Hierarchical Hybrid Modeling of Embedded Systems

Rajeev Alur

Systems Design Research Lab University of Pennsylvania www.cis.upenn.edu/~alur/

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Programming Interacting Autonomous Robots

Low level

Analysis of vision data Control laws for legs Wireless cards

Current programming

How to implement Go-to-ball Real-time scheduling

High level

Modes: Attack, Defend How to switch? Strategies Communication to collaborate



Trends in Model-Based Design

- □ Emerging notations: UML-RT, Stateflow
 - Visual, Hierarchical, Object oriented
 - Simulation, code generation
- □ Formal models and Model checking tools
- Programming languages (Esterel, FRP...)
- □ Control engineer's tools (Matlab...)
- **Design/Simulation environments**
 - SHIFT, Ptolemy-II, Modelica

Guiding Themes for CHARON

- Integrated modeling of control program and physical environment (hybrid)
 - Programming language technology in Control tools
 - Continuous modeling in Programming environments
- Foundations for hybrid systems in presence of concurrency, hierarchy, exceptions ...
 - Compositionality, refinement,
- □ Models need to be analyzable
 - Model checking
 - Exploit modeling constructs for efficiency

CHARON Team

Rajeev Alur (CLS) Vijay Kumar (MEAM) I nsup Lee (CLS) George Pappas (EE)	Joel Esposito Yerang Hur Franjo I vancic P K Mishra Usa Sammapun
Research Associates Thao Dang Salvatore La Torre Supratik Mukhopadhyay	Collaborators Rafael Fierro (U Oklahoma) Radu Grosu (SUNY StonyBrook)

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Talk Outline

- ✓ Motivation
- **CHARON Summary**
- **Compositional Refinement**
- Model Checking via Predicate Abstraction
- **Conclusions**

Hybrid Modeling

State machines + Dynamical systems



Behavioral Hierarchy



CHARON Language Features

- □ Individual components described as agents
 - Composition, instantiation, and hiding
- □ Individual behaviors described as modes
 - Encapsulation, instantiation, and Scoping
- □ Support for concurrency
 - Shared variables as well as message passing
- □ Support for discrete and continuous behavior
 - Differential as well as algebraic constraints
 - Discrete transitions can call Java routines

Syntax: Modes and Agents



Modes describe sequential behavior
Agents describe concurrency

Example: V2V model

Agent Vehicle



Agent VehiclePlant



Agent RegulationLayer



CHARON Toolkit

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What is Compositionality ?



Which properties are preserved?

Can we restrict reasoning to modified parts of design?

Mode should have a precise interface spec that would permit composition of behaviors

Mode Executions



(ctl,t,level,infusion,rate,h)

(dx,0,5.1,1,0.2,Maintain)			
	Flow Step		
(dx,10,15.1,3,0.2,Maintain)			
	Env Step		
(de, 10, 15.1, 5, 0.2, Maintain)			
	Discrete Mode Step		
(dx,10,15.1,	5,0.1,Compute)		

Semantics of modes

□ Semantics of a mode consists of:

- control interface: entry and exit points
- data interface: global variables
- traces (sequences over observable states)
- □ Key Thm: Semantics is compositional
 - traces of a mode can be computed from traces of its sub-modes

Refinement

Refinement is trace inclusion



- Same control points and global variables
- Guards and constraints are relaxed

Sub-mode refinement



Compositional Reasoning



Sub-mode refinement

Context refinement

Automated refinement checking for discrete systems Pipelined processors, Network protocols

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Model Checking of Hybrid Systems

Goal: Given an initial region, compute whether a bad state can be reached

- Existing tools: HyTech, d/dt, Checkmate
- Key step is to compute Reach(X) for a given set X under dx = Ax+Bu (expensive !!!)



Polyhedral Flow Pipe Approximations



A. Chutinan and B. H. Krogh, Computing polyhedral approximations to dynamic flow pipes, IEEE CDC, 1998

Predicate Abstraction

- I nput is a hybrid automaton and a set of k boolean predicates, e.g. x+y > 5-z.
- □ The partitioning of the concrete state space is specified by the user-defined *k* predicates.



Overview of the Approach



Why use this approach?

- Reach(X) needs to be computed only for abstract states X and not intermediate regions of unpredictable shapes/complexity
- No need to compute Reach (X). Goal is to find one new abstract state reachable from X, partial results are of great use
 - Simulate vertices
 - Consider time-slices at discrete times
- Our focus is on search strategies to make progress in the abstract state-space
- □ Initial implementation in C++ with promising results

Case-study: V2V

- Platoon controller for 3 vehicles scenario
- First step: make the model linear (feedback linearization)
- □ Initial predicates from model description. E.g.
 - from LowLevelController:
 - $u_{isl} < 0, u_{isl} > 0$
 - from BrakeControl:
 - u_isl < max_brake_pressure</pre>
 - from ThrottleControl:
 - u_isl < max_throttle</pre>
- Safety of the controller verified using the tool (17 predicates, 4 continuous vars, < 1 min)</p>

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Ongoing Activities

- Distributed simulation
- Qualitative abstractions of hybrid systems
- Model-based test generation
- Accurate event detection for simulation
- Exploiting hierarchy for efficient simulation
- □ Applications/Case-studies
 - Multiple autonomous robots
 - MoBIES challenge problems
 - Animation
 - Biomolecular networks...

Wrap-Up

- □ Modeling and Analysis in symbiosis
- □ Common themes in many modeling proposals
 - Hierarchy
 - Concurrency and Communication
 - Component interfaces
 - Formal semantics and compositionality
 - Integration: Discrete + Continuous + Stochastic
- Automating formal verification is hard, but not impossible, and there is steady progress
- Biological systems: emerging application for modeling with similarities to embedded software