Graph Coloring: Background and Assignment

Networked Life
NETS 112
Fall 2014
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A Little Experiment

• Consider the network in which you are connected to the people sitting to your left, right, in front of and behind you
  - if you have no neighbors, change your seat
  - if no one is sitting to your (say) left, you’re just “missing“ that neighbor
• Hold up 1, 2, or 3 fingers
• Your “goal” is to pick a different number than all your neighbors
  - can change your number whenever you want, as many times as you want
• While your number matches at least one neighbor, raise your other hand
Graph Coloring

- We are given a graph or network $G$ with $N$ vertices (e.g. $N \sim 100$)
- We are given $K$ values or “colors” (e.g. $K = 3$)
- We would like to find a labeling of vertices by colors such that for every edge $(u,v)$ in $G$, $u$ and $v$ have *different colors*
- In general, for any given $G$, this problem is harder the smaller $K$ is
- *Chromatic number* $K(G) = \text{smallest } K \text{ for which there is a solution}
How Does Structure of G Influence K(G) (and Good Algorithms?)

- How many colors are always enough?
- If some vertex has degree D, must K(G) ≥ D?
- What’s a sufficient condition to force K(G) ≥ 3? More generally?
- Is it a necessary condition?
- How did I know 2 colors was enough for the grid?
- What’s a good algorithm if you know K(G) = 2?
- Is there any “local” property/algorithm determining K(G)/solutions?
A Famous, Important and Notorious Problem

- **Optimization** (minimizing scarce resource): e.g. exam scheduling
  - one vertex for each Penn class: NWLife, CIS 120, Intro to Chocolate,…
  - draw an edge between two classes if there is a student taking both
  - colors = final exam time slots
  - solution ensures no student has simultaneous exams!
  - would like to minimize the length of exam period…
- **Cartography**: 4-color theorem
- **Graph Coloring** as a model of social differentiation
- **Graph Coloring** is a computationally hard problem:
  - algorithm given graph (list of vertices and edges) as input
  - output a proper coloring (solution) with smallest number of colors \( K(G) \)
  - best known algorithms not much better than exhaustive search
  - running time scale exponentially in \( N \), e.g. \( 10^N \)
  - For \( N=100 \): \( 10^{100} \gg \) number of protons in the universe
  - centralized or “birds-eye” computation, vs distributed & local
- **Even significant relaxations remain intractable**
  - e.g. allow (much) more than \( K(G) \) colors
Your Assignment

• Later in the course, we will study experiments on human subjects solving graph coloring from local, distributed information
• For now, you are asked to try some coloring problems on your own
• Will employ a web app we have designed and developed
• Your score will give points for simply finding solutions, with bonus points for finding them quickly
• No collaboration/collusion of any kind --- not even discussion
• Assignment is due (app will close) in one week (midnight Tue 9/16)