

Many problems in computer vision and Machine Learning can be formulated using **Markov Random Fields**

proposed approach:

- no restrictions** on the pairwise clique potentials
- complexity** $O(\#\text{edges} \#\text{labels}^2)$, **linear** in the description length of the clique potentials
- improved general optimality bounds**

MRFs and IQP/QP formulation

$$\max P(X) = \frac{1}{Z} \prod_{ij \in E} \Psi_{ij}(X_i, X_j) \prod_i \Phi_i(X_i)$$

$$X_i \in \{1, \dots, k\}$$

Proposition: QP is equivalent to IQP

(generalizes a result from Ravikumar & Lafferty, 2006)
given a solution to IQP, we can construct a solution to QP

IQP formulation

$$x_{ia} = 1 \text{ iff } X_i = a \quad \epsilon(x) = x^T W x + V^T x$$

$$\max \epsilon(x) \quad \text{s.t.} \quad Cx = 1, x \in \{0, 1\}^{nk}$$

↓ discrete set Ω_d

QP relaxation

$$\max \epsilon(x) \quad \text{s.t.} \quad Cx = 1, 0 \leq x \leq 1$$

non-convex! simplex Ω_s

Solving Affine Constrained Rayleigh Quotients

$$\max_x \frac{x^T W x + V^T x}{x^T x + \beta} \quad \text{s.t.} \quad Cx = b$$

EQUIVALENT to IQP for x binary

Linear Constraint: $Cx = 0$ (and $V = 0, \beta = 0$)

Yu and Shi, 2001

Affine Constraint: $Cx = b \iff \sum_a x_{ia} = 1$

\implies **SQP**

Inequality Constraint? $Cx \leq b$

Theorem: this is NP-HARD

Solution

1. rewrite as $\max_{x,t} \frac{x^T W x + t V^T x}{x^T x + \beta t^2} \quad \text{s.t.} \quad Cx = t b$ then $\max \frac{\bar{x}^T \bar{W} \bar{x}}{\bar{x}^T \bar{D} \bar{x}} \quad \text{s.t.} \quad \bar{C} \bar{x} = 0$

2. introduce $x' = \bar{D}^{1/2} \bar{x} \implies \max \epsilon_1(x') = \frac{x'^T W' x'}{x'^T x'} \quad \text{s.t.} \quad C' x' = 0$

Efficient computation with Sherman-Morrison formula

3. solve $P_C W' P_C x' = \lambda x'$ with $P_C = I - C'^T (C' C'^T)^{-1} C'$

Complexity (per eigensolver iteration): $O(nnz(W')) = O(nnz(W)) = O(k^2 |E|)$

Spectral Relaxation to the QP (SQP)

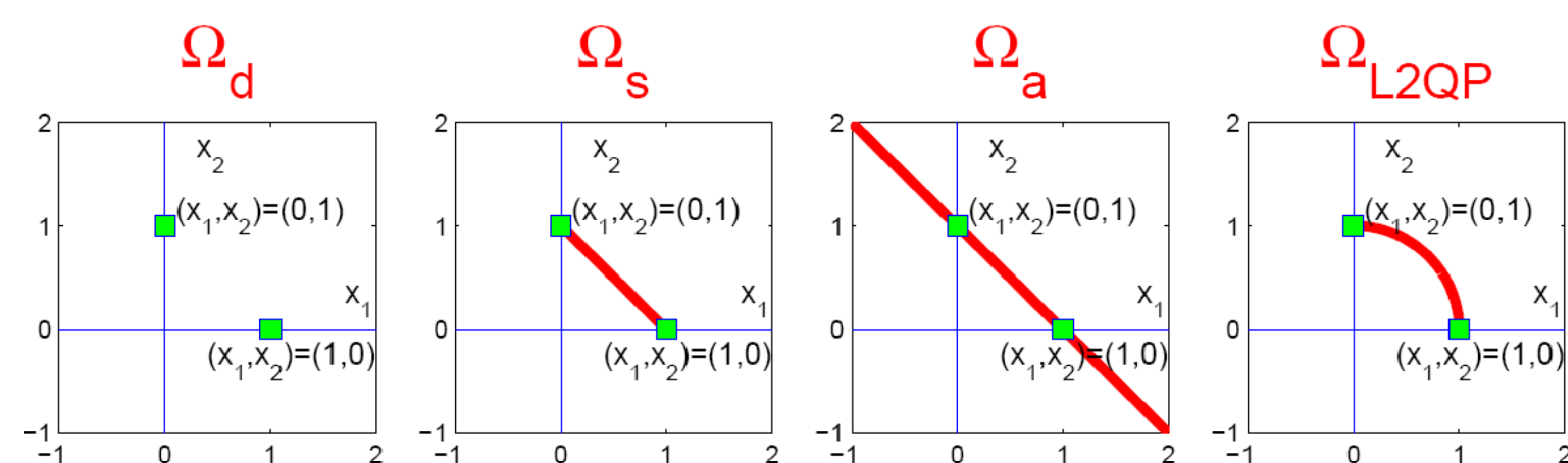
$$\max \epsilon_S(x) = \frac{x^T W x + V^T x}{x^T x + \beta} \quad \text{s.t.} \quad Cx = 1$$

normalization encourages $\|x\|$ small, and $x \in [0, 1]$ affine space Ω_a



still non-convex! but solvable

different constraints



Algorithm

- Input: clique potentials W, V
- Set $\beta = \hat{\beta}$
- Compute x_S from x' , first eigenvector of $W_{eq} = P_C W' P_C$
- Output upper bound $\epsilon^* \leq \frac{n+\beta}{x_S^T x_S + \beta} \epsilon(x_S)$
- Discretize using Relaxation Labeling $\implies x_d$
- Output lower bound $\epsilon^* \geq \epsilon(x_d)$

upper/lower bounds

General upper/lower bounds

$$\forall x \in \Omega_d, \quad \frac{1}{n+\beta} \epsilon(x) = \epsilon_S(x)$$

$$\forall x \in \Omega_s, \quad \frac{1}{n+\beta} \epsilon(x) \leq \epsilon_S(x) \leq \frac{1}{n/k + \beta} \epsilon(x)$$

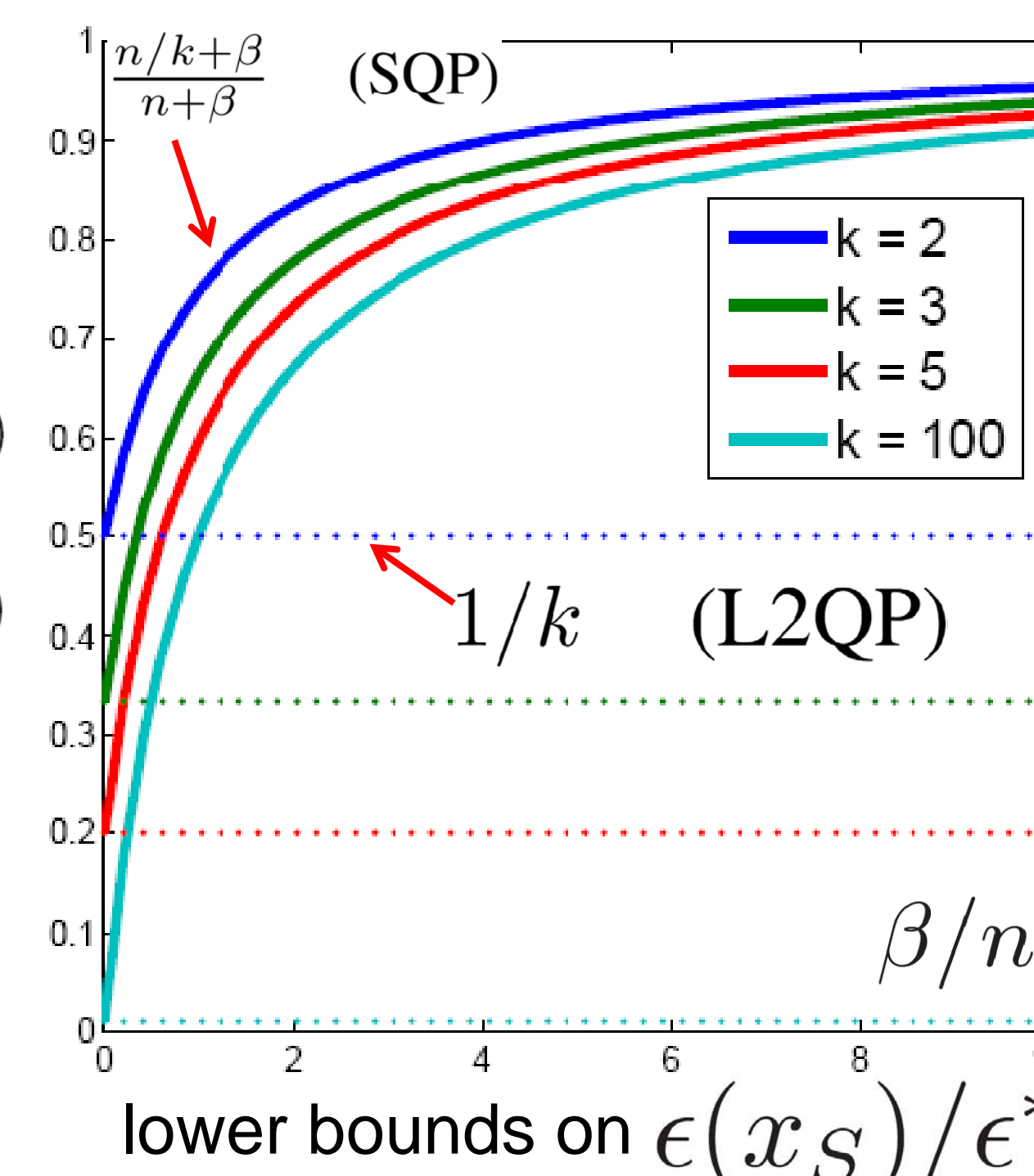
$$\forall x \in \Omega_a, \quad \epsilon_S(x) \leq \frac{1}{n/k + \beta} \epsilon(x)$$

($x^T x \geq n/k$, etc...)

From IQP \Leftrightarrow QP and the above, we get:

when $x_S \geq 0$,

$$\epsilon(x_S) \leq \epsilon(x_d) \leq \epsilon^* \leq \frac{n + \beta}{n/k + \beta} \epsilon(x_S)$$



Data dependent lower bound

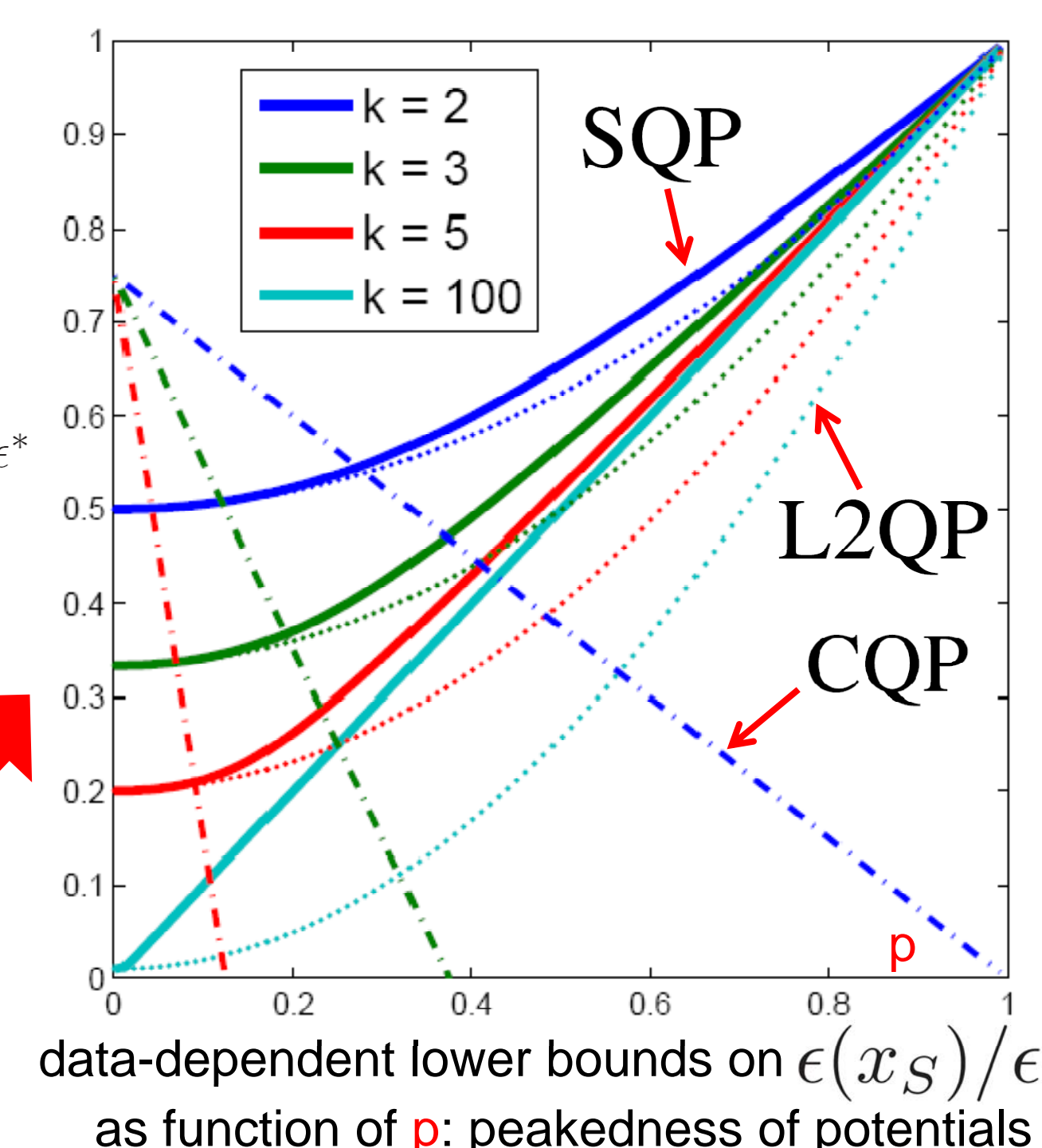
$$M = W + \text{diag}(V)$$

$$\forall i, x_{i1} = 1 \text{ and } 0 \text{ otherwise} \quad \mathbf{1}^T M_{1,1} \mathbf{1} = x^{*T} M x^* = \epsilon^*$$

$$M = \begin{bmatrix} M_{1,1} & M_{1,2:k} \\ M_{1,2:k}^T & M_{2:k,2:k} \end{bmatrix}$$

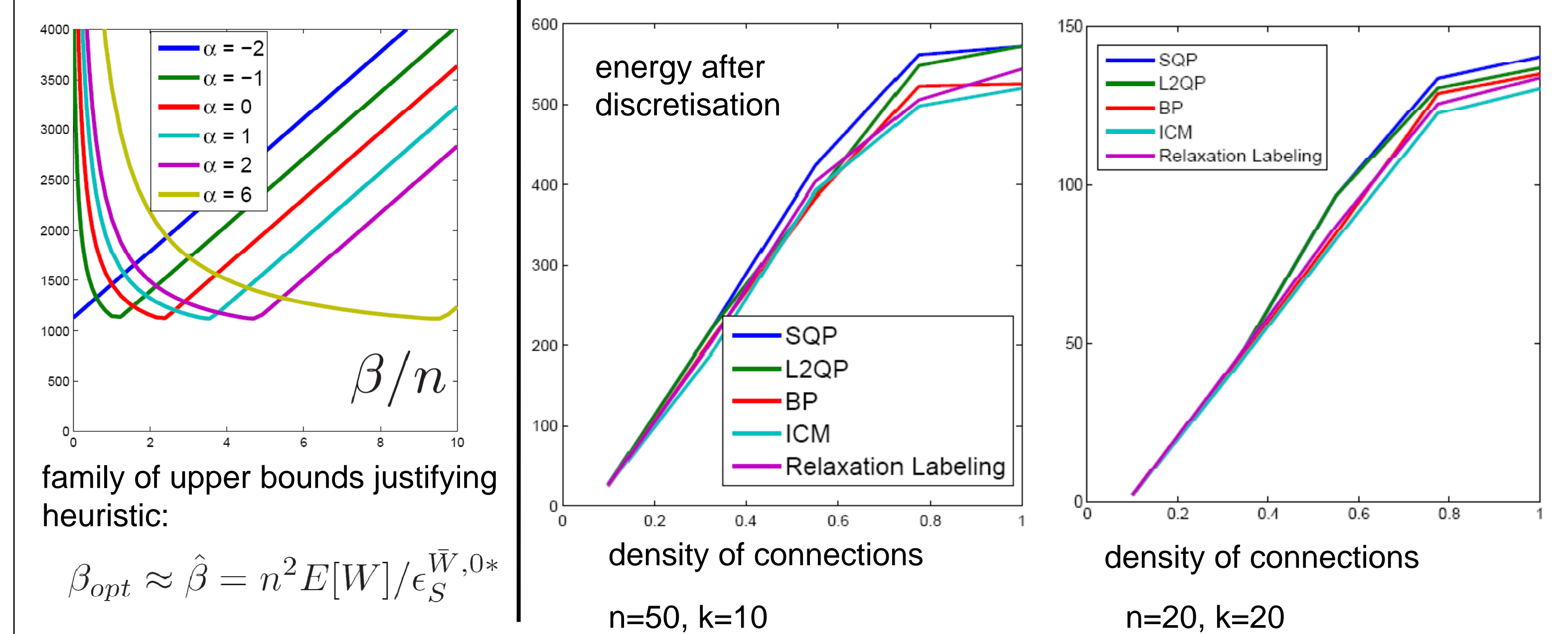
p : the largest element in $[0, 1]$ with $pE[M_{1,1}] \leq E[M_{2:k,2:k}]$ (peakedness)
 $pE[M_{1,1}] \leq E[M_{1,2:k}]$

$$\epsilon(x_S) \geq f(p, k) \epsilon^*$$



Experiments

Comparison between SQP, L2QP, BP, ICM, Relaxation Labeling on random MRF problems with controlled parameters



Previous work

Linear relaxations: LP, SDP, SOCP

QP rewritten as (linear) matrix inner product, approximating rank 1 constraint $x^T W x + V^T x = \langle X, W_{eq} \rangle$ where $X = [x; 1][x; 1]^T$

Quadratic relaxations: L2QP and CQP

CQP: convexification of objective using $(W, V) \implies (W - \text{diag}(D), V + D)$

L2QP: $Cx=1$ relaxed to $\sum_a x_{ia}^2 = 1$