

Unified Platform for Secure Networked Information System

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Motivation

- **Proliferation of new network architecture and protocols**
 - Overlay networks with new capabilities
 - Mobility, resiliency, anycast, multicast, anonymity, etc
 - Distributed data management applications
 - Network monitoring, publish-subscribe systems, content-distribution networks
- **Challenges - scalability and security threats**
- **Techniques proposed by security/networking community**
 - Distributed debugging: PIP [NSDI 06], FRIDAY [NSDI 07]
 - Forensics: IP traceback [SIGCOMM 00], IP forensics [ICNP 06]
 - Network accountability: PeerReview [SOSP 07], AIP [SIGCOMM 08]
 - Trust management: SD3 [Oakland 01], Delegation Logic [TISSEC 03]



Motivation

- **Problem: lacking generalized framework**

- Designed for specific security threats
- Implemented and enforced in different languages and environments
- Lack of cross-layer integration with existing distributed query processors

*A unified platform – **network protocol specification**, **security policy**, **support for a variety of techniques for secure networks***



Contributions

- **A unified declarative language:**
 - Declarative networking: network protocol specifications
 - Access control languages: logic for security policies
 - Securing network routing (S-BGP), DHTs, p2p query processing
- **Authenticated distributed query processing**
 - Extension of existing database techniques for *authenticated* communication
 - Implementation in a declarative networking engine
- ***Network provenance***
 - Data provenance: explain the existence of a tuple in database
 - Relate to real-world use cases in secure networked information systems
- **Experimental evaluation on a local cluster and Planetlab testbed**



Outline of Talk

- Introduction
- **Unified Declarative Framework**
 - Background: Declarative Networking and Access Control Languages
 - Secure Network Datalog (SeNDlog)
- Authenticated Distributed Query Processing
- Network Provenance
- Experimental Evaluation
- Conclusion & Future Work

Background: Declarative Networking

■ Declarative query language for network protocols

- Network Datalog (NDlog) – distributed Datalog [SIGCOMM 05, SIGMOD 06]
- Compiled to distributed dataflows, executed by distributed query engine
- *Location specifiers* (@ symbol) indicate the source/destination of messages

■ Example: Network Reachability

r1: `reachable(@S,D) :- link(@S,D)`

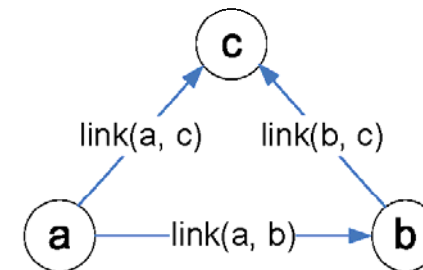
r2: `reachable(@S,D) :- link(@S,Z), reachable(@Z,D)`

`link(@a,b)` – “there is a link from node *a* to node *b*”

`reachable(@a,b)` – “node *a* can reach node *b*”

If there is a link from *S* to *D*, then *S* can reach *D*.

If there is a link from *S* to *Z*, AND *Z* can reach *D*, then *S* can reach *D*.



Node a	Node b
<code>link(@a, b)</code>	<code>link(@b, c)</code>
<code>link(@a, c)</code>	<code>reachable(@b, c)</code>
<code>reachable(@a, c)</code>	



Background: Access Control Languages

- **Access control is broadly defined as:**

- Enforce security policies in a multi-user environment
- Assigning credentials to principals to perform actions

- **Declarative interface:**

- Analyzing and implementing security policies
- Several runtime systems based on distributed Datalog/Prolog

- **Binder [Oakland 02]: a simple representative language**

- **Context:** each principal has its own context where its rules and data reside
- **Authentication:** “says” construct (credentials, signatures)

At alice:

b1: access(P,O,read) :- good(P).

b2: access(P,O,read) :- bob says access(P,O,read).

- “In alice's context, any principal P may access object O in read mode if P is good (b1) or, bob says P may do so (b2 - delegation)”

Secure Network Datalog (SeNDlog)

- Rules within a context
 - Untrusted network
 - Predicates in rule body in local context
- Authenticated communication
 - “says” construct
 - *Import predicate*: “X says p”
 - X asserts the predicate p.
 - *Export predicate*: “X says p@Y”
 - X exports the predicate p to Y.

```
r1: reachable(@S,D) :- link(@S,D).
r2: reachable(@Z,D) :- link(@S,Z),
    reachable(@Z,D).
```

⇓ *localization rewrite*

At S:

```
s1: reachable(@S,D) :- link(@S,D).
s2: linkD(D,S)@D :- link(S,D).
s3: reachable(Z,D)@Z :- linkD(@S,Z),
    reachable(@S,D).
```

⇓ *authenticated communication*

At S:

```
s1: reachable(@S,D) :- link(@S,D).
s2: S says linkD(D,S)@D :- link(S,D).
s3: S says reachable(Z,D)@Z :-
    Z says linkD(@S,Z),
    W says reachable(@S,D).
```




Example Protocols in SeNDlog

- **Secure network routing**

- Nodes import/export signed route advertisements from neighbors
- Advertisements include signed sub-paths (*authenticated provenance*)
- Building blocks for secure BGP

- **Distributed hash table overlay**

- Chord DHT – authenticate the node-join process
- Signed node identifiers to prevent malicious nodes from joining the DHT

- **P2P query processing – application layer**

- PIER - built upon Chord DHT
- Capability of *layered authentication*



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- **Authenticated Distributed Query Processing**
 - **Authenticated Pipeline Semi-Naïve**
 - **Dataflow Architecture**
- Network Provenance
- Experimental Evaluation
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Authenticated Query Processing

■ Semi-naïve Evaluation

- Standard technique for processing recursive queries
- Synchronous rounds of computation

■ Pipelined Semi-naïve Evaluation [SIGMOD 06]

- Asynchronous communication in distributed setting
- No requirement on expensive synchronous computation

■ Authenticated Semi-naïve Evaluation

- Modification for “says” construct, in p’s context:

$a :- d_1, \dots, d_n, b_1, \dots, b_m, p_1 \text{ says } a_1, p_2 \text{ says } a_2, \dots, p_o \text{ says } a_o.$

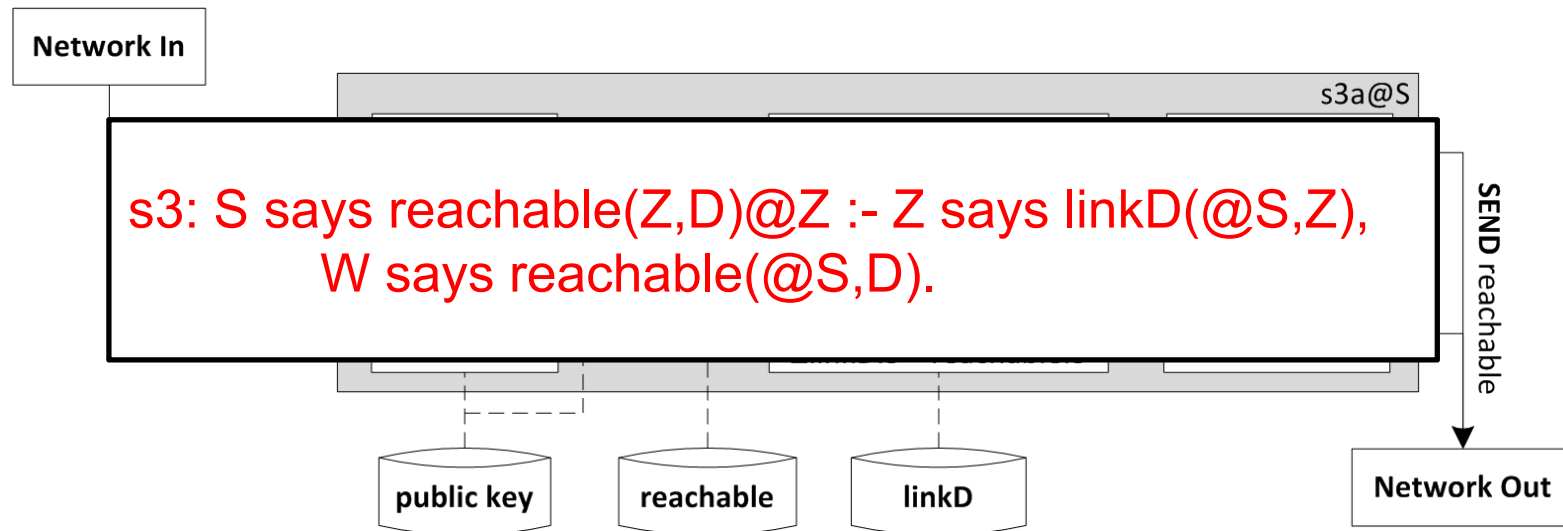
for k th *import predicate*, an authenticated delta rules is generated:

$p \text{ says } \Delta a :- d_1, \dots, d_n, b_1, \dots, b_m, p_1 \text{ says } a_1, \dots, p_k \text{ says } \Delta a_k, \dots, p_o \text{ says } a_o.$

Architectural Overview of Dataflow

■ Dataflow Architecture

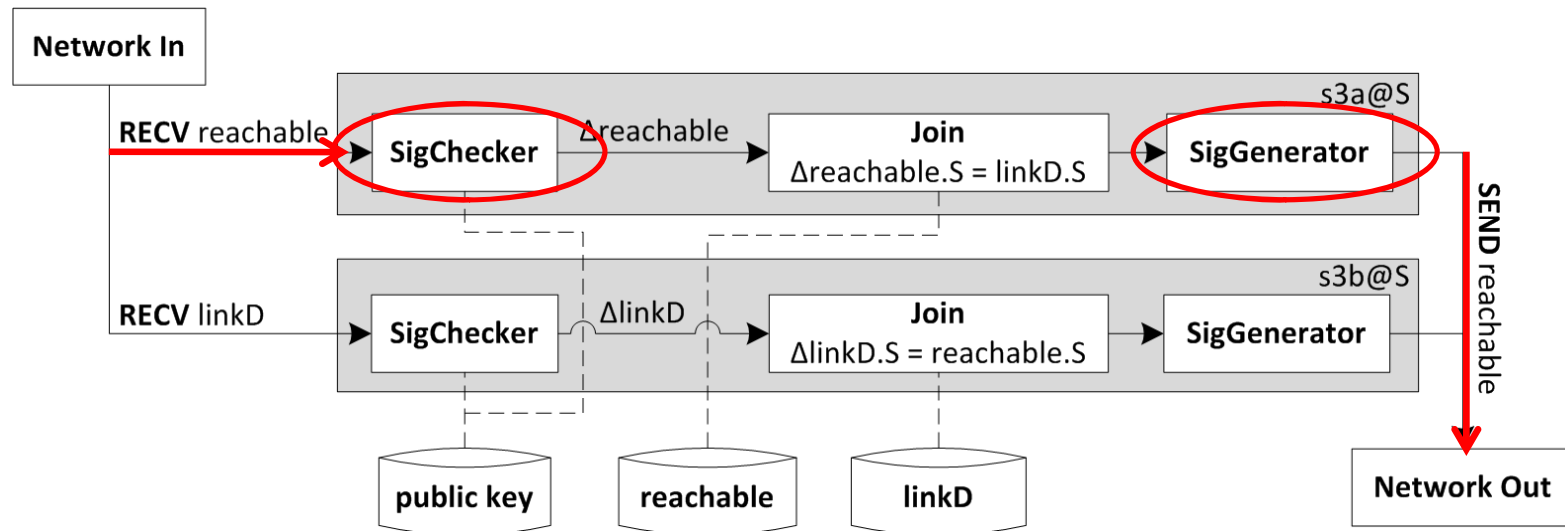
- Based on the P2 declarative networking system [<http://p2.cs.berkeley.edu/>]
- Additional modules to support authenticated communication



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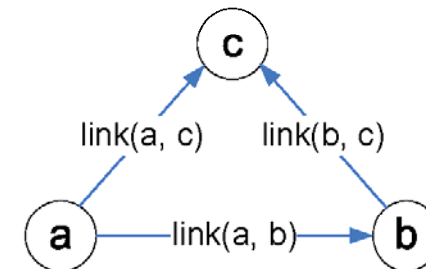
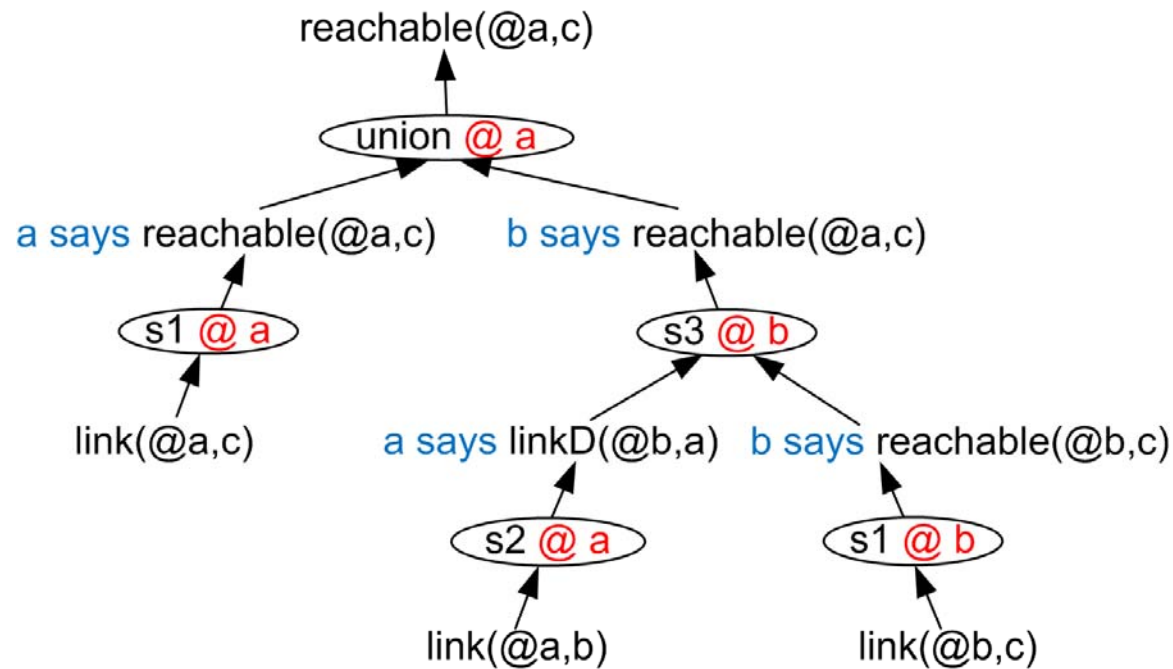


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- Authenticated Distributed Query Processing
- **Network Provenance**
 - Network Provenance
 - Wide Application of Network Provenance
- Experimental Evaluation
- Conclusion & Future Work

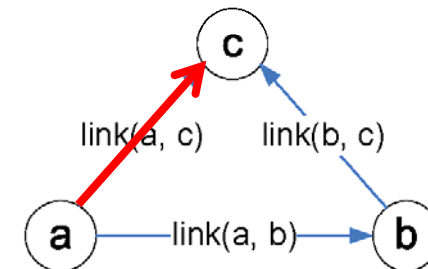
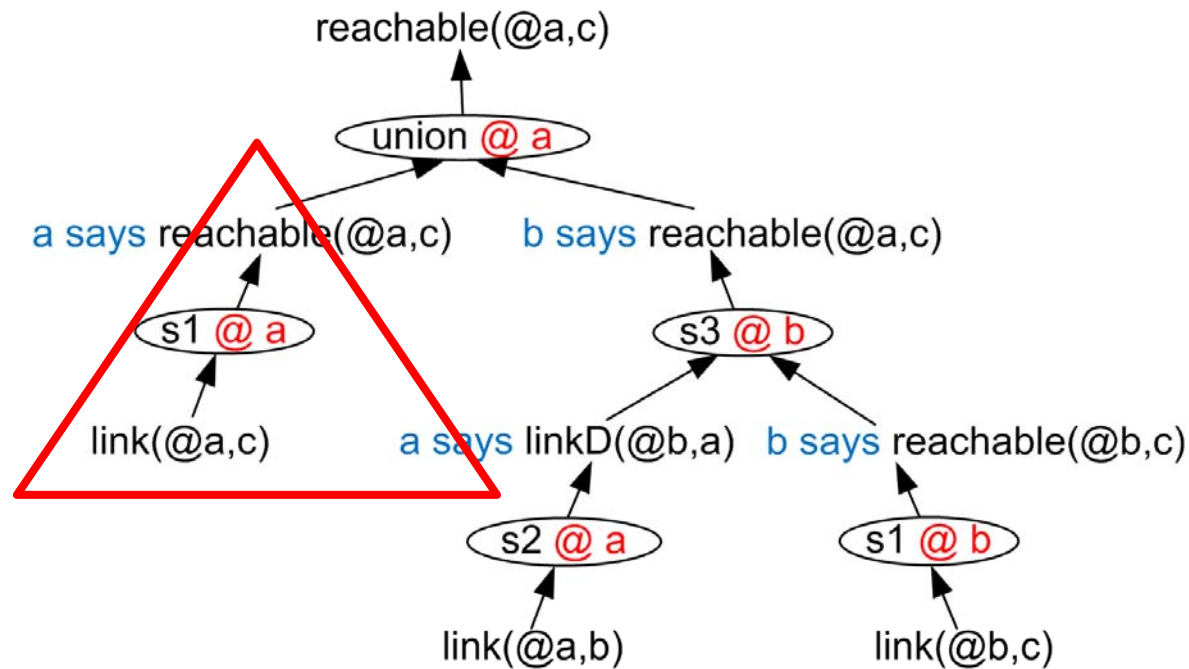
Network Provenance

- Naturally captured within declarative framework
- Explain the existence of any network state
- Similar notion in security community: *proof-trees*



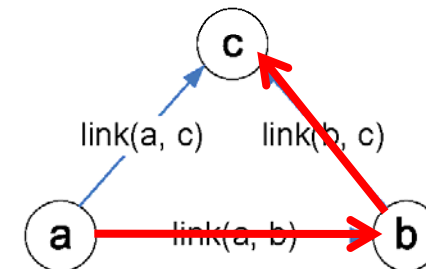
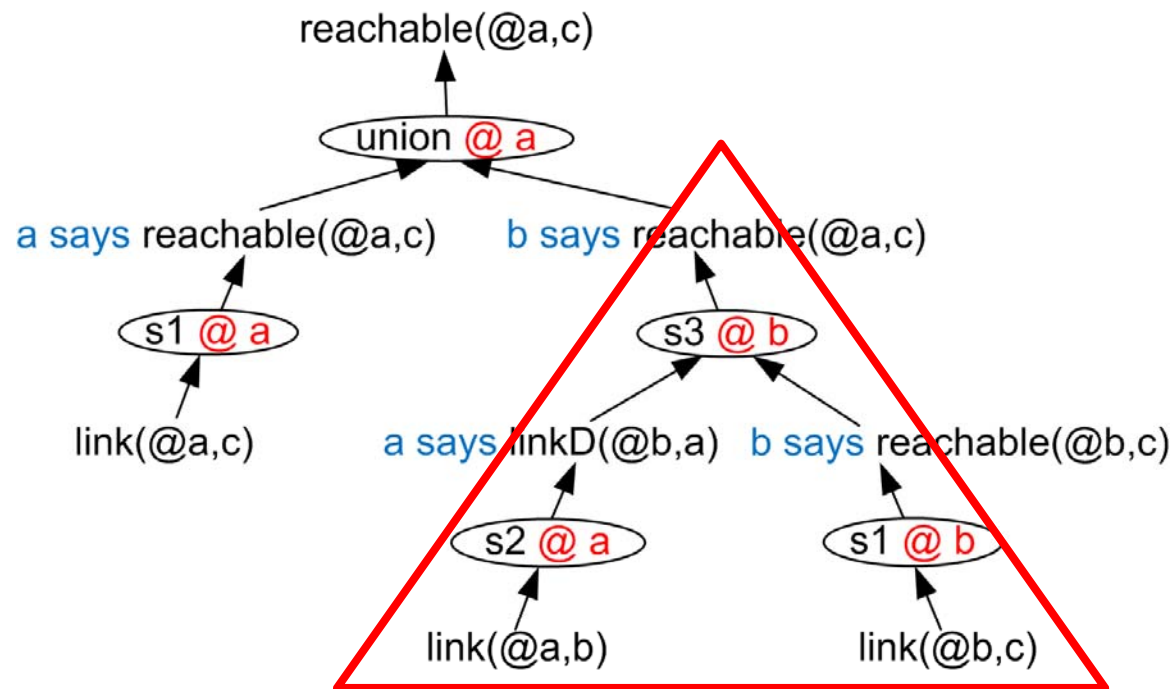
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Wide Application of Network Provenance

Provenance Taxonomy	Distributed Debugging	Forensics	Network Accountability	Trust Management
Derivation Tree / Algebra Expr.	Both	Derivation Tree	Both	Algebra Expr.
Local / Distributed	Both	Both	Both	Local
Online / Offline	Online	Offline	Offline	Online
Boolean/ Quantifiable	Both	Boolean	Boolean	Both

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Experimental Setup

- **P2 declarative networking system**

- Extensions for security and provenance support

- **Workload**

- Path-vector – network routing
- Chord – distributed hash table
- PIER – p2p query processing

- **Test-bed**

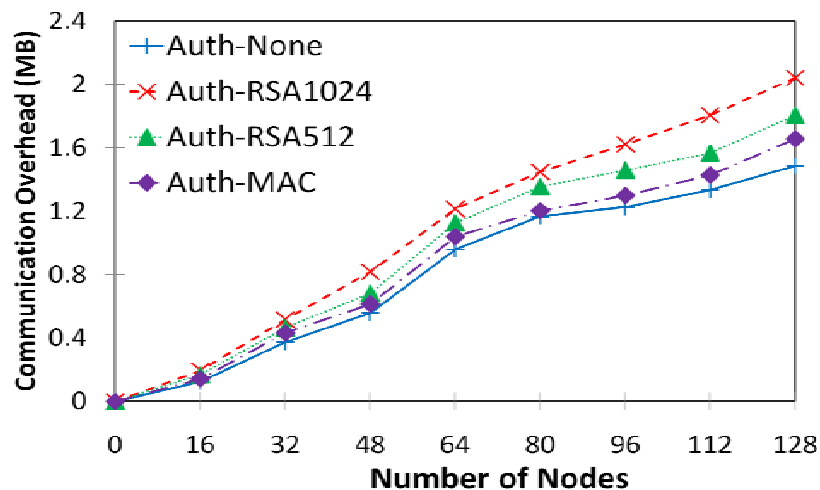
- A local cluster with 16 quad-core machines
- Planetlab testbed with 80 nodes

- **Metrics**

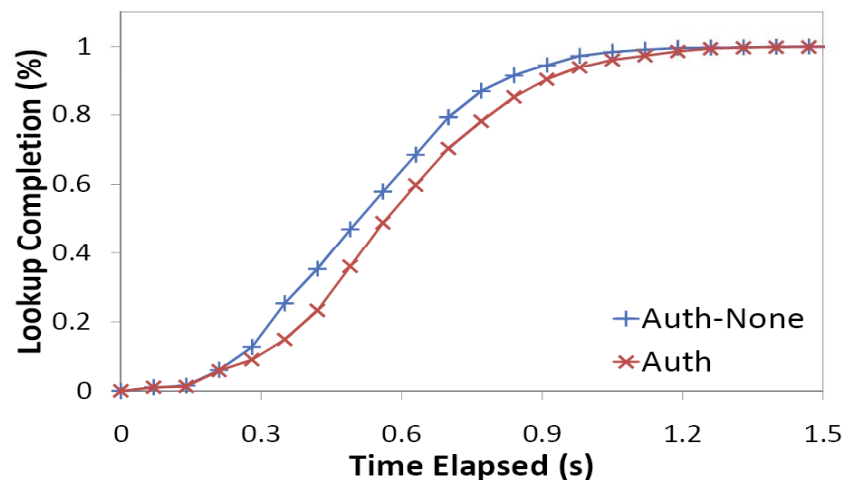
- Communication overhead
- Query completion time / lookup latency

Feasibility Study of SeNDlog

Path-vector on cluster



Lookup latency for Chord on cluster



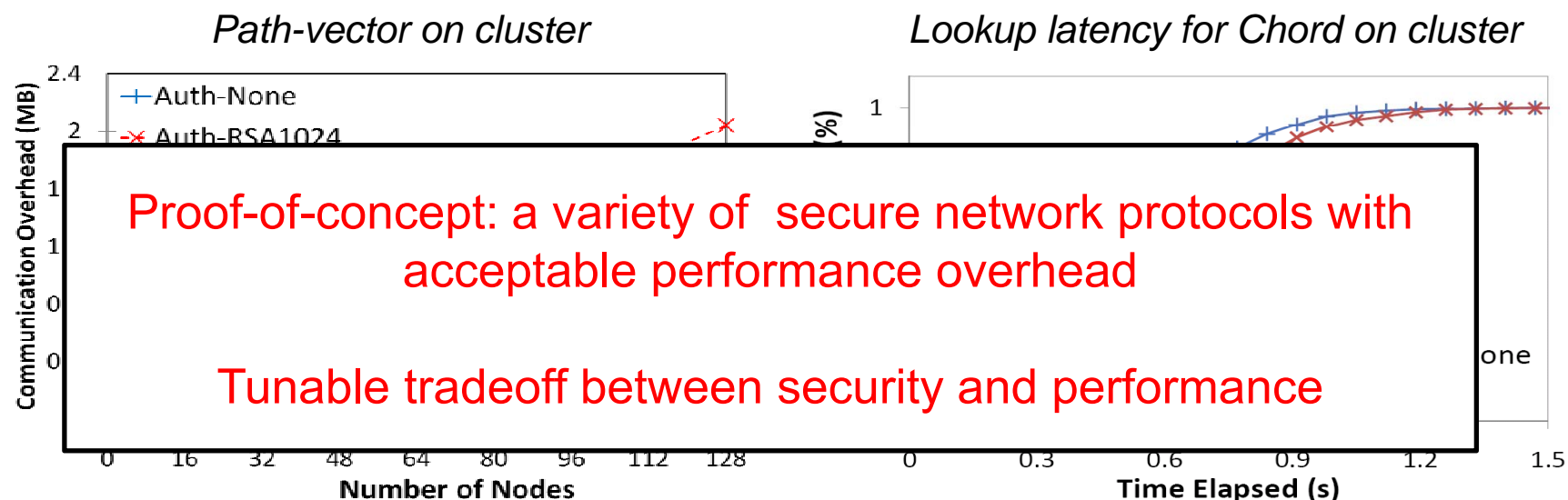
□ Path-vector protocol

- 128 nodes, 6 neighbors per node
- Auth-HMAC – 10% increase
- Auth-RSA512 – 20% increase
- Auth-RSA1024 – 40% increase

□ Chord DHT protocol

- 128 Chord nodes, random lookups
- Auth (with RSA1024) – less than 10% increase to finish 50% lookups

Feasibility Study of SeNDlog



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Other Evaluation Results

- **Planetlab Experiments**

- WAN effect: high inter-node latency, low availability of computation resources
- Relative overhead increase is amortized

- **Validate the feasibility of network provenance**

- Packet delivery: routing tables pre-computed using SeNDLog programs
- Use local network provenance to trace the route taken by a packet
- Acceptable performance: 0.05s increase in packet delivery latency in LAN



Conclusion & Future Work

■ Conclusion

- SeNDlog: Unified language for declarative networking and access control
- Authenticated query processing techniques for distributed settings
- Support for network provenance

■ Future Work

- Possible language extensions
 - Secrecy / encrypted facts
 - Restricted delegation / “speaks-for” primitive
- Optimizations opportunities
 - Performance / security tradeoff
 - Bandwidth optimization for network provenance



Thank You ...