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# CSE331: Introduction to Networks and Security

Lecture 3  
Fall 2006

# Announcements

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- First project: Due: 25 Sept. 2006
  - <http://www.cis.upenn.edu/~cse331/project1.html>
  - Please e-mail [savi@seas.upenn.edu](mailto:savi@seas.upenn.edu) with your project group by Weds. 13 Sept.
  - If you need a group, send e-mail
- Please put "cse331" in the subject of all course-related e-mail
- Prof. Zdancewic will be away Sept. 18 & 20.
  - Class will be taught by Peng Li



# Plan for Today:

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- Wrap up buffer overflows discussion.
- Take a step back and begin talking about networks.
  - Basic network architecture
  - Terminology
  - Simple performance characteristics

# Tool support for C/C++

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- Extensions to gcc that do array bounds checking
- Link against "safe" versions of libc (e.g. libsafe)
- Test programs with tools such as Purify, Splint, valgrind
- Compile programs using tools such as:
  - Stackguard and Pointguard (Cowan et al., immunix.org)
- Research compilers:
  - Ccured (Necula et al.)
  - Cyclone (Morrisett et al.)
- Binary rewriting techniques
  - Software fault isolation (Wahbe et al.)

# Defeating Buffer Overflows

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- Use a typesafe programming language
  - Java/C# are not vulnerable to these attacks
  - Garbage collection eliminates most memory management problems
- Some operating systems move the start of the stack on a per-process basis:
  - E.g. eniac-l
  - Doesn't provide complete protection (but it does make things harder)

# The Four Major Networks

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- Telephone
  - Television
  - Radio
- } Special Purpose
- Internet (grew out of ARPANET—late 1960's)
  - Starting to see hybrids...
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- Computer networks
    - General purpose programmable hardware
    - Support many different applications



# How to build such a network?

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- Connectivity
- Efficient Resource Sharing
- Functionality
- Performance
- *Security*

# Requirement: Connectivity

- Goal of a network is to get information from one place to another
  - Source
  - Destination } Specified by an *address*
  - *Nodes or Hosts*
- Network paths
  - Can be *direct* or *indirect*
  - Can be *static* or *dynamic*

# Connectivity: Direct Links

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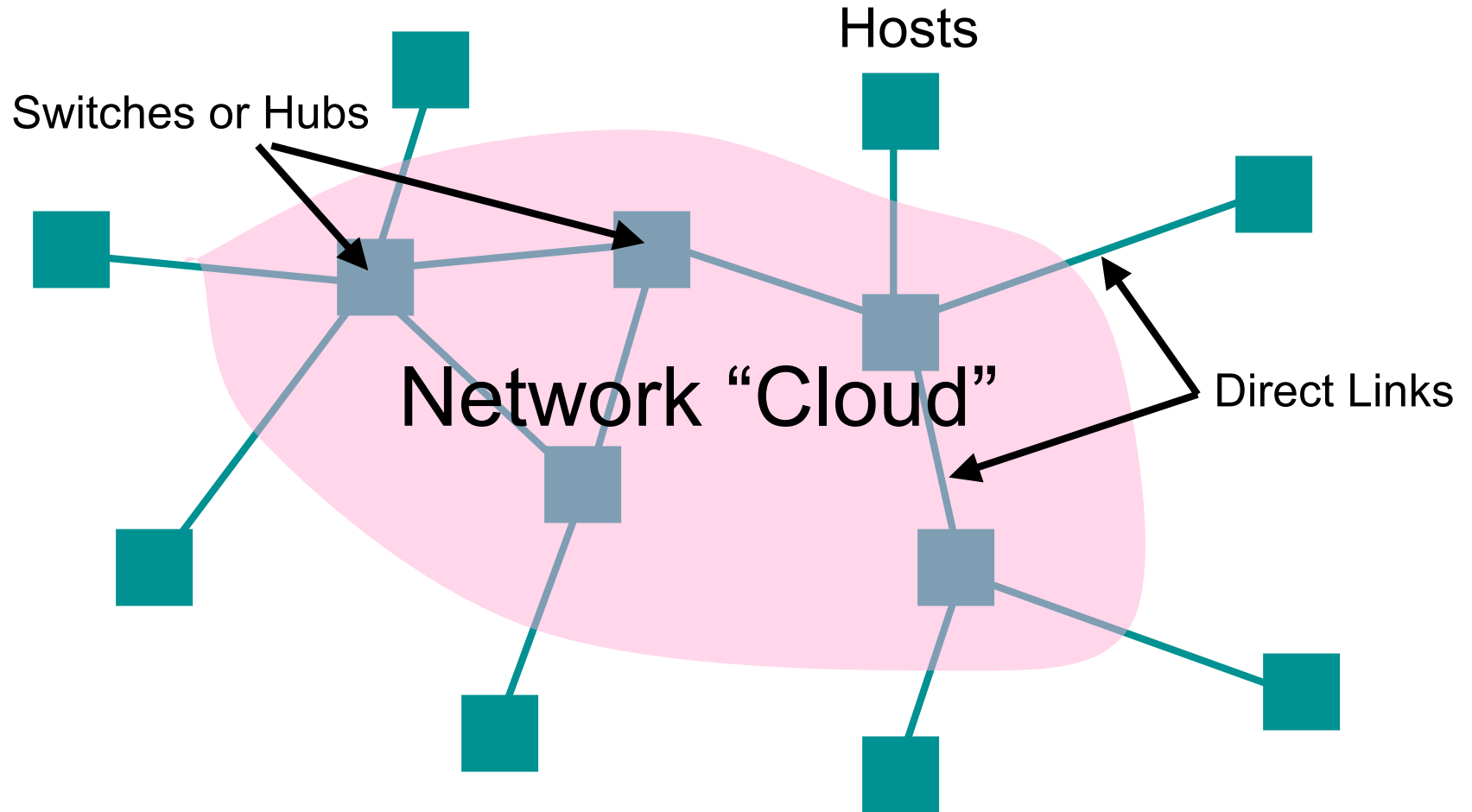


Point to Point  
e.g. telephone

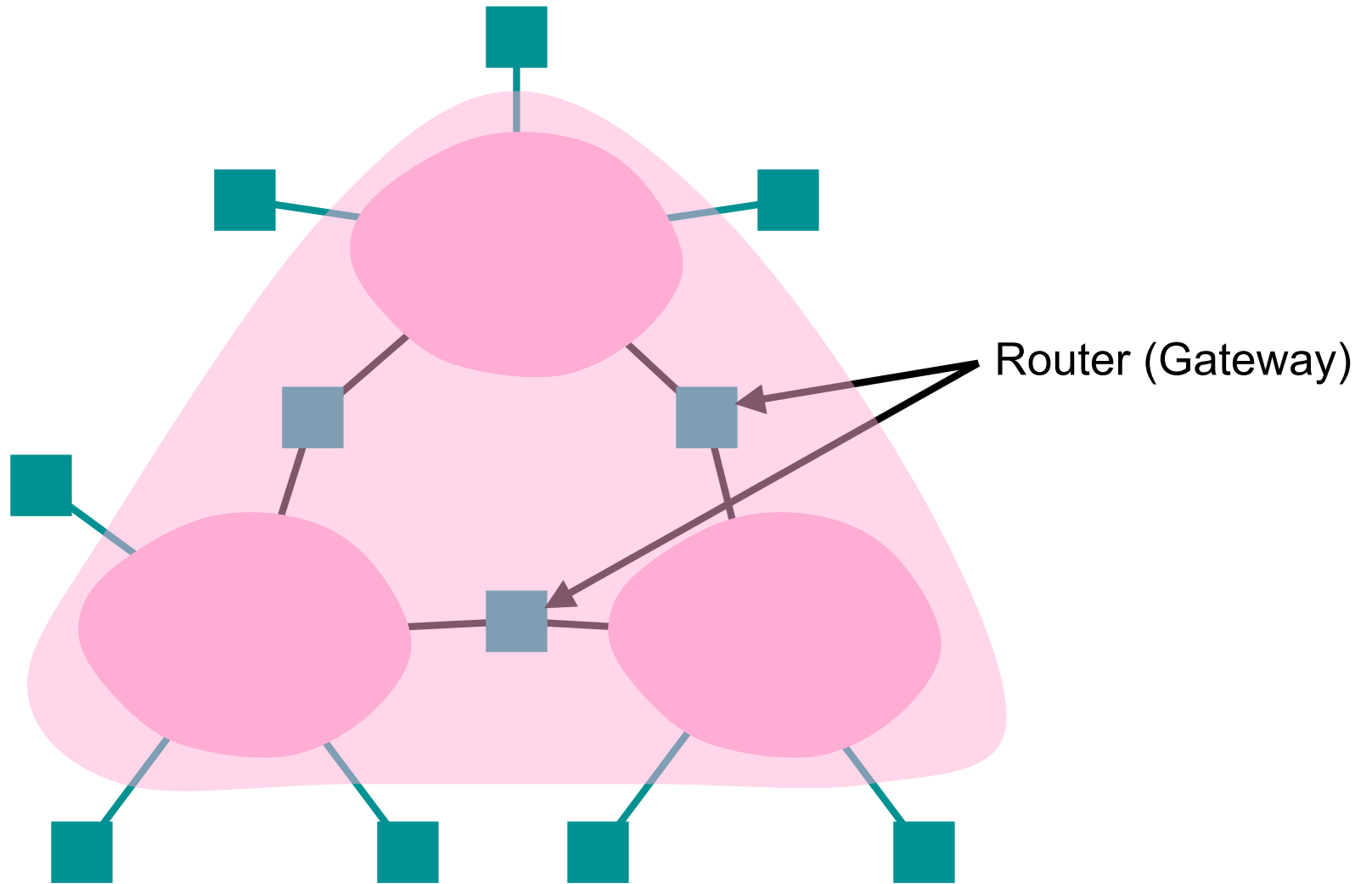


Multiple Access  
e.g. Ethernet

# Connectivity: Switched Networks

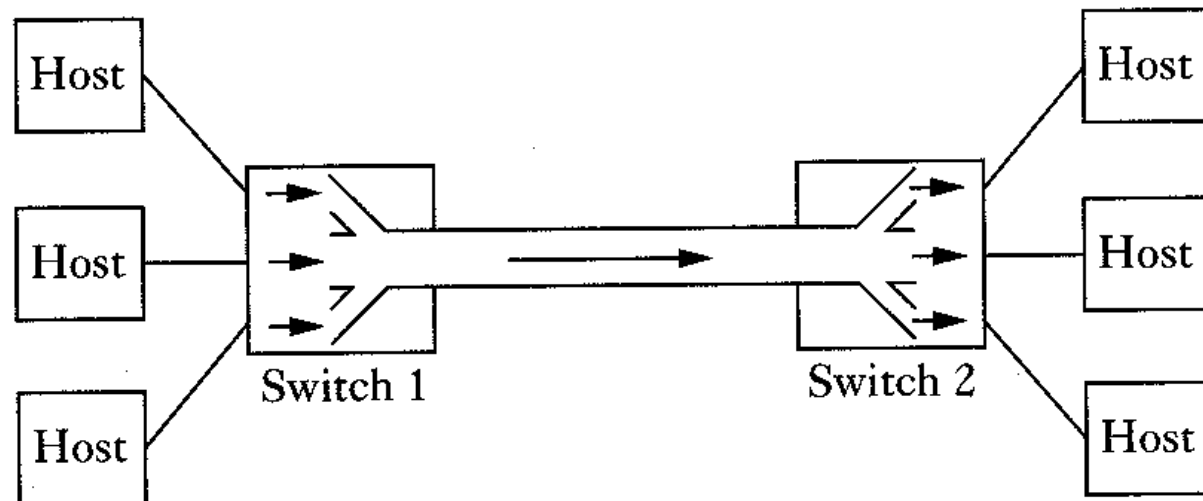


# Connectivity: Internetworks



# Resource Sharing: Multiplexing

- How can multiple hosts share the network if they want to use it at the same time?
  - Sharing links
  - Sharing switches





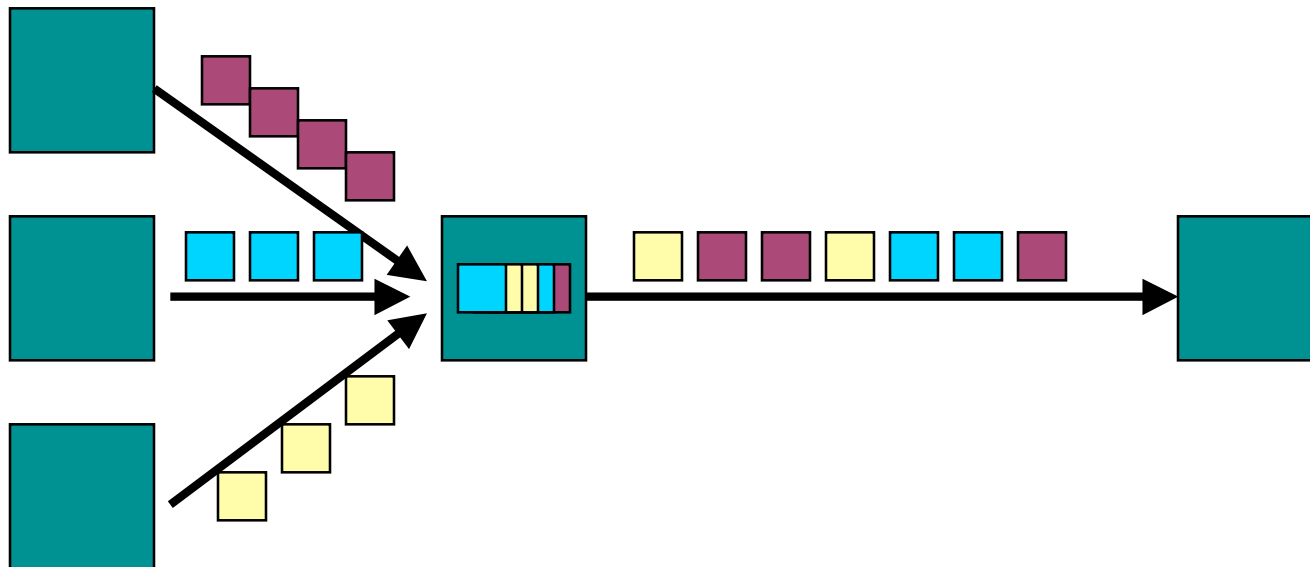
# Multiplexing: STDM & FDM

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- Synchronous Time-division Multiplexing (STDM)
  - “Time sharing”
  - Divide time into equal sized quanta
  - Round-robin
- Frequency-division Multiplexing (FDM)
  - Transmit all flows at different frequencies
  - Radio or Television
- Limitations:
  - Wasted resources
  - Maximum # flows can't be changed

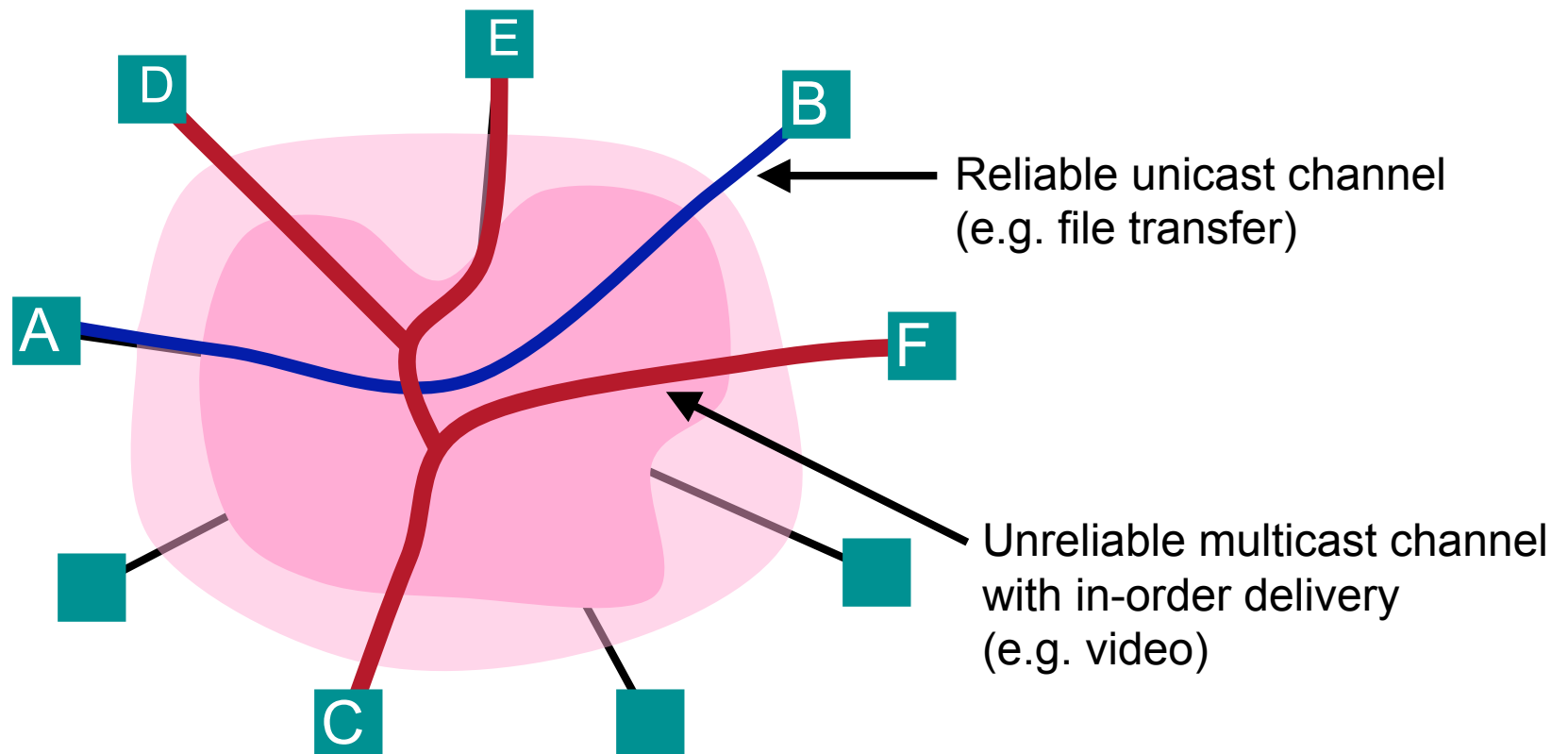
# Statistical Multiplexing

- Data is partitioned into *packets*
- Routing decision is made per packet
- Better resource usage than STDM
- Fairness? Congestion?



# Functionality

- Different applications require different services





# Functionality & Dealing with Failure

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- Fairness
- Congestion
- Quality of Service
- Bit or burst errors
- Link or node outages

# Performance

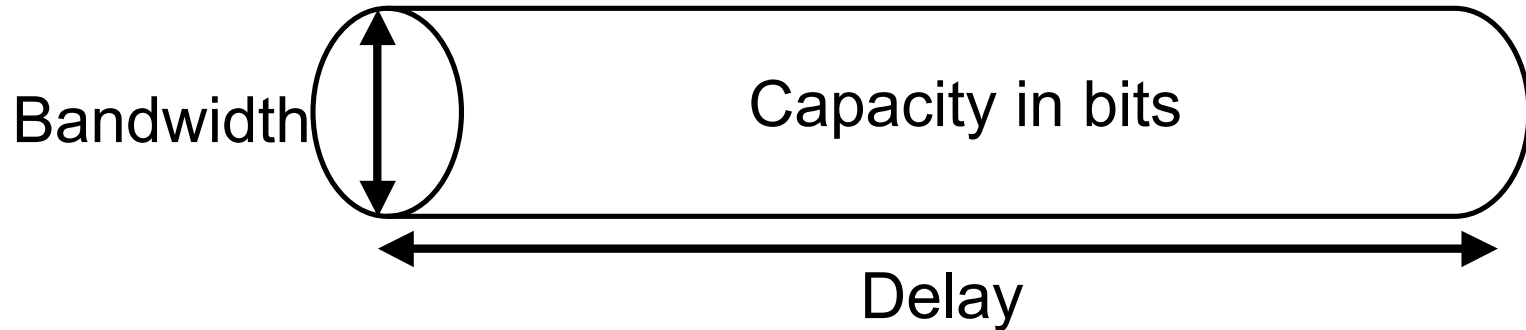
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- *Bandwidth* (throughput)
  - The number of bits that can be transmitted over the network in a certain period of time.
  - Measured in bits/sec
- *Latency* (delay)
  - How long it takes a single bit to propagate from one end of the network to the other.
  - Measured in seconds
- Round Trip Time (RTT)
  - How long it takes for a bit to get from one end of the network to the other *and back*.

# Connectivity: Direct Link Technologies

Wired Ethernet	10, 100 Mbps, 1, 10 Gbps
SONET fiber Synchronous Optical Network	up to 9.6 Gbps
CATV Cable TV	1-6 Mbps, asymmetric
ADSL Asymmetric Digital Subscriber Line	Downstream: 1.5-55.2 Mbps Upstream: 16-640 Kbps
ISDN Integrated Services Digital Network	64 Kbps*n with bonding
POTS Plain Old Telephone Service	56 Kbps
Wireless Ethernet	2, 11, 22, ... Mbps
Infrared IrDA	115 Kbps to 4 Mbps
CDPD Cellular Digital Packet Data	19.2 Kbps

# Performance: Delay x Bandwidth



Delay x Bandwidth determines the number of bits that can be “in flight”.

For efficient resource usage: keep the pipe full.

# Key Equations

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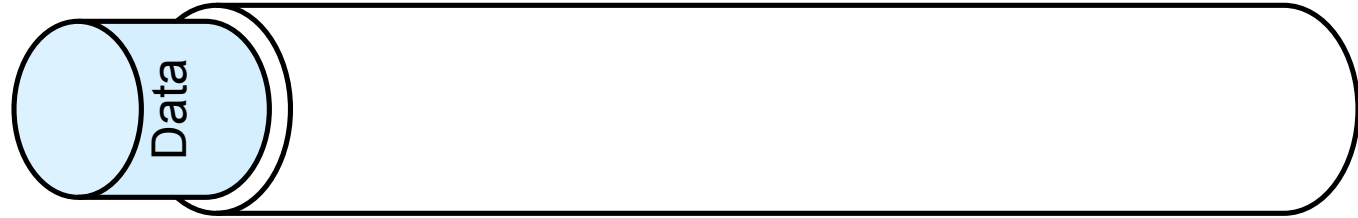
Latency = Propagation + Transmit + Queue

Propagation = Distance / SpeedOfLight

Transmit = Size / Bandwidth

# Total Latency: Direct Link

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Data moves through the link at the speed of light.

Time

0

Data ready to be sent

# Total Latency: Direct Link

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Data moves through the link at the speed of light.

Time

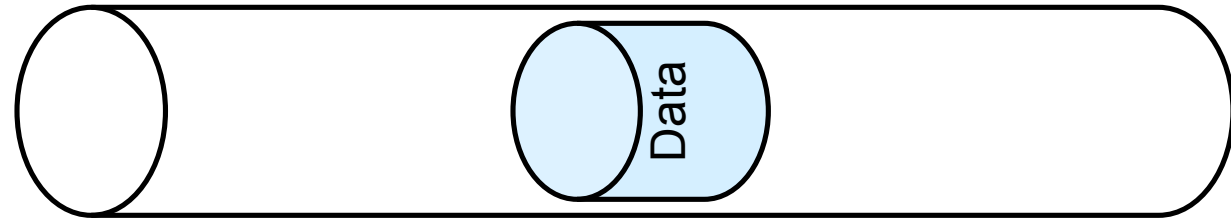
0

$t = \text{Size}/\text{Bandwidth}$

Data ready to be sent

Data in the link

# Total Latency: Direct Link



Data moves through the link at the speed of light.

Time

0

$t = \text{Size}/\text{Bandwidth}$

$t+k$

Data ready to be sent

Data in the link

Data traveling through the link

# Total Latency: Direct Link



Data moves through the link at the speed of light.

Time

0

Data ready to be sent

$t = \text{Size}/\text{Bandwidth}$

Data in the link

$t+k$

Data traveling through the link

$\text{prop} = \text{Distance}/\text{LightSpeed}$

First bit arrives at destination

# Total Latency: Direct Link



Data moves through the link at the speed of light.

Time

0

Data ready to be sent

$t = \text{Size}/\text{Bandwidth}$

Data in the link

$t+k$

Data traveling through the link

$\text{prop} = \text{Distance}/\text{LightSpeed}$

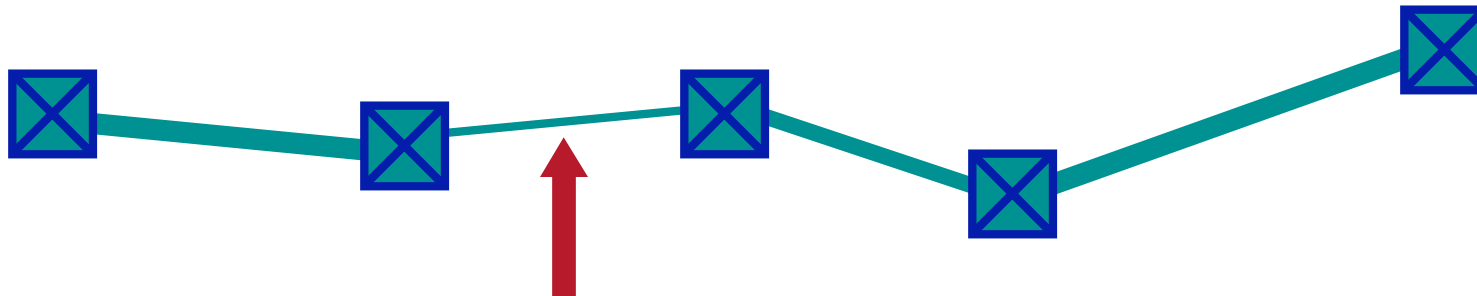
First bit arrives at destination

$\text{prop} + t$

Last bit arrives at destination

# Paths Are Made of *Links*

- Links are interconnected by zero or more *network elements*, e.g., switches, routers, hubs, bridges, etc.
- Path delay is sum of link delays plus queuing (switching) delays
- Path throughput = *bottleneck link* t'put



# Tradeoffs

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- RTT from Penn to Stanford is approx. 100ms
- 1.4 GHz workstation
  - 140 million cycles elapsed in that time
- Data compression
  - Trades machine cycles for bandwidth
- (Question: Why is RTT important?)

# Bandwidth vs. Latency

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- Which is the better deal:
  - Improve your *bandwidth* from 1 Mbps to 100 Mbps, or
  - Improve your *RTT* from 100 ms to 1 ms?
- The answer depends on what you need to send.

# Latency Bound

- Send 1 byte

Transmit Time	
1 Mbps	8 $\mu$ s
100 Mbps	.08 $\mu$ s

Perceived Latency	100 ms	1 ms	
1 Mbps	100.008 ms	1.008 ms	99%
100 Mbps	100.00008 ms	1.00008 ms	99%
	.008%	.8%	

# Bandwidth Bound

- Send 25 MB

Transmit Time	
1 Mbps	3.5 min
100 Mbps	21 sec

Perceived Latency	100 ms	1 ms	
1 Mbps	210.1 sec	210.001 sec	.05%
100 Mbps	21.1 sec	21.001 sec	.5%
	90%	90%	

# Some Units and Measurements

- Mbps =  $10^6$  bits/sec
- byte = 8 bits
- KB =  $2^{10}$  bytes (= 8,192 bits)
- MB =  $2^{20}$  bytes (= 8,388,608 bits)
- ms =  $10^{-3}$  seconds
- $\mu$ s =  $10^{-6}$  seconds
  
- Speed of light:
  - Vacuum :  $3 \times 10^8$  m/sec
  - Copper or Fiber:  $2 \times 10^8$  m/sec