, -1), (1, 0), (-1, 0)\}$  with a uniform probability distribution. Define the random variable  $X$  to denote the first value in the ordered pair and the random variable  $Y$  to denote the second value.

$$\begin{aligned} E[X] &= E[Y] = \sum_{x \in X} x \Pr[X = x] \\ &= (0)(1/2) + (-1)(1/4) + (1)(1/4) \\ &= 0 \\ E[XY] &= \sum_{z \in XY} z \Pr[XY = z] \\ &= (0)(1) = 0 \\ \text{Cov}(X, Y) &= 0 - 0 = 0 \end{aligned}$$

Thus,  $X$  and  $Y$  are uncorrelated. However, they are not independent:

$$\begin{aligned} \Pr[X = 1 \cap Y = 1] &= 0 \\ \Pr[X = 1] &= \frac{1}{4} \\ \Pr[Y = 1] &= \frac{1}{4} \\ \Pr[X = 1] \times \Pr[Y = 1] &= \frac{1}{16} \neq 0 \end{aligned}$$

- (b) Applying the Linearity of Expectation (LOE), we see:

$$\begin{aligned} E[(X - E[X])(Y - E[Y])] &= E[XY - X \cdot E[Y] - E[X] \cdot Y + E[X]E[Y]] \\ &= E[XY] - E[X \cdot E[Y]] - E[E[X] \cdot Y] + E[E[X]E[Y]] \\ &\hspace{15em} \text{(by LOE)} \end{aligned}$$

Observing that  $E[S]$  is a constant for any random variable  $S$ , and noting that  $E[c] = c$  for all constants  $c \in \mathbb{R}$ , we see:

$$\begin{aligned}
 &= E[XY] - E[Y]E[X] - E[X]E[Y] + E[X]E[Y] \\
 &= E[XY] - E[X]E[Y] \\
 &= \text{Cov}(X, Y) \qquad \qquad \qquad (\text{by definition})
 \end{aligned}$$

Now applying this result gives:

$$\begin{aligned}
 \text{Cov}(X, X) &= E[(X - E[X])(X - E[X])] \\
 &= E[(X - E[X])^2] \\
 &= \text{Var}(X)
 \end{aligned}$$

(c) We begin by expanding the definition of variance:

$$\begin{aligned}
 \text{Var}(X + Y) &= E[(X + Y)^2] - E[X + Y]^2 \\
 &= E[X^2 + 2XY + Y^2] - (E[X] + E[Y])^2 \qquad \qquad \qquad (\text{by LOE}) \\
 &= E[X^2] + 2E[XY] + E[Y^2] - (E[X]^2 + 2E[X]E[Y] + E[Y]^2) \qquad (\text{by LOE})
 \end{aligned}$$

Rearranging this gives:

$$\begin{aligned}
 &= E[X^2] - E[X]^2 + E[Y^2] - E[Y]^2 + 2E[XY] - 2E[X]E[Y] \\
 &= \text{Var}(X) + \text{Var}(Y) + 2\text{Cov}(X, Y)
 \end{aligned}$$

**Problem 4:**

Waley is prepping for an Easter Egg hunt. He has a basket of  $n$  eggs, and independently for each egg, paints it either red with probability  $\frac{1}{3}$  or blue otherwise. He defines an equivalence relation  $\rho$  where  $x \rho y$  if and only if  $x$  and  $y$  are the same color. What is  $E[|\rho|]$ ?

**Solution:**

Call the set of eggs  $E$ . Consider the sample space:

$$\Omega = \{R, B\}^n$$

where each outcome lists the colors of the eggs in a fixed order. We define:

Let  $R_x$  be the event that egg  $x$  is colored red,  $x \in E$ .

Let  $B_x$  be the event that egg  $x$  is colored blue,  $x \in E$ .

Let  $S_{x,y}$  be the event that eggs  $x$  and  $y$  have the same color,  $x, y \in E$ .

Let  $|\rho|$  be the random variable denoting the cardinality of  $\rho$ .

Let  $I_{S_{x,y}}$  be an indicator random variable for  $S_{x,y}$ ,  $x, y \in E$ .

We seek  $E[|\rho|]$ . All elements of  $\rho$  are ordered pairs  $(x, y)$  where  $x, y \in E$ , and  $x$  and  $y$  have the same color. We thus see that:

$$|\rho| = \sum_{x,y} I_{S_{x,y}} = \sum_x I_{S_{x,x}} + \sum_{x \neq y} I_{S_{x,y}}$$

By the Linearity of Expectation, we have that:

$$\begin{aligned} E[|\rho|] &= E \left[ \sum_x I_{S_{x,x}} + \sum_{x \neq y} I_{S_{x,y}} \right] \\ &= \sum_x E[I_{S_{x,x}}] + \sum_{x \neq y} E[I_{S_{x,y}}] \\ &= \sum_x \Pr[S_{x,x}] + \sum_{x \neq y} \Pr[S_{x,y}] \end{aligned}$$

Note that  $\Pr[S_{x,x}] = 1$ ,  $\forall x \in E$ , since an egg is always the same color as itself. For the case where  $x \neq y$ , we see that  $S_{x,y} = (R_x \cap R_y) \cup (B_x \cap B_y)$ . Since these events are disjoint, we can add their probabilities with the Sum Rule:

$$= \sum_x 1 + \sum_{x \neq y} \Pr[R_x \cap R_y] + \Pr[B_x \cap B_y]$$

Observing that  $R_x \perp R_y$  and  $B_x \perp B_y$  when  $x \neq y$ , since Waley colors eggs independently:

$$\begin{aligned} &= n + \sum_{x \neq y} (\Pr[R_x] \times \Pr[R_y] + \Pr[B_x] \times \Pr[B_y]) \\ &= n + \sum_{x \neq y} \left( \frac{1}{3} \times \frac{1}{3} + \frac{2}{3} \times \frac{2}{3} \right) \\ &= n + \frac{5}{9}n(n-1) = \frac{5n^2 + 4n}{9} \end{aligned}$$