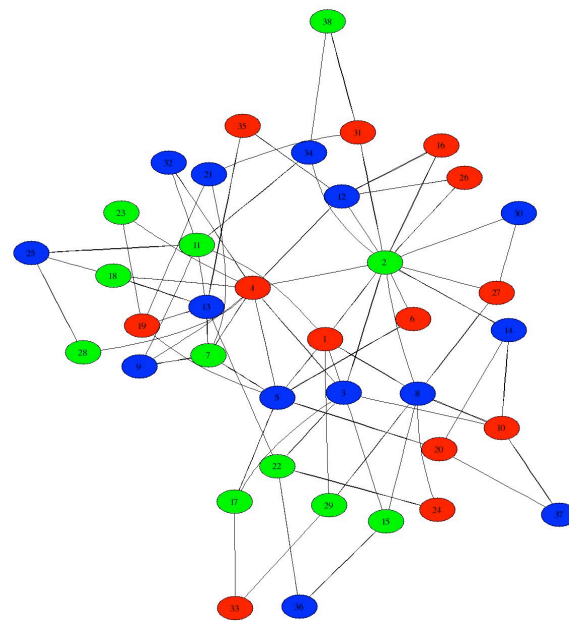


# Graph Coloring: Background and Assignment

Networked Life  
NETS 112  
Fall 2014  
Prof. Michael Kearns

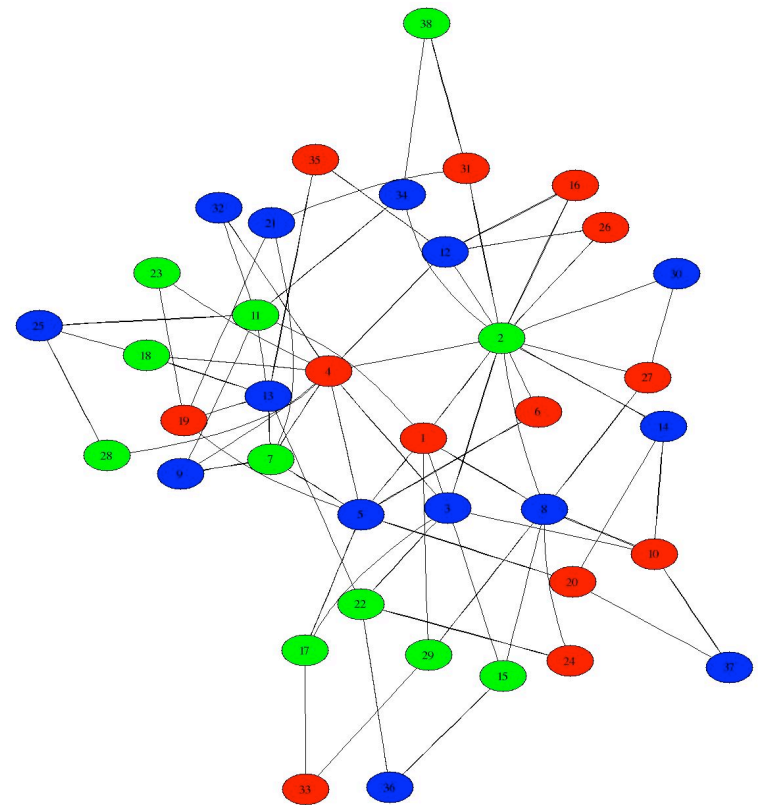
# A Little Experiment

- Consider the network in which you are connected to the people sitting to your left, right, in front of and behind you
  - if you have no neighbors, change your seat
  - if no one is sitting to your (say) left, you're just "missing" that neighbor
- Hold up 1, 2, or 3 fingers
- Your "goal" is to pick a *different* number than all your neighbors
  - can change your number whenever you want, as many times as you want
- While your number matches at least one neighbor, raise your other hand



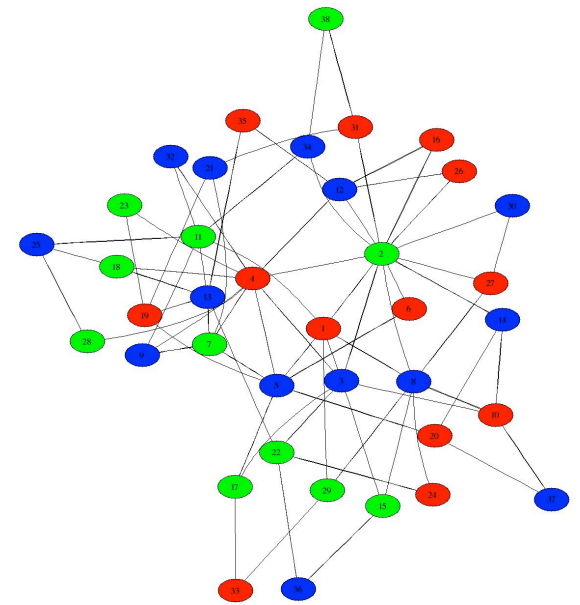
# Graph Coloring

- We are given a graph or network  $G$  with  $N$  vertices (e.g.  $N \sim 100$ )
- We are given  $K$  values or "colors" (e.g.  $K = 3$ )
- We would like to find a labeling of vertices by colors such that for every edge  $(u,v)$  in  $G$ ,  $u$  and  $v$  have *different colors*
- In general, for any given  $G$ , this problem is harder the smaller  $K$  is
- *Chromatic number*  $K(G) =$  smallest  $K$  for which there is a solution



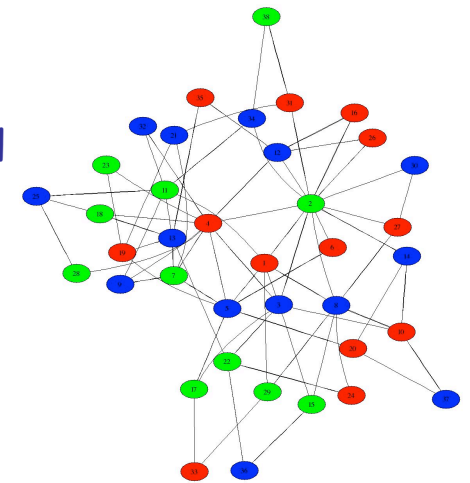
## How Does Structure of $G$ Influence $K(G)$ (and Good Algorithms?)

- How many colors are always enough?
- If some vertex has degree  $D$ , must  $K(G) \geq D$ ?
- What's a sufficient condition to force  $K(G) \geq 3$ ? More generally?
- Is it a necessary condition?
- How did I know 2 colors was enough for the grid?
- What's a good algorithm if you know  $K(G) = 2$ ?
- Is there any "local" property/algorithm determining  $K(G)$ /solutions?



# A Famous, Important and Notorious Problem

- Optimization (minimizing scarce resource): e.g. exam scheduling
  - one vertex for each Penn class: NWLife, CIS 120, Intro to Chocolate,...
  - draw an edge between two classes if there is a student taking both
  - colors = final exam time slots
  - solution ensures no student has simultaneous exams!
  - would like to minimize the length of exam period...
- Cartography: 4-color theorem
- Graph Coloring as a model of social differentiation
- Graph Coloring is a computationally hard problem:
  - algorithm given graph (list of vertices and edges) as input
  - output a proper coloring (solution) with smallest number of colors ( $K(G)$ )
  - best known algorithms not much better than exhaustive search
  - running time scale exponentially in  $N$ , e.g.  $\sim 10^N$
  - For  $N=100$ :  $10^{100} \gg$  number of protons in the universe
  - centralized or "birds-eye" computation, vs distributed & local
- Even significant relaxations remain intractable
  - e.g. allow (much) more than  $K(G)$  colors



# Your Assignment

- Later in the course, we will study experiments on human subjects solving graph coloring from local, distributed information
- For now, you are asked to try some coloring problems on your own
- Will employ a web app we have designed and developed
- Your score will give points for simply finding solutions, with bonus points for finding them quickly
- No collaboration/collusion of any kind --- not even discussion
- Assignment is due (app will close) in one week (midnight Tue 9/16)

