Symbolic-Numeric Algorithms for Constraint Solving:

Overview, Challenges, Applications, Future

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NSF Workshop on Symbolic Computation for Constraint Satisfaction Problems, 14 November 2008 – p. 1/30

Outline of the presentation

- Overview of Symbolic-Numeric Algorithms
- Challenges
- Applications
- Future directions?

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Context: Continuous Constraints

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 - Same as discrete ones, except for the domains of their variables: continuous

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 - Same as discrete ones, except for the domains of their variables: continuous
- **•** Example:
 - *a chemistry problem:*

 $14 * z1^{2} + 6 * z1 * z2 + 5 * z1 - 72 * z2^{2} - 18 * z2 = 850 * z3 - 2.0e - 9,$ $0.5 * z1 * z2^{2} + 0.01 * z1 * z2 + 0.13 * z2^{2} + 0.04 * z2 = 4.0e4,$ 0.03 * z1 * z3 + 0.04 * z3 = 850

 $z_i \in [-1000; +1000]$

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- What do we do with these intervals during the search to find solutions?
 - Intervals are contracted (possibly discarded) along the search

Symbolic-Numeric Algorithms

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Symbolic-Numeric Algorithms

- Numeric part of the algorithms = the interval computations to contract the domains
- Symbolic part of the algorithms = the design of these contractors, and more

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Historically: the <u>Newton</u> method

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- Limitations:
 - What if $0 \in f'(D)$? \rightsquigarrow no contraction
- Reasons for this:
 - D is so large that f is not monotonic on D
 - f'(D) is overestimated, hence including a 0 while it is not part of the range

Let us assume we have a linear constraint to solve.

Let our constraint be:

$$a_1x_1 + a_2x_2 + \ldots + a_nx_n = b$$

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Natural approach: the Gauss-Seidel method:

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Limitation: What if the constraint is not linear?

Let us assume we have a nonlinear constraint to solve.

Let our constraint be:

$$\cos(x) + 3ztan(y) = 2 \iff x = \cos^{-1}(2 - 3ztan(y))$$

= symbolic inversion of (nonlinear) constraints

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- Different approaches to such inversion:
 - Kearfott 1991: decomposition of the constraints into primitive, easily invertible constraints
 - Ceberio & Granvilliers 2000: recursive inversion of terms
 - Hansen & Walster 2003: inversion of f as f = g h where g is easy to invert

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$$x^2 + y^2 - 1 = 0$$
 with $x, y \in [0, 2]$
 $y = (1 - x^2)^{1/2} = (1 - [0, 4])^{1/2} = [-1, 1]$

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- Problem with this approach:
 - Result = intersection of approximations
 - What we aim at = approximation of intersection

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Challenges

- Weaknesses of interval computations: Overestimation leads to slower solving processes
- Weaknesses of the domain contractions (\equiv constraint consistency techniques):

The locality of reasonings leads to slower solving processes

Efficiency of interval computations:

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 - Decomposition of intervals into smaller ones that are evaluated
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 - <u>But:</u> expensive

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$$C: \begin{cases} c_1: & x+y+x^2+xy+y^2 &= 0\\ c_2: & x+t+xy+t^2+x^2 &= 0\\ c_3: & y+z+x^2+z^2 &= 0\\ c_4: & x+z+x^2+y^2+z^2+xy &= 0 \end{cases}$$

defined over $E = [-100, 100]^4$,

4 solutions reached in 140ms using realpaver [Granvilliers, 2002].

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<pre>{</pre>	$lc_1:$ $lc_2:$ $lc_3:$	$egin{array}{c} x \ x \end{array}$	+y	+z	+t	$+u_1$ $+u_1$ $+u_1$	$+u_2$ $+u_2$	+	u_3	$+u_4$	$+u_5$	= 0 = 0 = 0
l	$lc_4:$	x	U	+z		$+u_1$	$+u_{2}$	+	u_3		$+u_{5}$	= 0
							ſ	u_1	=	x^2		
								u_2	=	xy		
			and th	ne abs [.]	tracte	d syste	m: {	u_3	=	y^2		
								u_4	=	t^2		
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	lc'_3 :			-z	$-u_5$			+x	$+u_2$	$+ u_{3}$	= 0
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The new system is solved in 240ms!!

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Let us consider again the same problem. We begin with the linearized system:

ſ	$lc_1:$	x	+y			$+u_1$	$+u_{2}$	$+u_{3}$			= 0
J	$lc_2:$	x			+t	$+u_1$	$+u_{2}$		$+u_4$		= 0
	$lc_3:$		y	+z		$+u_1$				$+u_5$	= 0
l	$lc_4:$	x		+z		$+u_1$	$+u_2$	$+u_3$		$+u_{5}$	= 0

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First step of elimination

ſ	lc_3 :		y	+z		$+u_1$				$+u_5$	=	0
J	$lc_1:$	x	+y			$+u_1$	$+u_2$	$+u_3$			=	0
	$lc_2:$	x			+t	$+u_1$	$+u_{2}$		$+u_4$		=	0
	$lc'_4:$	-x	+y				$-u_2$	$-u_3$			=	0

Control criterion: controls the densification of the "linear" system

User linear part: 0

Abstracted linear part: -2

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Second step of elimination

ſ	$lc_3:$		y	+z		$+u_1$				$+u_5$	=	0
J	lc_4' :	-x	+y				$-u_2$	$-u_3$			=	0
	$lc'_1:$		2y			$+u_1$					=	0
l	lc_2 :	x			+t	$+u_1$	$+u_2$		$+u_4$		=	0

Control criterion: controls the densification of the "linear" system

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Abstracted linear part: -1

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Third step of elimination

ſ	$lc_3:$		y	+z	$+u_1$				$+u_5$	—	0
J	$lc'_4:$	-x	+y			$-u_2$	$-u_3$			=	0
	$lc_1':$		2y		$+u_1$					=	0
	$lc_2':$	-x	+2y	-t		$-u_2$		$-u_4$		=	0

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User linear part: +1

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Triangularized system

$$\begin{cases} lc'_{2}: & -t & -x & -u_{2} & -u_{4} & +2y & = & 0\\ lc'_{4}: & -x & -u_{2} & -u_{3} & +y & = & 0\\ lc_{3}: & +u_{5} & +z & +u_{1} & +y & = & 0\\ lc'_{1}: & +u_{1} & +2y & = & 0 \end{cases}$$

Concretization: nonlinear terms are restored

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Concretization phase

$$\begin{cases} c'_1: & -t & -x & -xy & -t^2 & +2y & = & 0 \\ c'_2: & -x & -xy & -y^2 & +y & = & 0 \\ c'_3: & z^2 & +z & +x^2 & +y & = & 0 \\ c'_4: & x^2 & +2y & = & 0 \end{cases}$$

Post-processing: *simplification of the system using specific constraints*

$$x_i = f(x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_n)$$

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Post-processing:
$$x^2 = -2y$$

	$c_{1}^{T}:$	-t	-x	-xy		$-t^{2}$				+2y	=	0
C_{π}	c_2^T :		-x	-xy	$-y^2$					+y	=	0
$\mathcal{O}T$.	$c_3^{T'}$:						z^2	+z		-y	=	0
	c_4^T :								$+x^{2}$	+2y	=	0

Solving stage: 4 solutions reached in 10ms!

- Redundant constraints: "the more the better"
- Linearization: as the solving process goes, eliminate non-linear terms
- **Triangularization:** *Gaussian-like elimination*
- Common Sub-Expressions technique

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- Chemistry
- Aeronautics
- Software verification
- Bio-medical engineering
- Geosciences
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- **2.** Solve $S \land \neg P$
- 3. If there is no solution: the program is verified
- 4. Otherwise we have test cases for counter-examples

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 - 2. It can be translated into a pattern constraint problem: qualitative approach

Geosciences

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- Limitations...
 - The odds are that there will be no "satisfactory" solution
 - *•* The constraints are too strong: modeling a "too secure" network

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 - Keep the validity constraints: they model the MLS policy
 - S Relax the risk constraints: aim at the minimum risk
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Future directions of work? (1)

Existing challenges:

- Extension of global constraints
- Combining the current work on DAGs with triangularization approaches
- "New" kinds of interval arithmetic e.g., circular arithmetic of Siegfried Rump

Future directions of work? (2)

New challenges?

- Robustness of solutions
- Use of numeric tensors, n-d arrays to better represent the dependence between variables' "values"
 (discussions from NSF-sponsored workshop, CoProD'08)
 Similar to shaving to some extent (forward looking)
 Can bring improvement thanks to the availability of information (n-d)

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 - The combination occurs at different stages: e.g., pre-processing, completely integrated
- Many challenges still to be addressed, including:
 - Pursuing existing research directions
 - *Exploring new representations: such as tensors, circular arithmetic*

The end

Thank you for your attention ANY QUESTIONS?

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- Enumeration is not an option...
- Algorithms based on intervals

Intervals are enumerated and the whole search space is covered

* Branch and Bound (B&B):

http://www-sop.inria.fr/coprin/logiciels/ALIAS/Movie/movie_undergraduate.mpg

***** More sophisticated consistency algorithms:

Box / Hull-consistencies and their combinations Result in Branch and Prune algorithms (B&P)

How is Branch and Prune different?

- Discarding boxes in no longer based on the constraints' evaluations only
- Consistency techniques are used to filter out / prune elements that do not satisfy the constraints

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Then, we have to bisect if we want more information but the domain to bisect is much smaller (25%) than with B& B

Example from chemistry:

 $14 * z_1^2 + 6 * z_1 * z_2 + 5 * z_1 - 72 * z_2^2 - 18 * z_2 = 850 * z_3 - 2.0e - 9,$ $0.5 * z_1 * z_2^2 + 0.01 * z_1 * z_2 + 0.13 * z_2^2 + 0.04 * z_2 = 4.0e4,$ $0.03 * z_1 * z_3 + 0.04 * z_3 = 850$

 $z_i \in [-1000; +1000]$

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Solutions:

OUTER BOX 1 $z_1 \in [124.7643488370932, 124.7643488370934]$ $z_2 \in [25.28546066708894, 25.28546066708898]$ $z_3 \in [224.6935300130224, 224.6935300130227]$ OUTER BOX 2 $z_1 \in [131.7475644308769, 131.7475644403561]$ $z_2 \in [-24.62787829918984, -24.62787829860991]$ $z_3 \in [212.9030823228049, 212.9030823233716]$

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- Duplicates... This is a real problem

Most of them are local...

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Some results

Problem	v	Initial Pb.		Triangul. Pb.	
		Time	Sol.	Time	Sol.
Bratu	7	1.10	3	0.60	4
	8	0.70	2	0.10	2
	10	2.30	2	0.10	2
	13	20.50	6	0.10	2
	14	46.40	11	0.20	2
	15	94.40	12	0.20	2

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+ Recent work of Gilles Trombettoni et al. on combining these ideas within the consistency techniques.