Model-driven Test Generation

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Outline and scope

- Classification of model-driven testing
- Conformance testing for communication protocols
- Coverage-based testing
  - Coverage criteria
  - Coverage-based test generation
- Can we do more? (open questions)

Testing classification

- By component level
  - Unit testing
  - Integration testing
  - System testing
- By abstraction level
  - Black box
  - White box
  - Grey box ???

Testing classification

- By purpose
  - Functional testing
  - Performance testing
  - Robustness testing
  - Stress testing
- Who performs testing?
  - Developers
  - In-house QA
  - Third-party

Functional testing

- An implementation can exhibit a variety of behaviors

- For each behavior, we can tell whether it is correct or not

- A *test* can be applied to the implementation and accept or reject one or more behaviors
  - The test fails if a behavior is rejected

- A *test suite* is a finite collection of tests
  - Testing fails if any test in the suite fails

Formal methods in testing

- "Testing can never demonstrate the absence of errors, only their presence."
  Edsger W. Dijkstra

- How can formal methods help?

- Add rigor!
  - Reliably identify what should to be tested
  - Provide basis for test generation
  - Provide basis for test execution

Model-driven testing

- Rely on a model of the system
  - Different interpretations of a model
- Model is a requirement
  - Black-box conformance testing
  - QA or third party
- Model is a design artifact
  - Grey-box unit/system testing
  - QA or developers

Conformance testing

- A specification prescribes legal behaviors
- Does the implementation conform to the specification?
  - Need the notion of conformance
- Not interested in:
  - How the system is implemented?
  - What went wrong if an error is found?
  - What else the system can do?

Test hypothesis

- How do we relate beasts of different species?
  - Implementation is a physical object
  - Specification is a formal object
- Assume there is a formal model that is faithful to implementation
  - We do not know it!
- Define conformance between the model and the specification
  - Generate tests to demonstrate conformance

Conformance testing with LTS

- Requirement is specified as a labeled transition system

- Implementation is modeled as an input-output transition system

- Conformance relation is given by `ioco`
  - [Tretmans96]
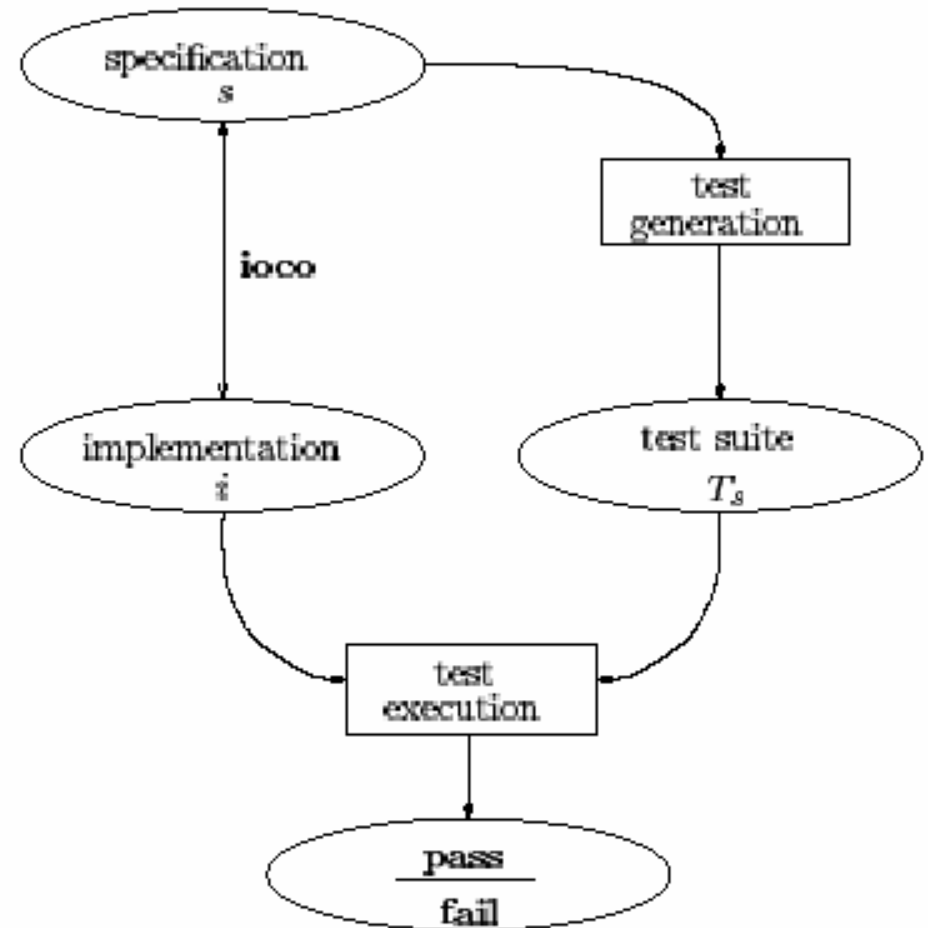 - Built upon earlier work on testing preorders

Historical reference

- Process equivalences:
  - Trace equivalence/preorder is too coarse
  - Bisimulation/simulation is too fine
- Middle ground:
  - Failures equivalence in CSP
  - may- and must-testing by Hennessy
  - Testing preorder by de Nicola
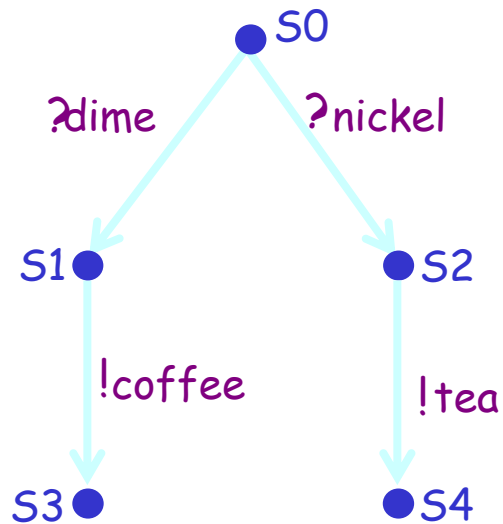 - They are all the same!
- Right notion but hard to compute

Testing architecture

- Implementation relation
- Test generation algorithm
- Test execution engine

Input-Output Transition Systems

S0

?dime   ?nickel

S1   S2

!coffee   !tea

S3   S4

dime, nickel   coffee, tea

from user to machine   from machine to user
initiative with user   initiative with machine
machine cannot refuse   user cannot refuse

input   output
L_I L_U

L_I = { ?dime, ?nickel }

L_U = { !coffee, !tea }

$$L_I \cap L_U = \varnothing \qquad L_I \cup L_U = L$$

Input-Output Transition Systems

?dime   ?nickel

?dime
?nickel

?dime
?nickel

!coffee   !tea

?dime
?nickel

?dime
?nickel

?dime
?nickel

L_I = { ?dime, ?nickel }
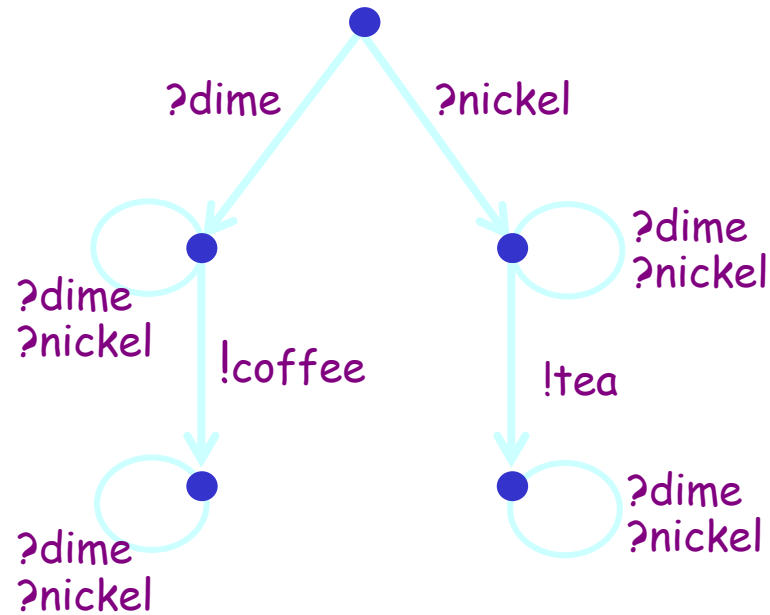
L_U = { !coffee, !tea }

Input-Output Transition Systems

$$\text{IOTS}\,(L_I, L_U) \;\subseteq\; \text{LTS}\,(L_I \cup L_U)$$

IOTS is LTS with Input-Output and always enabled inputs:

for all states s,
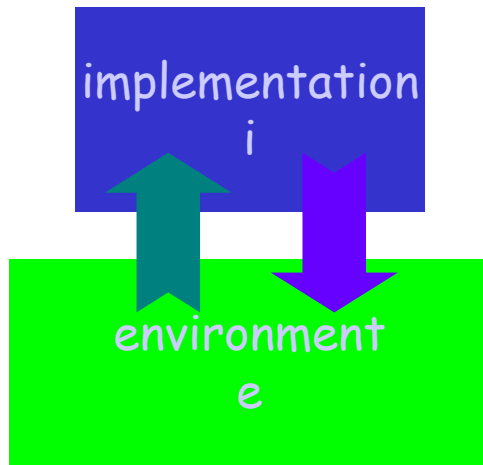
for all inputs $?a \in L_I$: $s \xRightarrow{\;?a\;}$

Preorders on IOTS

implementation i

imp

specification s

environment e

environment e

$$i \in IOTS(L_I, L_U)$$
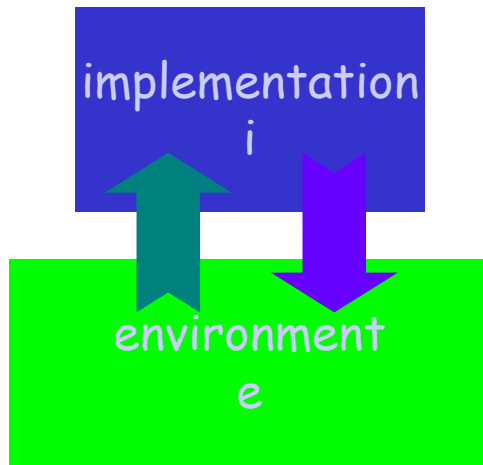
$$s \in LTS(L_I \cup L_U)$$

$$\mathbf{imp} \subseteq IOTS(L_I, L_U) \times LTS(L_I \cup L_U)$$

Observing IOTS where system inputs
interact with environment outputs, and vice versa
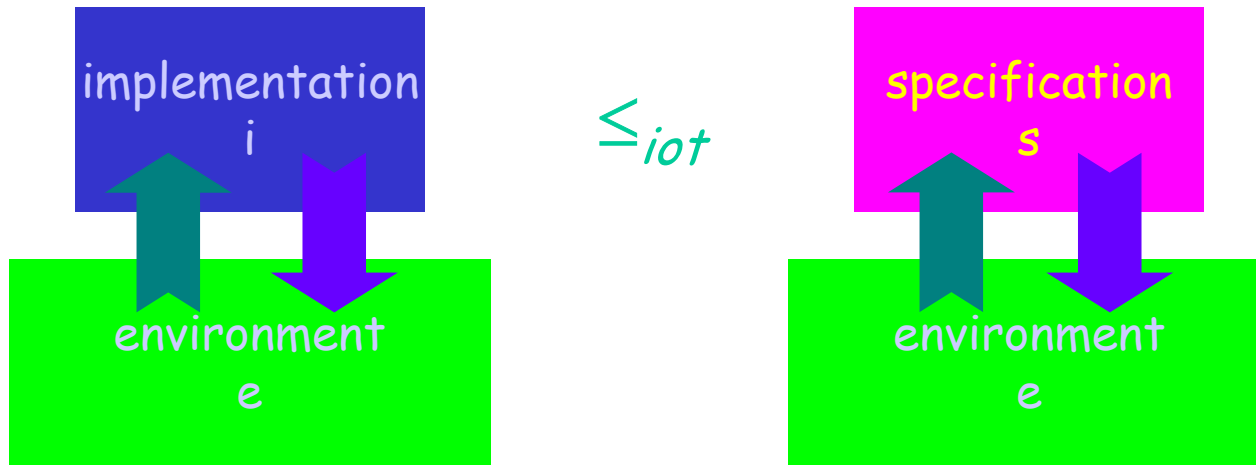
Preorders on IOTS

implementation
i

imp

specification
s

environment
e

environment
e

$i \in IOTS(L_I, L_U)$

$s \in LTS(L_I \cup L_U)$

$i \text{ imp } s \iff \forall e \in E . \text{ obs}(e, i) \subseteq \text{ obs}(e, s)$

$IOTS(L_U, L_I)$

Input-Output Testing Relation

$$i \in \mathrm{IOTS}(L_I, L_U) \qquad s \in \mathrm{LTS}(L_I \cup L_U)$$

$$i \leq_{iot} s \Leftrightarrow \forall\, e \in \mathrm{IOTS}(L_U, L_I)\, .\, \mathrm{obs}\,(\,e, i\,) \subseteq \mathrm{obs}\,(e, s\,)$$

$$\mathrm{obs}\,(\,e, p\,) \;=\; (\,\mathrm{traces}\,(e\|p\,),\; \mathrm{qtraces}\,(e\|p\,)\,)$$

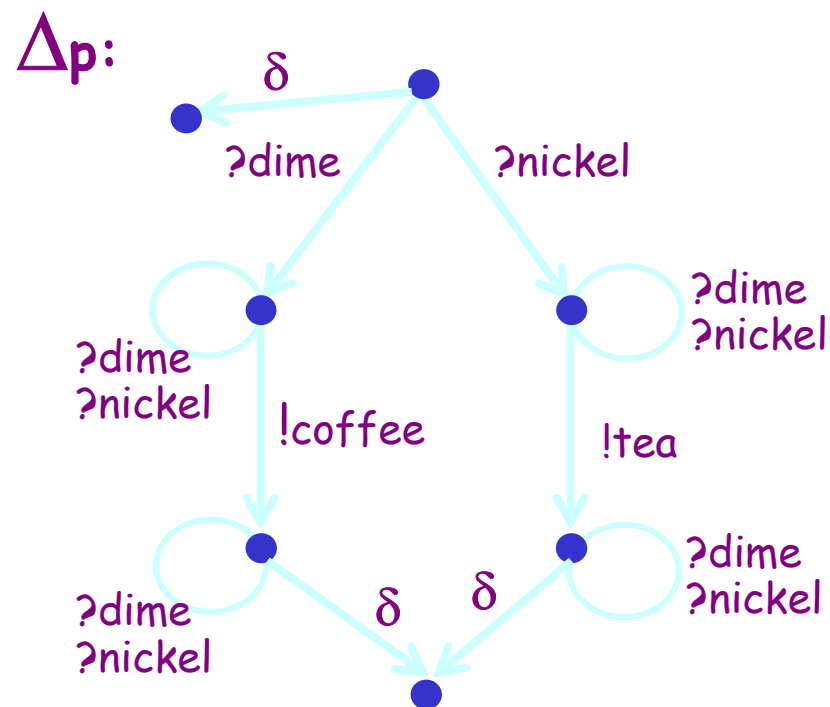$$\mathrm{qtraces}(p) = \sigma \in L^*.\ p \text{ after } \sigma \text{ refuses } L_U$$
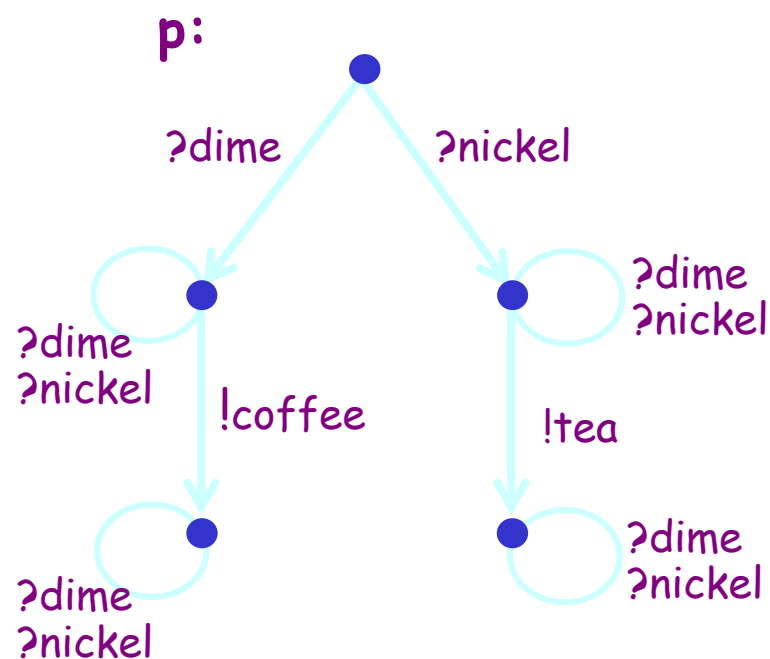
Testing preorders – a side note

- One of the reasons for using IOTS over LTS is that $\leq_{iot}$ is computationally simpler than conventional testing preorder
  - Testing preorder requires us to compare sets of pairs (trace, refusal set)
  - At the same time $\leq_{iot}$ allows us to use inclusion of weakly quiescent traces:
    - inputs can never be refused by i, and outputs can never be refused by e
    - i after σ refuses A $\Rightarrow$ A = $\varnothing$ or A = L_U

Representing quiescence

- Extend IOTS with quiescent transitions
  - deterministic δ-trace automata

p:

?dime   ?nickel

?dime
?nickel   !coffee   !tea   ?dime
?nickel

?dime
?nickel   ?dime
?nickel

Δ**p:**

δ

?dime   ?nickel

?dime
?nickel   !coffee   !tea   ?dime
?nickel

?dime
?nickel δ δ ?dime
?nickel

Conformance relation **ioconf**

- Finally…

  $i \leq_{iot} s \Leftrightarrow \forall \; \sigma \in L^*.out(\Delta i \; after \; \sigma \;) \subseteq out(\Delta s \; after \; \sigma \;)$

- Allow underspecification

  – restrict to traces of **s**

  i **ioconf** $s =_{def}$

  $\forall \sigma \in traces(\Delta s) \cap L^*.out(\Delta i \; after \; \sigma) \subseteq out(\Delta s \; after \; \sigma)$

- **ioconf**$_F$: use arbitrary F instead of traces of **s**

- Conformance relation **ioco** accounts for *repetitive* quiescence

Test cases

- A test case is a deterministic IOTS(L_U,L_I) with finite behaviors
  - Note reversed inputs and outputs
  - Do not allow choice between outputs or between input and output
- Verdict function v: S $\rightarrow$ {fail,pass}
- Test run: i **passes** t $=_{def}$
  (i||t) after σ deadlocks $\Rightarrow$
  v(t after σ)=pass

pass

?dime

fail

!coffee

!tea

pass

fail

Test generation

- Test suite T_s for a specification s is complete:
  i **ioconf** s iff $\forall t \in Ts$. i **passes** t

- Test suite T_s is sound if
  i **ioconf** $s \Rightarrow \forall t \in Ts$. i **passes** t

- Complete test suites are usually infinite
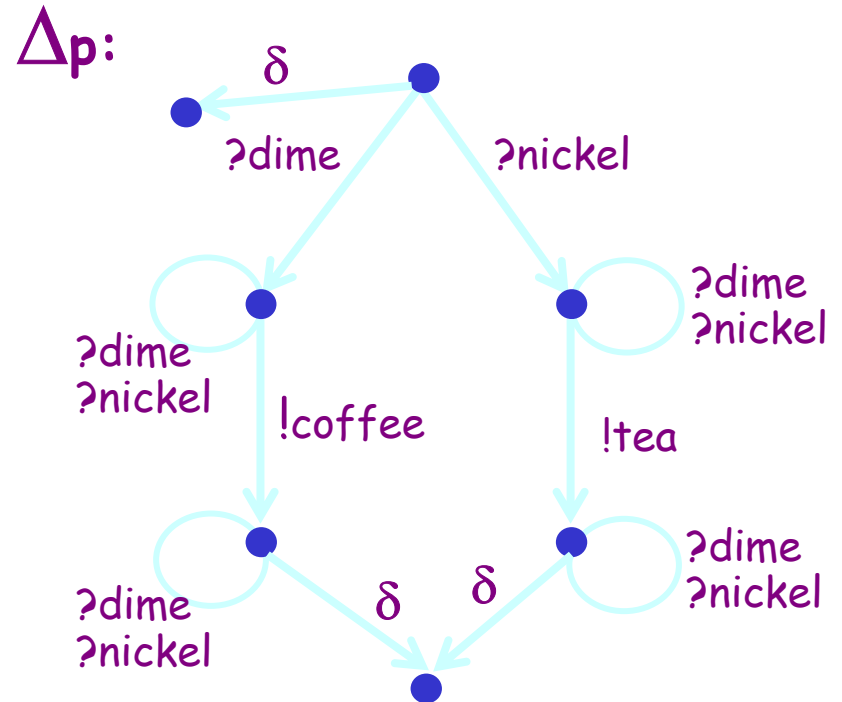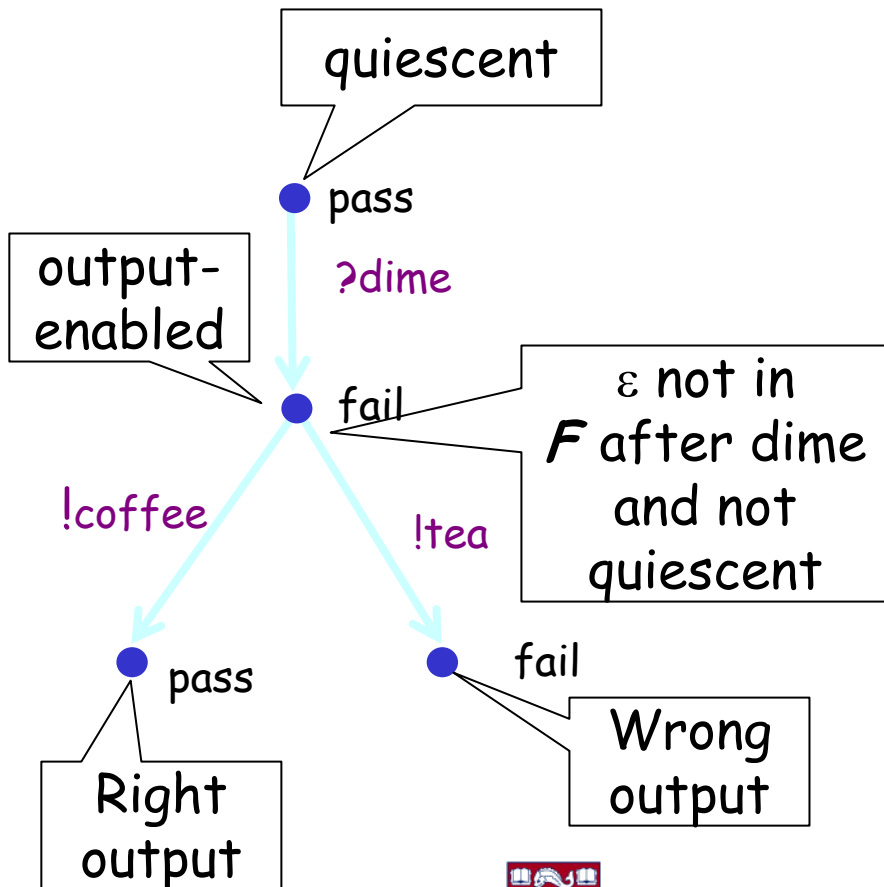
  – Aim at generating sound test suites

Test generation algorithm

- Gen(Δ_s, $\boldsymbol{F}$)
  - Choose non-deterministically:
  1. $t = \text{stop}$ and $v(t) = \text{pass}$
  2. $t = a$. Gen(Δ'_s, $\boldsymbol{F}$ after a), with $\Delta_s \xrightarrow{a} \Delta'_s$

     $v(t) = \text{pass}$
  3. $t = \sum \{x.\text{stop} \mid x \in L_U, x \notin out(\Delta_S)\} + \sum \{x.t_x \mid x \in out(\Delta_S)\}$

     $v(t) = \text{pass if } \delta \in out(\Delta_S) \vee \varepsilon \in F; \text{ otherwise fail}$

     $v(\text{stop}) = \text{fail if } \varepsilon \notin F; \text{ otherwise pass}$

Example

- **F** = dime.coffee

quiescent

pass

?dime

output-enabled

fail

ε not in **F** after dime and not quiescent

!coffee

!tea

pass

Right output

fail

Wrong output

Δp:

δ

?dime

?nickel

?dime
?nickel

?dime
?nickel

!coffee

!tea

?dime
?nickel

?dime
?nickel

δ

δ

Test purposes

- Where does *F* come from?
- Test purposes:
  - Requirements, use cases
  - Automata, message sequence charts
- Test purposes represent "interesting" or "significant" behaviors
  - Define "interesting" or "significant"…
- Can we come up with test purposes automatically?

Summary: conformance testing

- Advantages:
  - Very rigorous formal foundation
  - Size of the test suite is controlled by use cases
- Disadvantages:
  - How much have we learned about the system that passed the test suite?
  - Does not guarantee coverage

Coverage-based testing

- Traditional:
  - Tests are derived from the implementation structure (code)
- Model-driven:
  - Cover the model instead of code
  - Model should be much closer to the implementation in structure
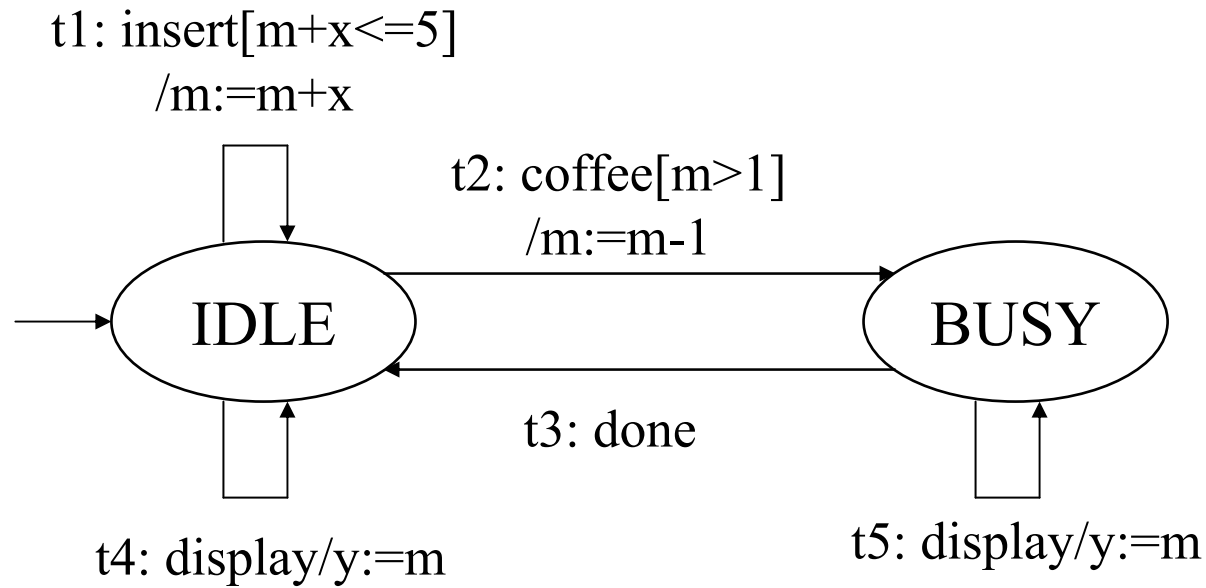- Relies on coverage criteria

Coverage criteria and tests

- [HongLeeSokolskyUral02]
- Control flow:
  - all-states
  - all-transitions
- Data flow:
  - all-defs
  - all-uses
  - all-inputs
  - all-outputs
- Test is a linear sequence of inputs and outputs

Specifications: EFSM

- Transition systems equipped with variables
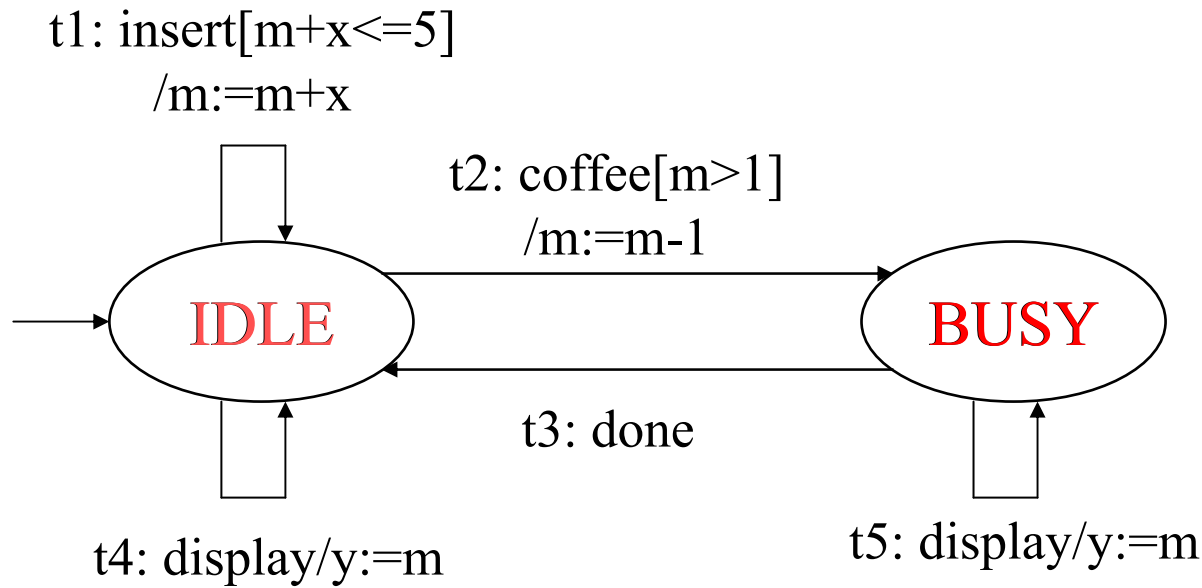- Transitions have guards and update blocks

t1: insert[m+x<=5]
   /m:=m+x

t2: coffee[m>1]
   /m:=m-1

IDLE

BUSY

t3: done

t4: display/y:=m

t5: display/y:=m

Coverage criteria

- Each coverage criterion is represented by a set of temporal logic formulas
  - WCTL: a subset of CTL
    - Atomic propositions $p_1,...,p_n$
    - Temporal operators EX, EU, EF
    - Conjunctions: at most one non-atomic conjunct
    - Negations is applied only to atomic propositions
    - Unrestricted disjunctions
    - E.g.: $EF(p_1 \& EF p_2)$
  - WCTL formulas have linear witnesses

All-states coverage criterion

- Requires every state be covered at least once
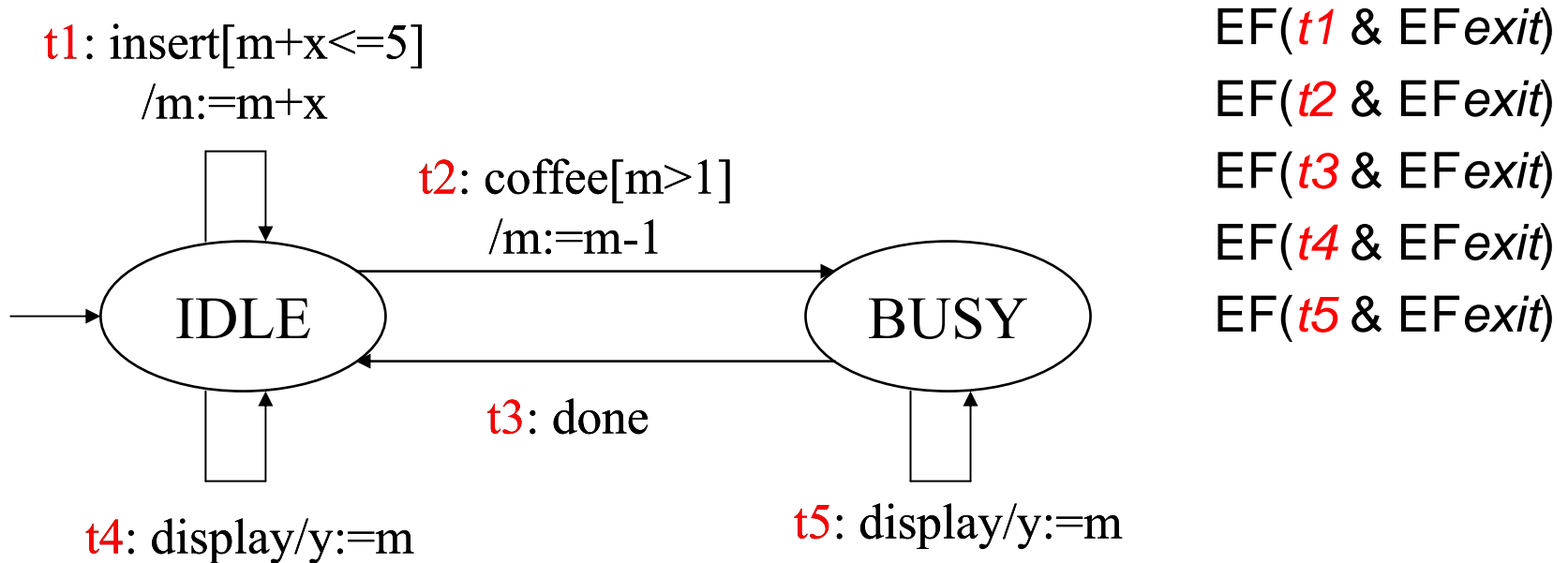- With every state s, associate EF(s & EF$exit$)

t1: insert[m+x<=5]
/m:=m+x

t2: coffee[m>1]
/m:=m-1

IDLE

BUSY

t3: done

t4: display/y:=m

t5: display/y:=m

EF(*idle* & EF*exit*)
EF(*busy* & EF*exit*)

All-transitions coverage criterion

- Requires every transition be covered at least once

- With every transition t, associate $EF(t \& EFexit)$

t1: insert[m+x<=5] /m:=m+x

t2: coffee[m>1] /m:=m-1

IDLE   BUSY

t3: done

t4: display/y:=m

t5: display/y:=m

EF($t1$ & EF$exit$)
EF($t2$ & EF$exit$)
EF($t3$ & EF$exit$)
EF($t4$ & EF$exit$)
EF($t5$ & EF$exit$)

Data flow: definitions and uses

- **Definition**: a value is assigned to a variable
- **Use**: a value of a variable is used in an expression
- Variables are defined and used in transitions
- **Definition-use pair**: (v,t,t')
  - v is defined by t
  - v is used by t'
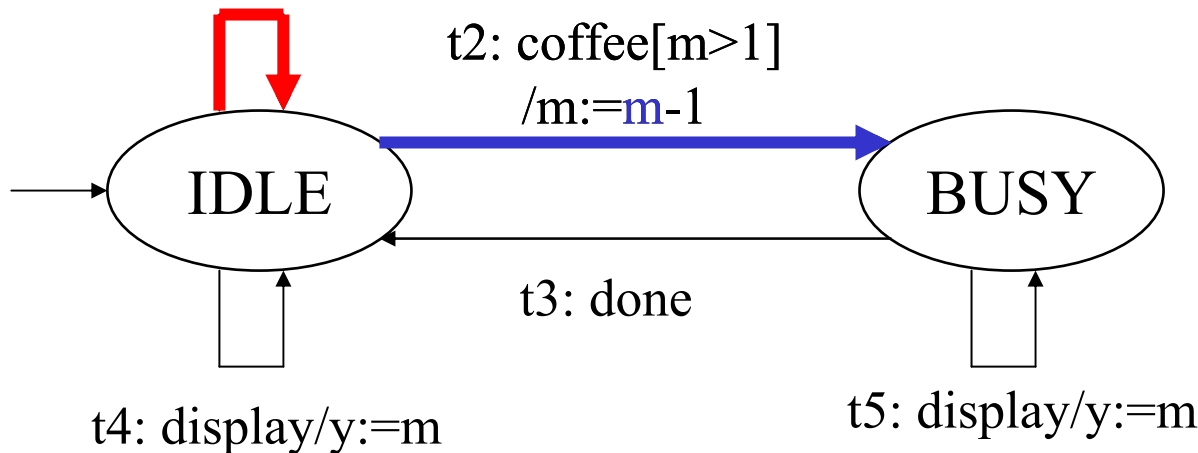 - There is a path from t to t' free from other definitions of v

Covering a du-pair

- With a du-pair (v, t, t'), *associate*
  - EF(t & EXE[!def(v) U (t' & EF$exit$)])
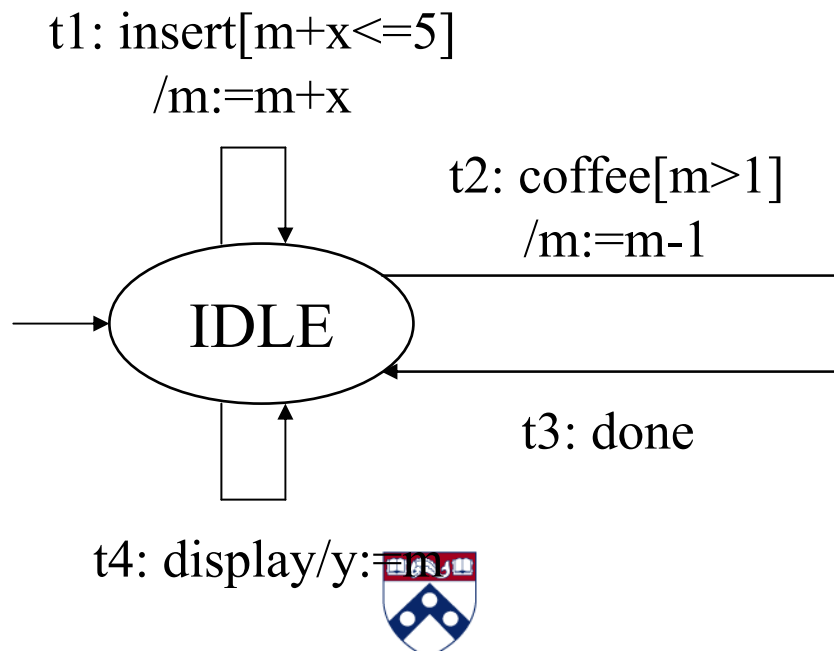 - def(v) : disjunction of all transitions that define v

t1: insert[m+x<=5]
/m:=m+x

EF($t1$ & EXE[!($t1$ | $t2$) U ($t2$ & EF$exit$)])

t2: coffee[m>1]
/m:=m-1

IDLE

BUSY

t3: done

t4: display/y:=m

t5: display/y:=m

Data-flow coverage criteria

- ## All-defs coverage criterion: a definition-clear path
  – from **every** definition to **some** use

- ## All-uses coverage criterion: a definition-clear path
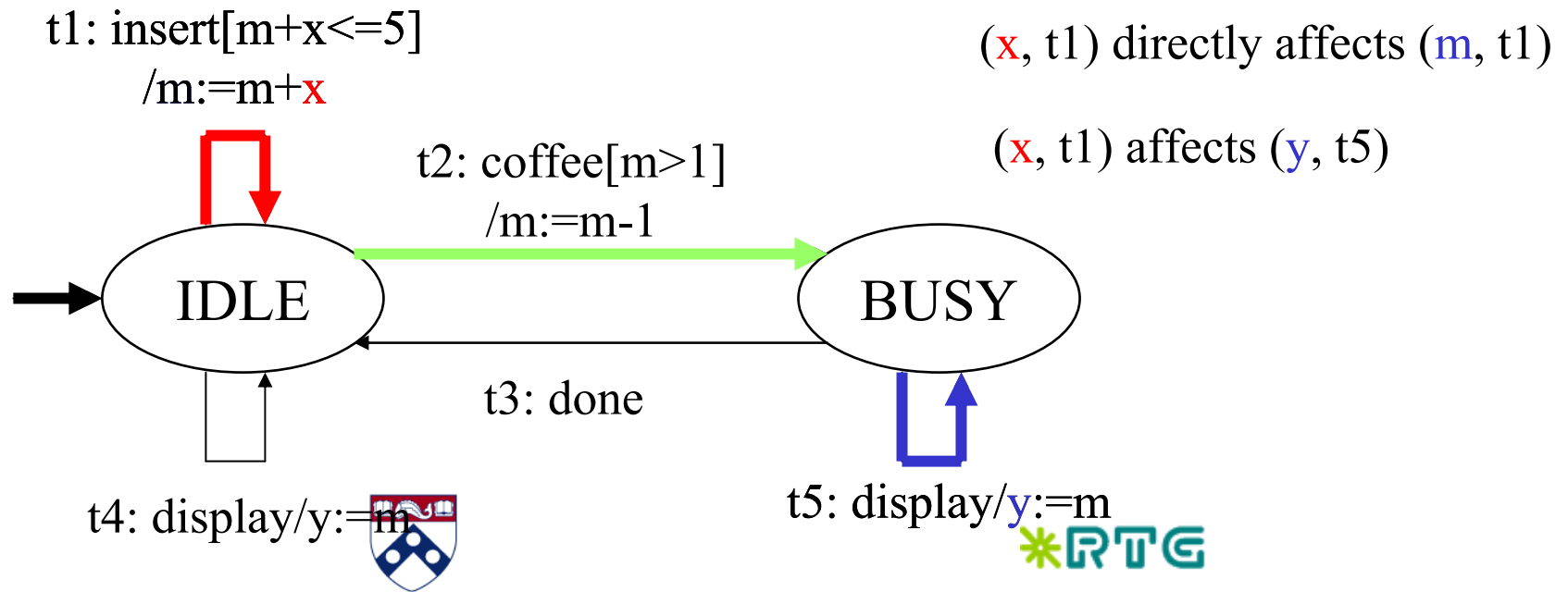 – from **every** definition to **every** use

t1: insert[m+x<=5]
/m:=m+x

t2: coffee[m>1]
/m:=m-1

IDLE

t3: done

t4: display/y:=m

All-uses coverage criterion
EF($t1$ & EXE[!($t1$ | $t2$) U ($t1$ & EF$exit$)])
EF($t1$ & EXE[!($t1$ | $t2$) U ($t2$ & EF$exit$)])
EF($t1$ & EXE[!($t1$ | $t2$) U ($t4$ & EF$exit$)])
EF($t1$ & EXE[!($t1$ | $t2$) U ($t5$ & EF$exit$)])
EF($t2$ & EXE[!($t1$ | $t2$) U ($t1$ & EF$exit$)])
EF($t2$ & EXE[!($t1$ | $t2$) U ($t2$ & EF$exit$)])
EF($t2$ & EXE[!($t1$ | $t2$) U ($t4$ & EF$exit$)])
EF($t2$ & EXE[!($t1$ | $t2$) U ($t5$ & EF$exit$)])
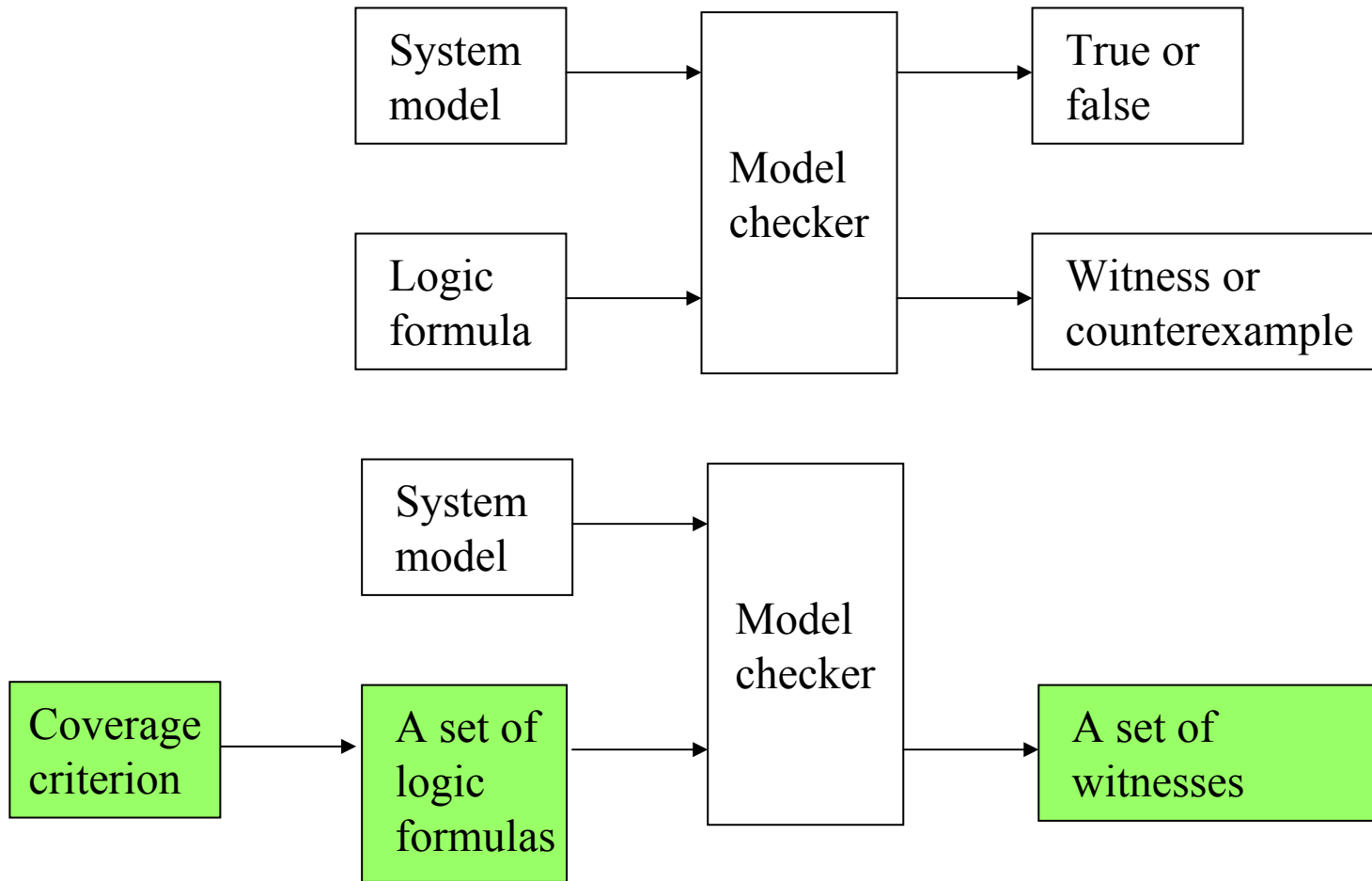
Data flow chains

- Affect pair (v, t, v', t'): the value of v used by t affects the value of v' defined at t'
  - Either $t=t'$ ((v,t) directly affects (v',t')) or
  - there is a du-pair (v'',t,t'') s.t. (v,t) directly affects (v'',t) and (v'',t'') affects (v',t')

t1: insert[m+x<=5]
  /m:=m+x

t2: coffee[m>1]
  /m:=m-1

IDLE

BUSY

t3: done

t4: display/y:=m

t5: display/y:=m

(x, t1) directly affects (m, t1)

(x, t1) affects (y, t5)

Data flow chain coverage

- Affect pair (v, t, v', t')
  - May consist of an arbitrary number of definition-use pairs
  - We extend CTL with least fixpoint operators
    - Alternatively, we can use (alternation-free) mu-calculus
- All-inputs coverage criterion
  - Requires a path from **every** input to **some** output be covered at least once
- All-outputs coverage criterion
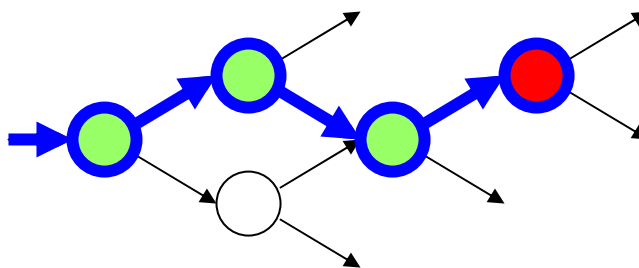 - Requires a path from **every** input to **every**

Test Generation

Test Generation

- Generating a witness for a formula
  - Cost: the length of a witness
  - A minimal-cost witness for a formula
    - Existing model checkers generate a minimal-cost witness by breadth-first search of state space
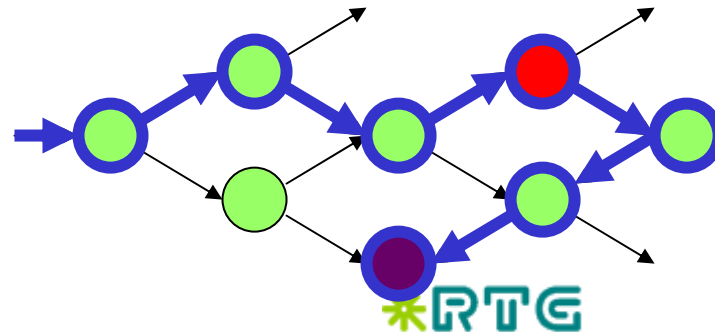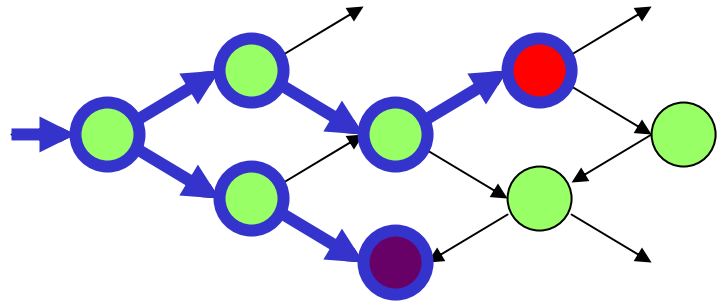
E[● U ●]

Test Generation

- Costs
  - The total length of witnesses or
  - The number of witnesses
- Both optimization problems are NP-hard

$E[\, \bigcirc \, U \, \bullet \,]$

$E[\, \bigcirc \, U \, \bullet \,]$

$E[\, \bigcirc \, U \, \bullet \,]$

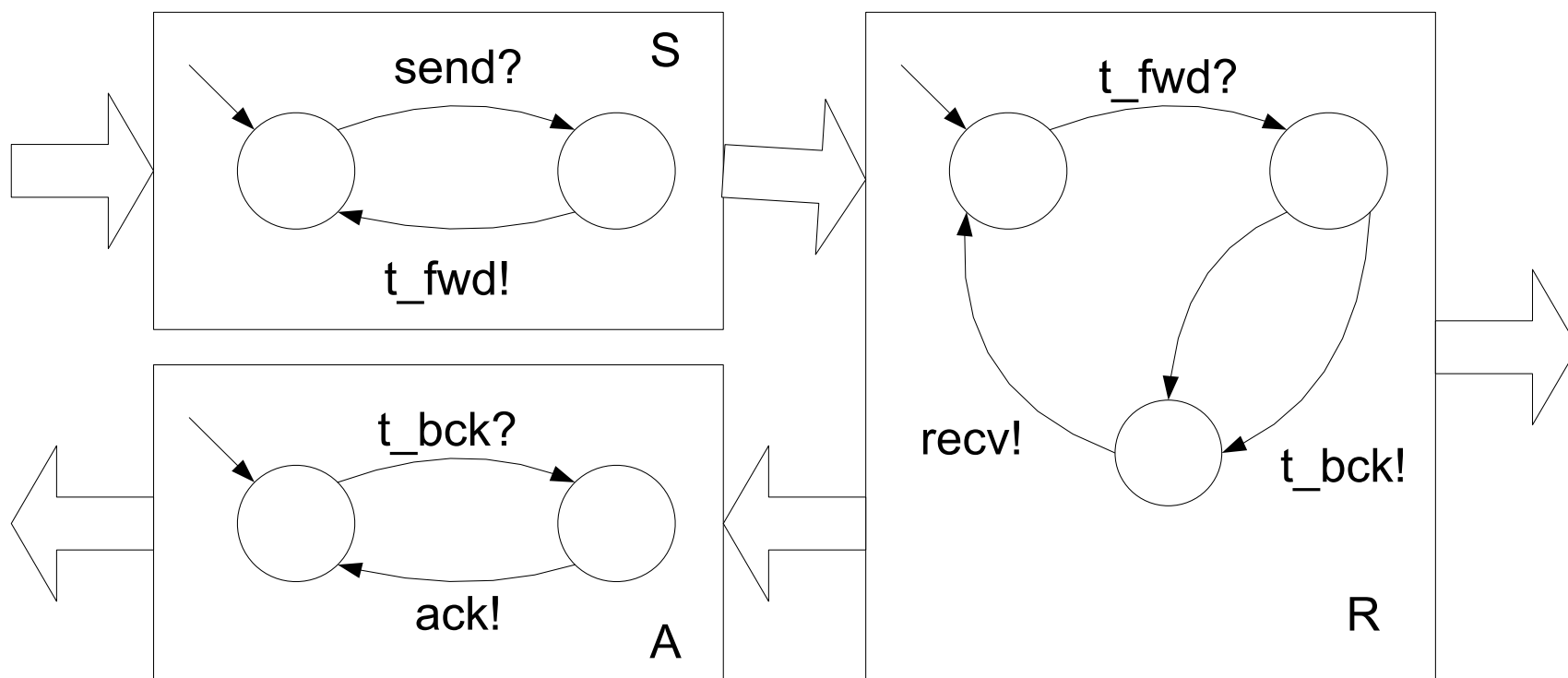$E[\, \bigcirc \, U \, \bullet \,]$

Coverage for distributed systems

- What if our system is a collection of components?
- Possible solutions:
  - Generate tests for each components
    - Clearly unsatisfactory; does not test integration
  - Generate tests from the product of component models
    - Too many redundant tests
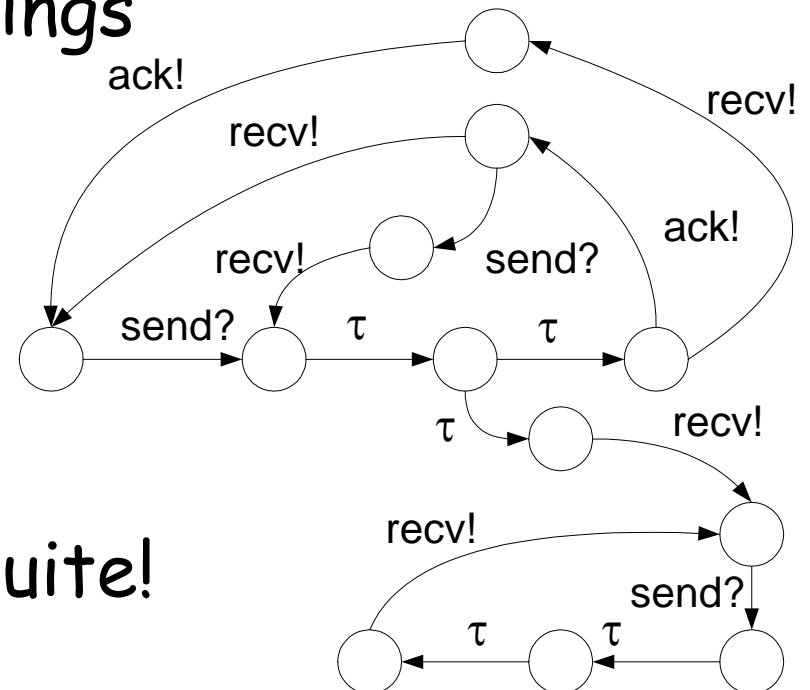- Non-determinism is another problem

Example

- Producer-consumer with acknowledgements

Covering product transition system

- Linear tests bring trouble:

  send?.ack!.recv!

  – May fail if the system chooses a different path

- Tests differ in interleavings of independent events

  – No need to test
    send?.ack!.recv!
    send?.recv!.ack!
    separately

- State explosion in test suite!

Partial orders for test generation

- Use event structures instead of transition systems [Heninger97]
- Test generation covers the event structure
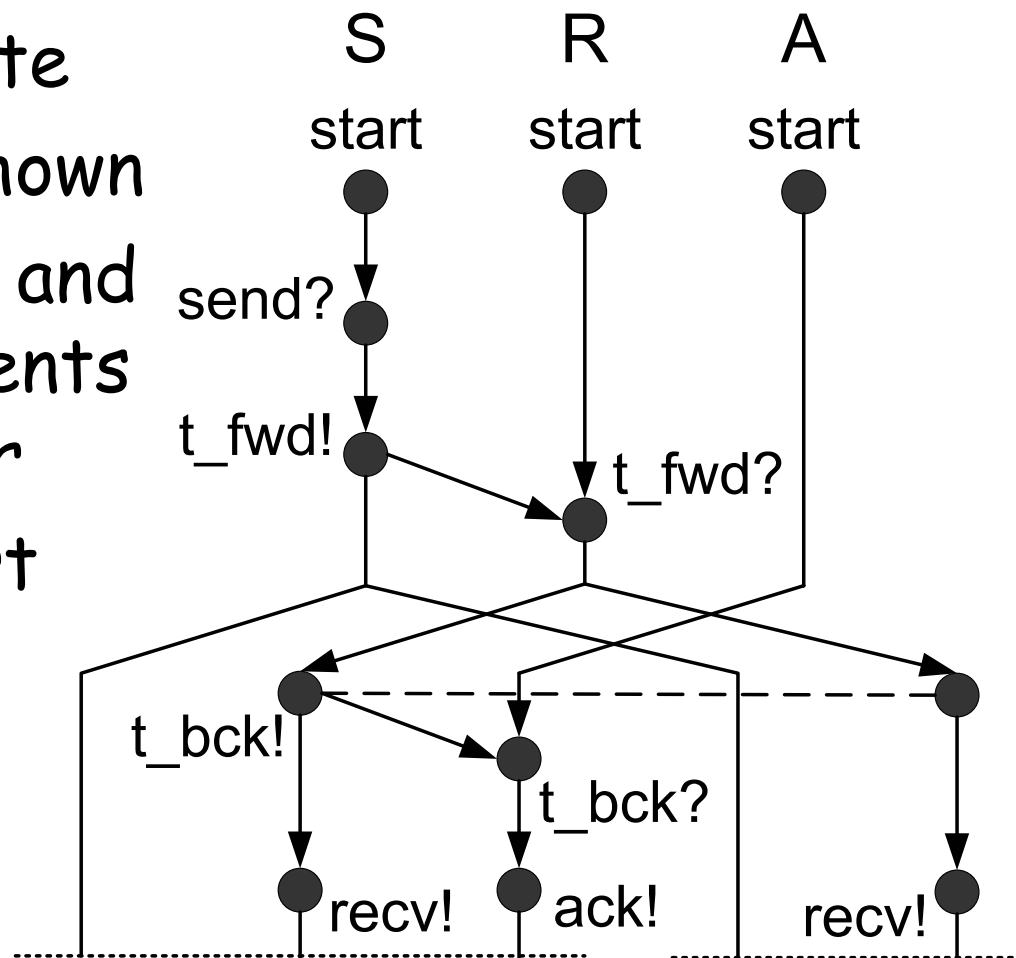- Allows natural generation of distributed testers

Prime event structures (PES)

- Set of events E

    - Events are occurrences of actions

- Causality relation $\quad \prec\, \subseteq E \times E$

    - Partial order

- Conflict relation $\quad \# \subseteq E \times E$

    - irreflexive and symmetric

- Labeling function $\quad l : E \to A$

- Finite causality $\qquad \{e' | e' \prec e\}$ is finite

- Conflict inheritance $\qquad e \# e' \wedge e' \prec e'' \Rightarrow e \# e''$
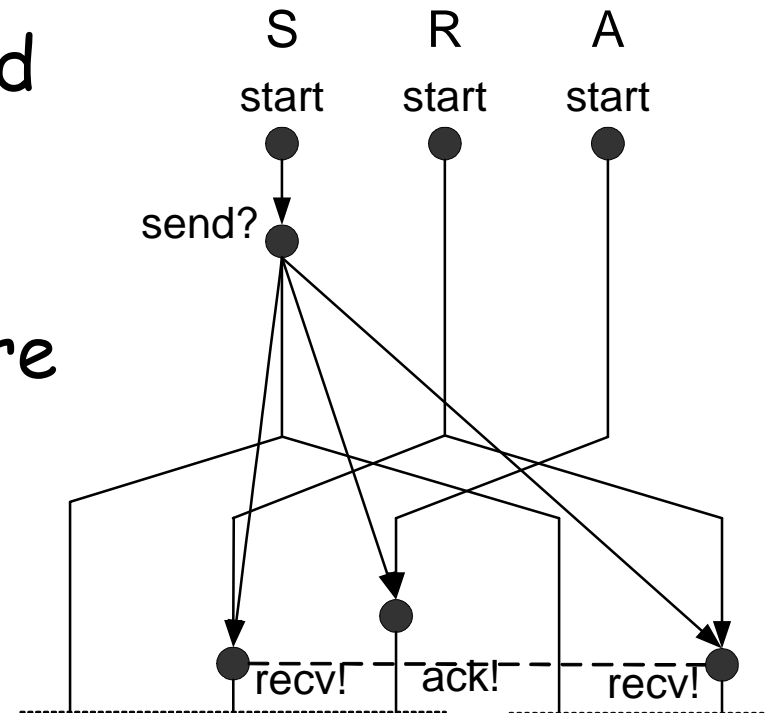
Producer-consumer PES

- Structure is infinite
  - Initial part is shown
- Causally unrelated and non-conflicting events can occur together
- Behaviors will start repeating
  - Can stop with finite structure

Test generation with PES

- Project PES on observable actions, propagating conflicts

- Every path in the PES should be covered

- Tests consist of distributed testers with coordination messages between tests

  - Coordination messages are inserted when there is a causal edge between locations

S   R   A

start   start   start

send?

recv!   ack!   recv!

Summary: coverage-based testing

- Advantages:
  - Exercise the specification to the desired degree
  - Does not rely on test purpose selection
- Disadvantages:
  - Large and unstructured test suites
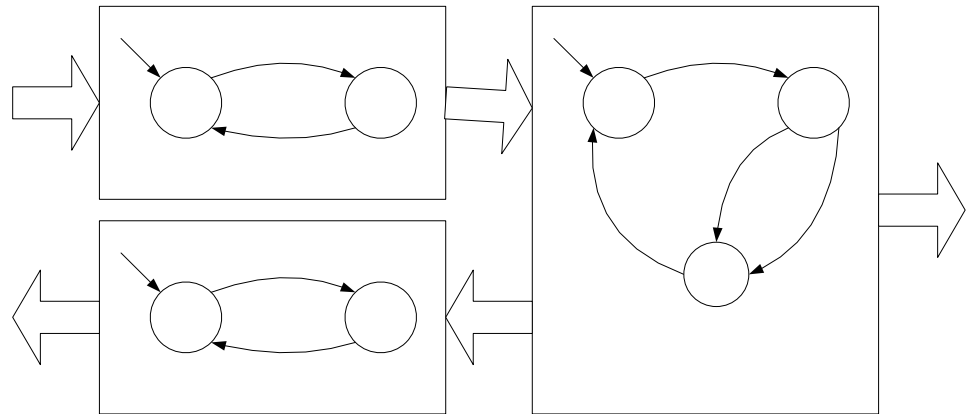 - If the specification is an overapproximation, tests may be infeasible

Generation of test purposes

- Recent work: [HenningerLuUral-03]
- Construct PES
- Generate MSC (message sequence charts) to cover PES
- Use MSC as test purposes in **ioco**-based test generation

Controllability of testing

- Conformance testing may not provide enough guarantees
  - With branching tests, test purpose behavior may be avoided
  - What if I never see ack?

- Problem: inherent uncertainty in the system

How to contain uncertainty?

- Avoidance (no need to increase control)
  - During testing, compute confidence measure
    - E. g., accumulate coverage
  - Stop at the desired confidence level
- Prevention (add more control)
  - Use instrumentation to resolve uncertainty
  - What to instrument?
    - Use model for guidance
- Anyone needs a project to work on?

✱RTG

Test generation tools

- TorX
  - Based on **ioco**
  - On-the-fly test generation and execution
  - Symbolic testing (data parameterization)
  - LOTOS, Promela, …
  - http://fmt.cs.utwente.nl/tools/torx/
- TGV
  - Based on symbolic **ioconf**
  - LOTOS, SDL, UML
  - http://www.irisa.fr/pampa/VALIDATION/TGV/TGV.html