

CIS 551 / TCOM 401

Computer and Network Security

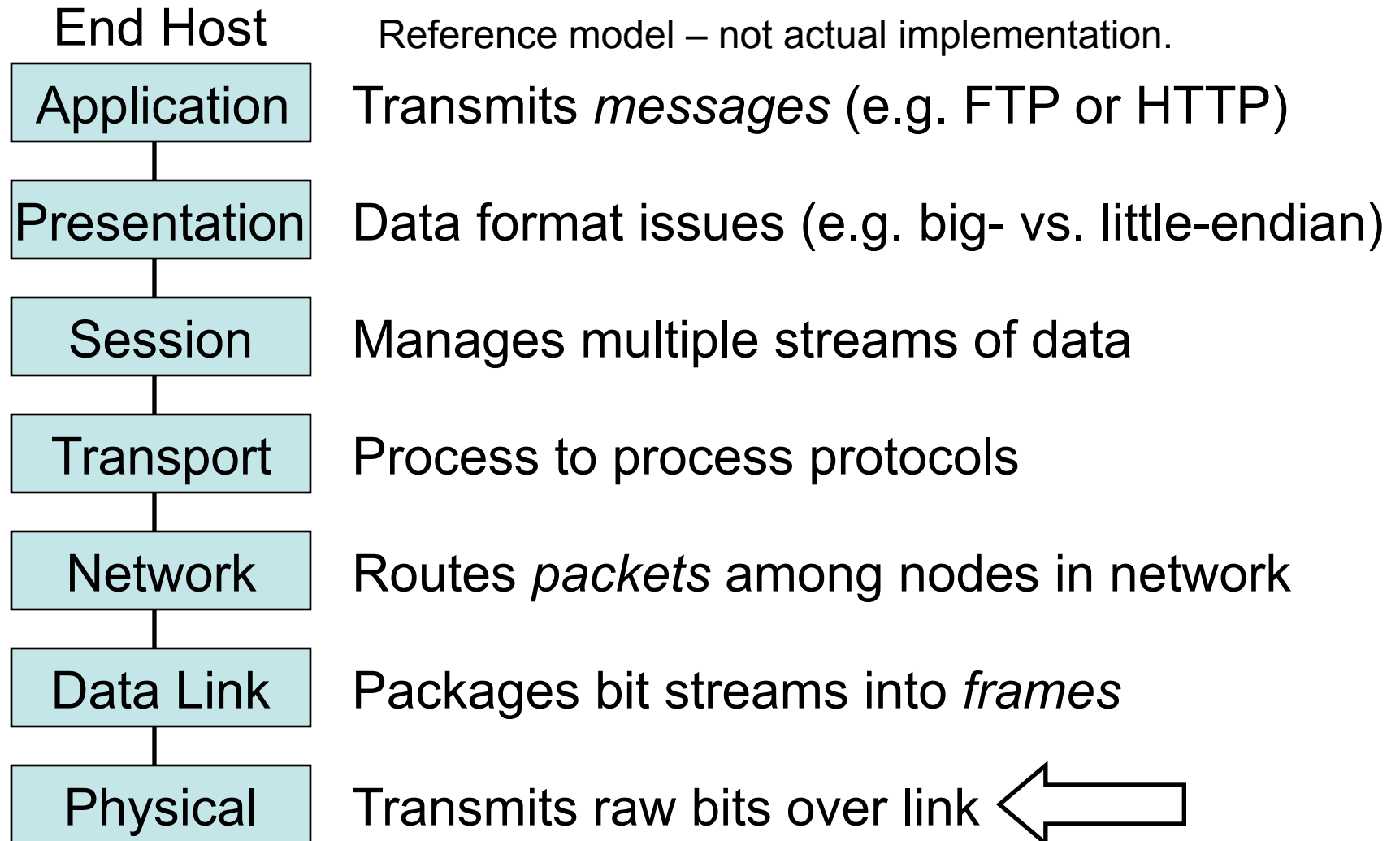
Spring 2009

Lecture 6

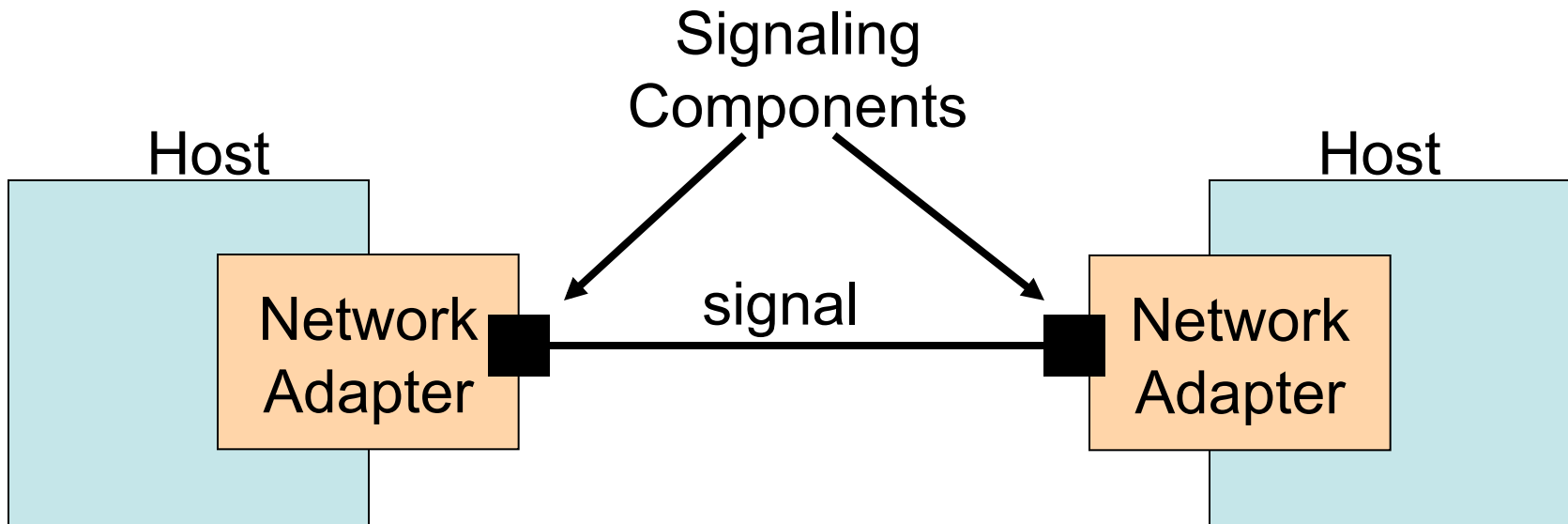
Announcements

- First project: Due: 6 Feb. 2009 at 11:59 p.m.
- <http://www.cis.upenn.edu/~cis551/project1.html>
- Plan for Today:
 - Networks: Lower layers

Open Systems Interconnection (OSI)



Signaling Components



Network adapters encode streams of bits into signals.

Simplification: Assume two discrete signals—high and low.

Practice: Two different voltages on copper link or different brightness of light on fiber link.
(leads to some interesting encoding issues)

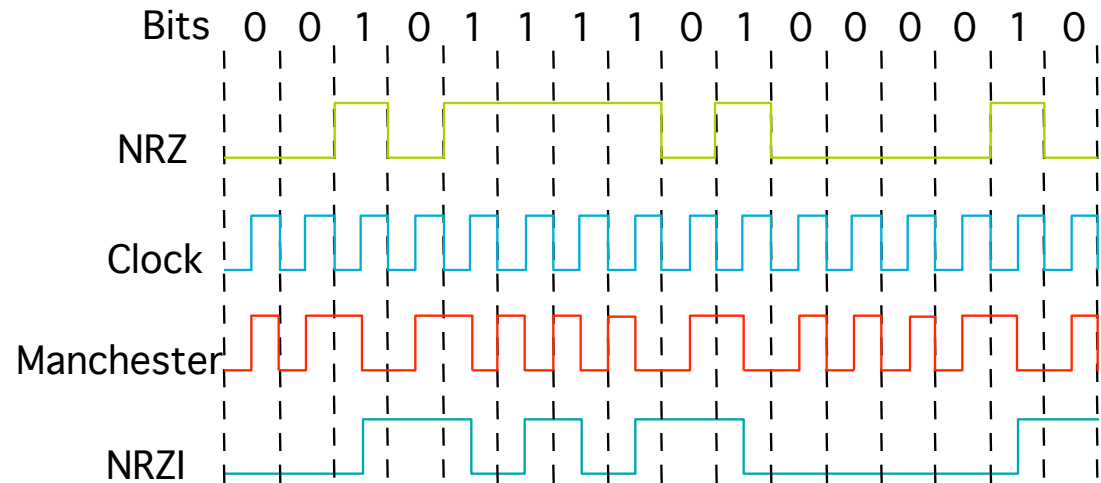
Not in this course

$$\nabla \cdot \vec{E} = \epsilon_0 \rho$$

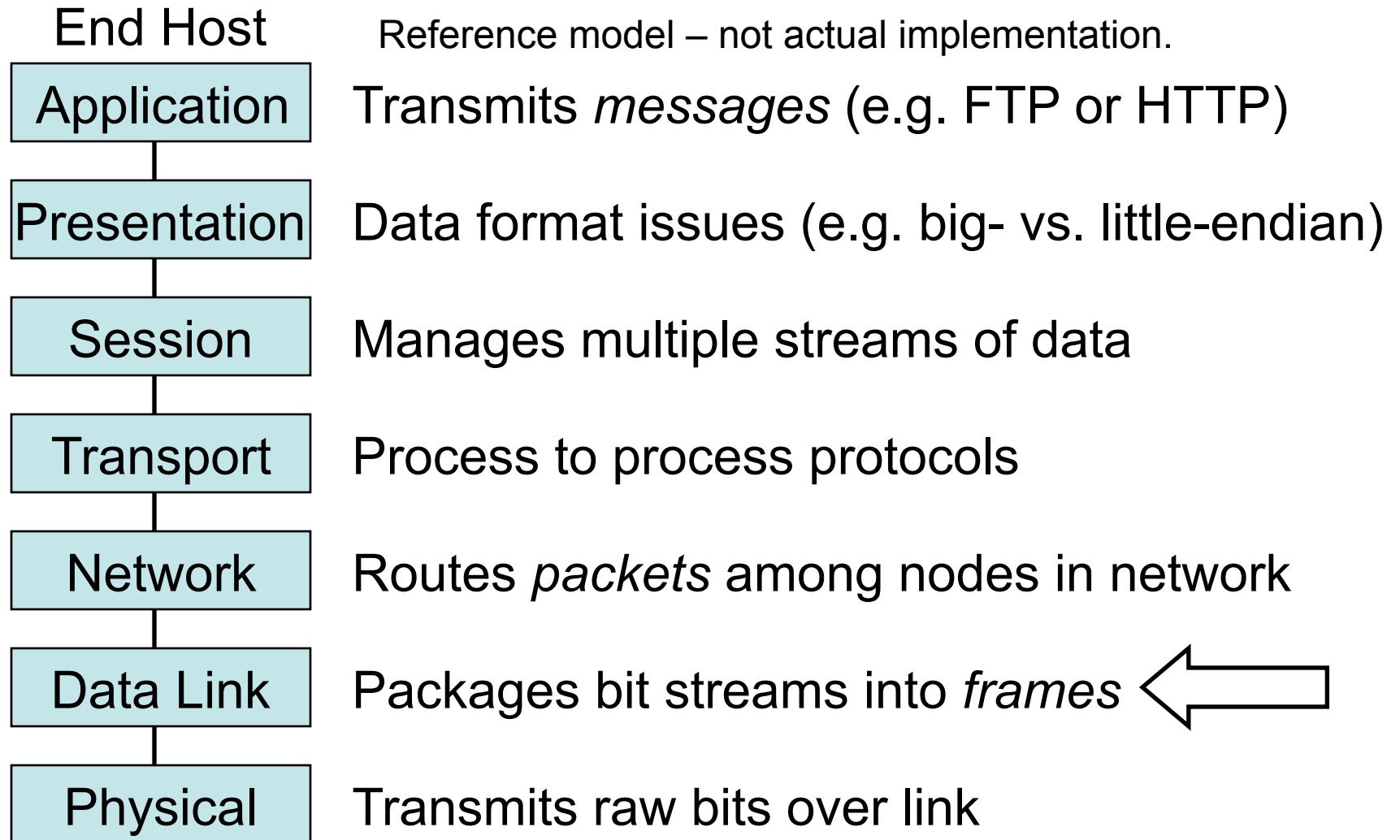
$$\nabla \times \vec{E} = -\frac{\partial \vec{B}}{\partial t}$$

$$\nabla \cdot \vec{B} = 0$$

$$\nabla \times \vec{B} = \mu_0 \vec{J} + \mu_0 \epsilon_0 \frac{\partial \vec{E}}{\partial t}$$



Open Systems Interconnection (OSI)

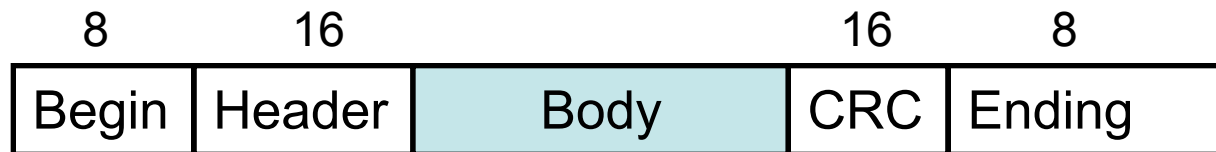


Framing

- Need a way to send blocks of data.
 - How does the network adapter detect when the sequence begins and ends?
 - Are there transmission errors in the data?
- *Frames* are link layer unit of data transmission
 - Byte oriented vs. Bit oriented
 - Point-to-point (e.g. PPP) vs. Multiple access (Ethernet)

A Multi-access, Bit-oriented Protocol

- Frames contain sequences of bits
 - Could be ASCII
 - Could be pixels from an image
- Frames read by many nodes
 - Address distinguishes intended recipient
- HDLC (High-level Data Link Control)
 - Begin and ending = 01111110
 - Uses *bit stuffing*: suffix five 1's with a 0



HDLC frame format

Problem: Error Detection & Correction

- Bit errors may be introduced into frames
 - Electrical interference
 - Thermal noise
- Could flip one bit or a few bits independently
- Could zero-out or flip a sequence of bits (*burst error*)

- How do you detect an error?

- What do you do once you find one?

Error Detection

- General principal: Introduce redundancy
- Trivial example: send two copies
 - High overheads: $2n$ bits to send n
 - Won't detect errors that corrupt same bits in both copies
- How can we do better?
 - Minimize overhead
 - Detect many errors
 - General subject: error detecting codes

Simple Error Detection Schemes

- Parity
 - 7 bits of data
 - 8th bit is sum of first seven bits mod 2
 - Overhead: 8n bits to send 7n
 - Detects: any odd number of bit errors
- Internet Checksum algorithm
 - Add up the words of the message, transmit sum
 - 16 bit ones-complement addition
 - Overhead: 16 bits to send n
 - Does not detect all two bit errors

Cyclic Redundancy Check

- Reading: Wikipedia entry on CRC
- Used in link-level protocols
 - CRC-32 used by Ethernet, 802.5, PKzip, ...
 - CRC-CCITT used by HDLC
 - CRC-8, CRC-10, CRC-32 used by ATM
- Better than parity or checksum
 - (e.g. 32 bits to send 12000)
- Simple to implement

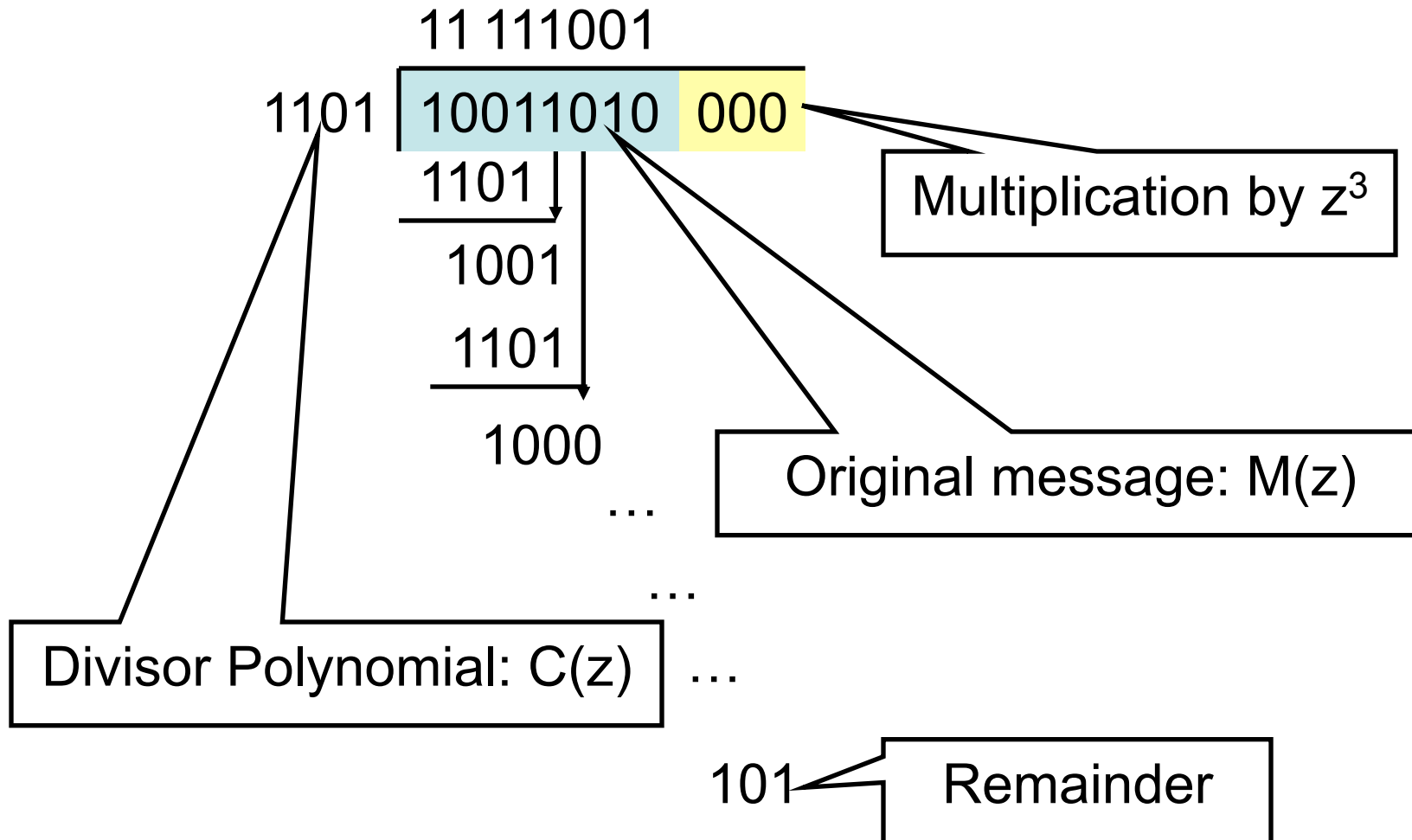
Cyclic Redundancy Check (CRC)

- Consider $(n+1)$ -bit message as a n -degree polynomial
 - Polynomial arithmetic modulo 2
 - Bit values of message are coefficients
 - Message = 10011010
 - Polynomial
$$M(z) = (1 \cdot z^7) + (0 \cdot z^6) + (0 \cdot z^5) + (1 \cdot z^4) + (1 \cdot z^3) + (0 \cdot z^2) + (1 \cdot z^1) + (0 \cdot z^0)$$
$$= z^7 + z^4 + z^3 + z^1$$

Cyclic Redundancy Check

- Sender and receiver agree on a *divisor polynomial* $C(z)$ of degree k
 - Example $k = 3$
 - $C(z) = z^3 + z^2 + 1$
 - Coefficients are 1101
- Error correction bits are remainder of
 $(M(z) \cdot z^k)$ divided by $C(z)$
- This yields a $n+k$ bit transmission polynomial $P(z)$ that is *exactly* divisible by $C(z)$

Example CRC Calculation



Example CRC Calculation

$$\begin{array}{r} Z^3 \cdot \text{Original Message } M(z) = \quad 10011010 \quad 000 \\ \text{Remainder} = \quad + \quad \quad \quad 101 \\ \hline \text{Transmitted message } P(z) = \quad 10011010 \quad 101 \end{array}$$

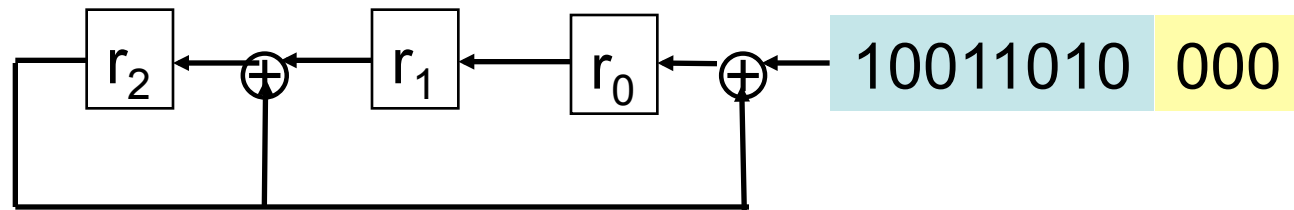
- Recipient checks that $C(z)$ evenly divides the received message.

CRC Error Detection

- Must choose a good divisor $C(z)$
 - There are many standard choices:
CRC-8, CRC-10, CRC-12, CRC-16, CRC-32
 - CRC-32: 0x04C11DB7
- All 1-bit errors as long as z^k and z^0 coefficients are 1
- All 2-bit errors as long as $C(z)$ has three terms
- Any odd number of errors if $(z+1)$ divides $C(z)$
- Any burst errors of length $\leq k$

CRC Implementations

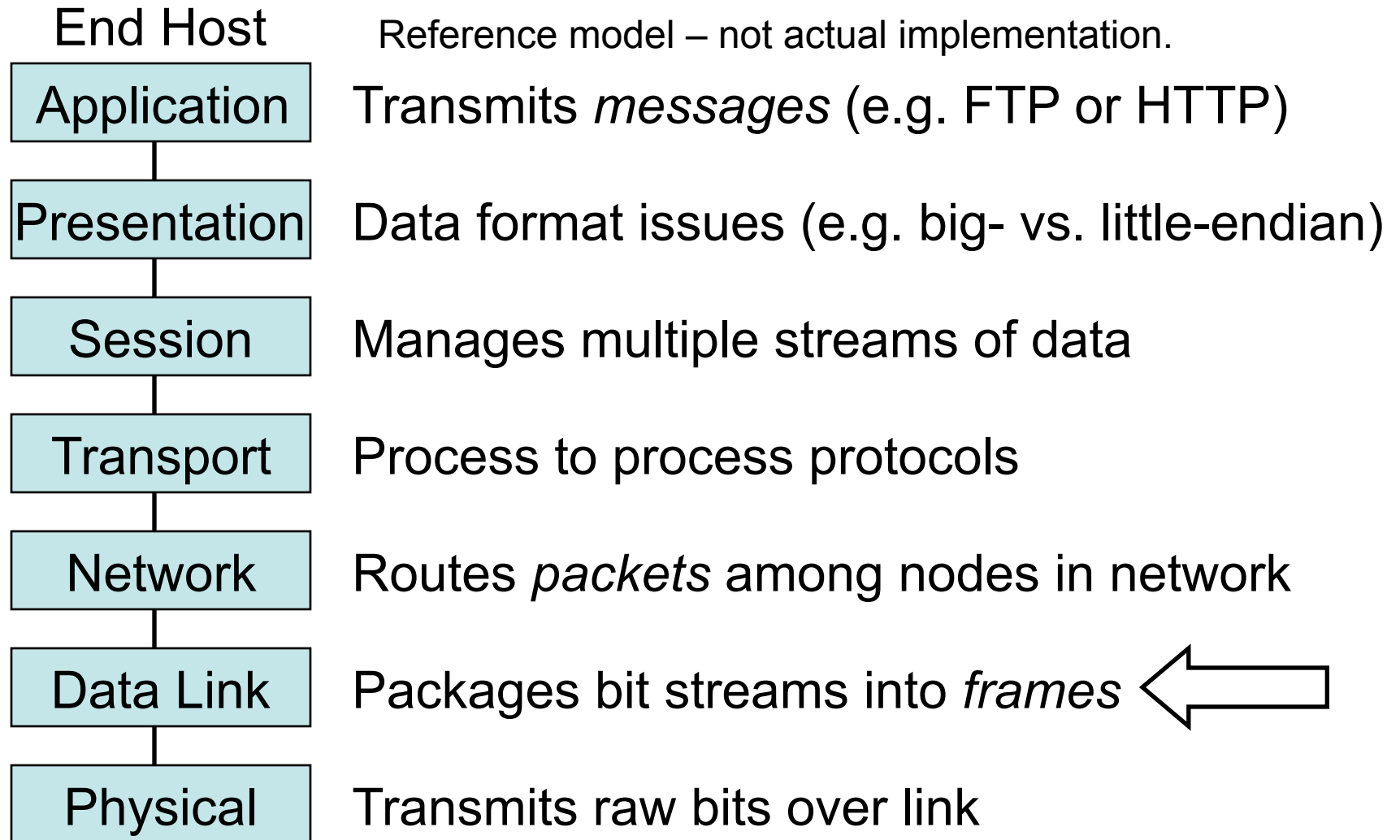
- Easy to implement in hardware
 - Base 2 subtraction is XOR
 - Simple k-bit shift register with XOR gates inserted before 1's in $C(z)$ polynomial
 - Message is shifted in, registers fill with remainder
- Example $C(z) = 1101$



Error Correction Codes

- Redundant information can be used to *correct* some errors
- Typically requires more redundancy
- Tradeoffs:
 - Error detection requires retransmission
 - Error correction sends more bits all the time
- Forward Error Correction is useful:
 - When errors are likely (e.g. wireless network)
 - When latency is too high for retransmission (e.g. satellite link)

Open Systems Interconnection (OSI)



IEEE 802 network standards

The IEEE 802 committee produces standards & specifications for Local Area Networks (LAN):

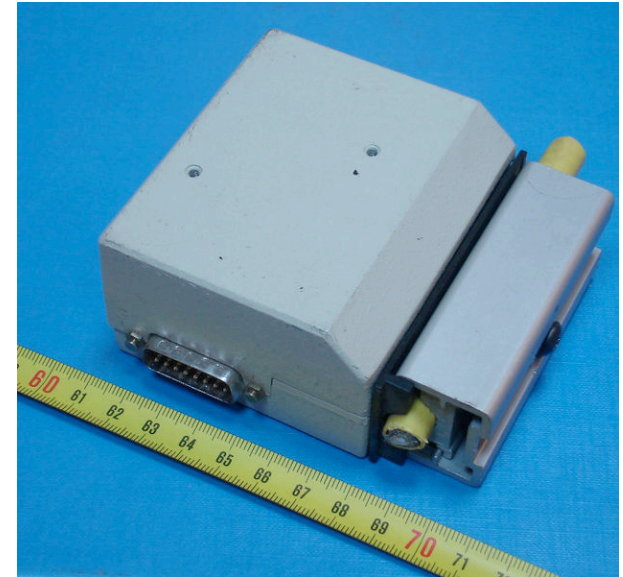
- **802.3 CSMA/CD Networks (Ethernet)**
- 802.4 Token Bus Networks
- 802.5 Token Ring Networks
- 802.6 Metropolitan Area Networks
- **802.11 Wireless LAN (Wifi) [Thursday]**

Ethernet (802.3)

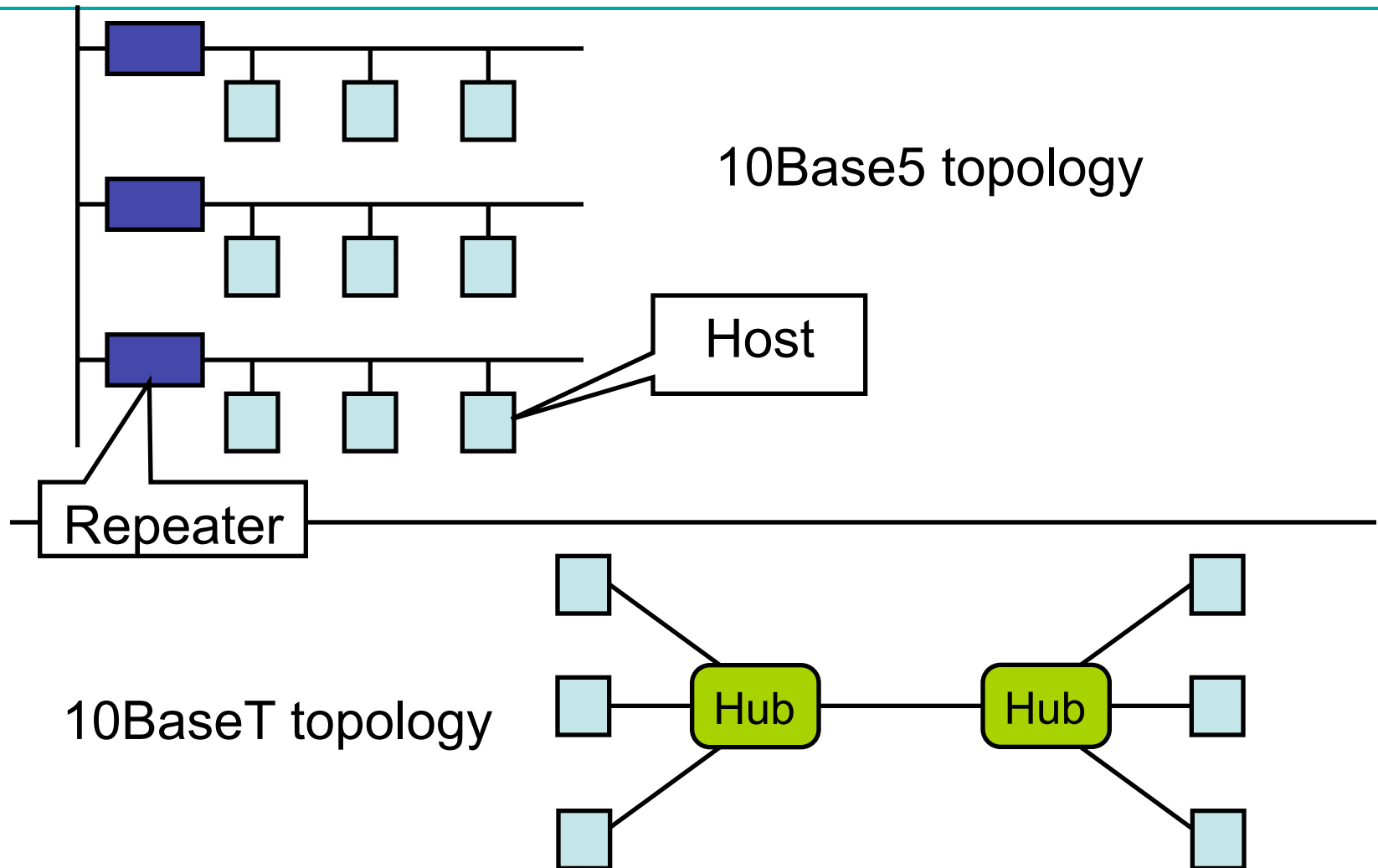
- A standard for local area networks (LAN)
- Developed in mid-70's at Xerox PARC
 - Descendent of Aloha, a U. of Hawaii radio packet network
 - DEC, Intel, and Xerox standard: 1978 for 10Mbps
 - IEEE 802.3 standard grew out of that
- Physical implementations:
 - 10Base5, 10BaseT, 100BaseT, 1000BaseT...
 - Speed: 10Mbps, 100Mbps, 1000Mbps, ...

Ethernet Physical links

- Originally used “Thick-net” 10Base5
 - 10 = 10Mbps
 - 5 = maximum of 500 meters segments
 - Up to 4 repeaters between two hosts
=2500m max
- More common: 10BaseT
 - 10 = 10Mbps
 - T = Twisted pair (typically Category 5),
Maximum of 100 meter segments
 - Connected via *hubs* (still 2500m max)
- Today’s standards: 100BaseT, 1000BaseT



Ethernet topologies



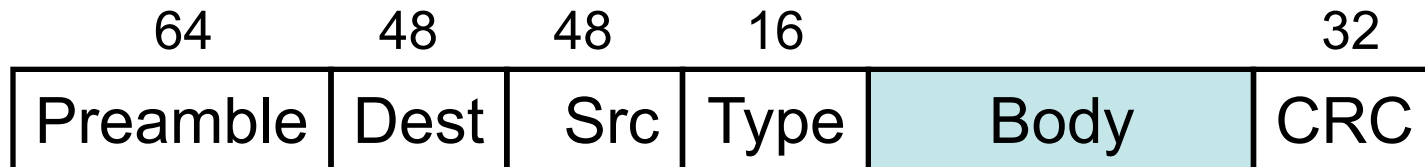
How the ethernet works

- The Ethernet link is *shared*
 - A signal transmitted by one host reaches *all* hosts
- Method of operation: **CSMA/CD**
 - Carrier Sense, Multiple Access, with Collision Detection
- Hosts competing for the same link are said to be in the same ***collision domain***
 - Good news: easy to exchange data
 - Bad news: have to regulate link access
- Protocol: *Media Access Control (MAC)*

Ethernet Addresses

- Every adapter manufactured has a unique address
 - 6 bytes (48 bits) usually written in Hex.
 - Examples: 00-40-50-B1-39-69 and 8:0:2b:e4:b1:2
 - Each manufacturer is assigned 24bit prefix
 - Manufacturer ensures unique suffixes

Ethernet Frame Format



- Preamble – repeating pattern of 0's & 1's
 - Used by receiver to synchronize on signal
- Dest and Src – Ethernet Addresses
- Type – demultiplexing key
 - Identifies higher-level protocol
- Body – payload
 - Minimum 46 Bytes
 - Maximum 1500 Bytes

Addresses in an ethernet frame

- All bits = 1 indicates a *broadcast* address
 - Sent to all adapters
- First bit = 0 indicates *unicast* address
 - Sent to only one receiver
- First bit = 1 indicates *multicast* address
 - Sent to a group of receivers

An Ethernet Adapter Receives:

- Frames addressed to the broadcast address
- Frames addressed to its own address
- Frames sent to a multicast address
 - If it has been programmed to listen to that address
- All frames
 - If the adapter has been put into *promiscuous mode*

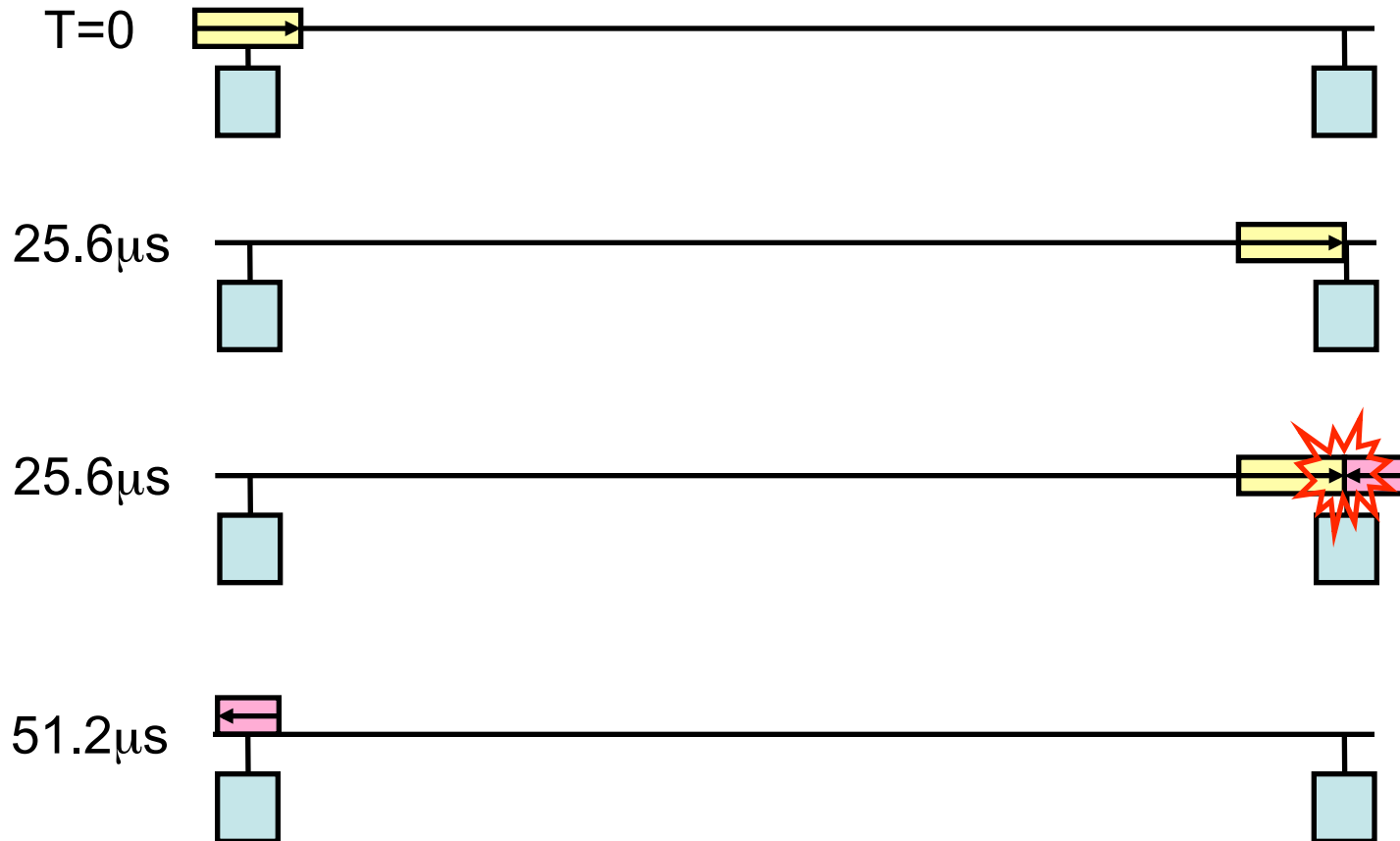
Ethernet Transmitter Algorithm

- If the link is idle transmit the frame immediately
 - Upper bound on frame size means adapter can't hog the link
- If the link is busy
 - Wait for the line to go idle
 - Wait for $9.6\mu\text{s}$ after end of last frame (sentinel)
 - Transmit the frame
- Two (or more) frames may *collide*
 - Simultaneously sent frames interfere

Collision Detection

- When an adapter detects a collision
 - Immediately sends 32 bit *jamming signal*
 - Stops transmitting
- A 10Mbps adapter may need to send 512 bits in order to detect a collision
 - Why?
 - 2500m + 4 repeaters gives RTT of $51.2\mu\text{s}$
 - $51.2\mu\text{s}$ at 10Mbps = 512 bits
 - Fortunately, minimum frame (excluding preamble) is 512 bits = 64 bytes
 - 46 bytes data + 14 bytes header + 4 bytes CRC

Ethernet Collision (Worst Case)



Exponential Backoff

- After it detects 1st collision
 - Adapter waits either 0 or 51.2 μ s before retrying
 - Selected randomly
- After 2nd failed transmission attempt
 - Adapter randomly waits 0, 51.2, 102.4, or 153.6 μ s
- After nth failed transmission attempt
 - Pick k in 0 ... 2ⁿ-1
 - Wait k x 51.2 μ s
 - Give up after 16 retries
(but cap n at 10)

Ethernet Security Issues

- Promiscuous mode
 - *Packet sniffer* detects all Ethernet frames
- Less of a problem in *switched* Ethernet
 - Why?